

Mathematical induction and universal quantification

In the following, x and x' range over the same domain; similarly, y and y' range over the same domain. Moreover, $<$ is a well-founded relation on the domain of x , i.e. for any P

$$(0) \quad (\underline{\forall}x :: P.x) \equiv (\underline{\forall}x :: P.x \Leftarrow (\underline{\forall}x' : x' < x : P.x')) \quad .$$

Consider now the obligation to prove for some R

$$(1) \quad (\underline{\forall}x,y :: R.x.y) \quad .$$

Two approaches are now possible.

(i) (1)

$$\begin{aligned} &= \{\text{nesting}\} \\ &(\underline{\forall}y :: (\underline{\forall}x :: R.x.y)) \\ &= \{(0) \text{ with } P.x := R.x.y; \text{ unnesting}\} \\ &(\underline{\forall}x,y :: R.x.y \Leftarrow (\underline{\forall}x' : x' < x : R.x'.y)) \end{aligned}$$

(ii) (1)

$$\begin{aligned} &= \{\text{nesting}\} \\ &(\underline{\forall}x :: (\underline{\forall}y :: R.x.y)) \\ &= \{(0) \text{ with } P.x := (\underline{\forall}y :: R.x.y); \text{ renaming dummy}\} \\ &(\underline{\forall}x :: (\underline{\forall}y :: R.x.y) \Leftarrow (\underline{\forall}x' : x' < x : (\underline{\forall}y' :: R.x'.y'))) \\ &= \{\Leftarrow Q \text{ distributes over } \underline{\forall}; \text{ unnesting}\} \\ &(\underline{\forall}x,y :: R.x.y \Leftarrow (\underline{\forall}x' : x' < x : (\underline{\forall}y' :: R.x'.y'))) \end{aligned}$$

Method (i) gives a shorter formula, but ends with a stronger proof obligation than method (ii),

which, therefore, is more powerful. The message was driven home to me at last week's session of the ETAC, where we considered for $R.x.y$

$$\neg x < y \vee \neg y < x$$

Approach (i) leads one into a dead end; approach (ii) leads one to observe for any x,y

$$\begin{aligned}
 & (\exists x': x' < x : (\exists y' : \neg x' < y' \vee \neg y' < x')) \\
 = & \quad \{ \text{trading, unnesting} \} \\
 & (\exists x', y' : \neg x' < x \vee \neg x' < y' \vee \neg y' < x') \\
 \Rightarrow & \quad \{ \text{instantiate with } x', y' := y, x \} \\
 & \neg y < x \vee \neg y < x \vee \neg x < y \\
 = & \quad \{ \text{pred. calc.} \} \\
 & \neg x < y \vee \neg y < x .
 \end{aligned}$$

Nuenen, 10 July 1990

prof. dr. Edsger W. Dijkstra
 Department of Computer Sciences
 The University of Texas at Austin
 Austin, TX 78712-1188
 USA