EVENT: Start with the library "c2".

```
;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;
;;                                                                            ;;
;;                    PREFIX ACCESSOR FUNCTIONS                               ;;
;;                                                                            ;;
;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;
```

```
;; These functions simply allow the accessing of various components
;; of Micro-Gypsy structures.
```

DEFINITION: loop-body (*stmt*) = cadr (*stmt*)

DEFINITION: prog2-left-branch (*stmt*) = cadr (*stmt*)

DEFINITION: prog2-right-branch (*stmt*) = caddr (*stmt*)

DEFINITION: signalled-condition (*stmt*) = cadr (*stmt*)

THEOREM: signalled-condition-expansion2
signalled-condition (cons (*signal*, *lst*)) = car (*lst*)

DEFINITION: if-condition (*stmt*) = cadr (*stmt*)

DEFINITION: if-true-branch (*stmt*) = caddr (*stmt*)

DEFINITION: if-false-branch (*stmt*) = cadddr (*stmt*)

DEFINITION: begin-body (*stmt*) = cadr (*stmt*)

DEFINITION: when-labels (*stmt*) = caddr (*stmt*)

DEFINITION: when-handler (*stmt*) = cadddr (*stmt*)

DEFINITION: call-name (*stmt*) = cadr (*stmt*)

THEOREM: call-name-expansion
call-name (cons (*proc-call*, cons (*name*, *y*))) = *name*

DEFINITION: call-actuals (*stmt*) = caddr (*stmt*)

DEFINITION: call-conds (*stmt*) = cadddr (*stmt*)

DEFINITION: formal-type (*exp*) = cadr (*exp*)

DEFINITION: formal-initial-value (*local*) = caddr (*local*)

```
;; current count 13
```

DEFINITION:
PREDEFINED-PROCEDURE-LIST
= '(mg-simple-variable-assignment
mg-simple-constant-assignment mg-simple-variable-eq
mg-simple-constant-eq mg-integer-le
mg-integer-unary-minus mg-integer-add
mg-integer-subtract mg-boolean-or mg-boolean-and
mg-boolean-not mg-index-array mg-array-element-assignment)

```
;; The procedure list is an alist of the form
;; ( ... (name proc-definition) ... ).  Notice that this will change when
;; I return to prefix form since the name is not in the car position in that
;; representation.
```

DEFINITION:
predefined-procp (*name*) = (*name* ∈ PREDEFINED-PROCEDURE-LIST)

DEFINITION:
user-defined-procp (*name*, *proc-list*)
= ((*name* ∉ PREDEFINED-PROCEDURE-LIST) ∧ definedp (*name*, *proc-list*))

DEFINITION:
defined-procp (*name*, *proc-list*)
= (predefined-procp (*name*) ∨ user-defined-procp (*name*, *proc-list*))

DEFINITION: fetch-def (*name*, *proc-list*) = assoc (*name*, *proc-list*)

DEFINITION:
fetch-called-def (*stmt*, *proc-list*) = fetch-def (call-name (*stmt*), *proc-list*)

```
;; (name data-formals cond-formals data-locals cond-locals body)
```

DEFINITION: def-name (*def*) = car (*def*)

EVENT: Disable def-name.

DEFINITION: def-formals (*def*) = cadr (*def*)

EVENT: Disable def-formals.

DEFINITION: def-conds (*def*) = caddr (*def*)

EVENT: Disable def-conds.

DEFINITION: def-locals (*def*) = cadddr (*def*)

EVENT: Disable def-locals.

DEFINITION: def-cond-locals (*def*) = caddddr (*def*)

EVENT: Disable def-cond-locals.

DEFINITION: def-body (*def*) = cadddddr (*def*)

EVENT: Disable def-body.

DEFINITION: array-elemtype (*type*) = cadr (*type*)

EVENT: Disable array-elemtype.

DEFINITION: array-length (*type*) = caddr (*type*)

EVENT: Disable array-length.

```
;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;
;;                                                                              ;;
;;                              THE RECOGNIZER                                  ;;
;;                                                                              ;;
;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;
```

DEFINITION:
RESERVED-NAMES-LIST = `'(leave normal routineerror)`

DEFINITION: reserved-word (*wd*) = (*wd* ∈ RESERVED-NAMES-LIST)

DEFINITION:
ok-mg-namep (*ident*)
= (litatom (*ident*)
∧ (45 ∉ unpack (*ident*))
∧ (¬ reserved-word (*ident*)))

EVENT: Disable ok-mg-namep.

```
;; Notice that the word size is selected to be exactly that of
;; Piton.  This eliminates some of the difficulties of the mappings.
```

DEFINITION: MG-WORD-SIZE = 32

EVENT: Introduce the function symbol *mg-max-ctrl-stk-size* of 0 arguments.

AXIOM: mg-max-ctrl-stk-size-small-naturalp
(MG-MAX-CTRL-STK-SIZE $\in$ **N**)
$\wedge$ (MG-MAX-CTRL-STK-SIZE $<$ exp (2, MG-WORD-SIZE))

EVENT: Introduce the function symbol *mg-max-temp-stk-size* of 0 arguments.

AXIOM: mg-max-temp-stk-size-numberp
(MG-MAX-TEMP-STK-SIZE $\in$ **N**)
$\wedge$ (MG-MAX-TEMP-STK-SIZE $<$ exp (2, MG-WORD-SIZE))

```
;; Notice also the I can simply use the Piton function
;; small-integerp rather than define a specific MG notion of what
;; is an acceptable integer.
```

DEFINITION: MAXINT = (exp (2, MG-WORD-SIZE $-$ 1) $-$ 1)

DEFINITION: MININT = ($-$ exp (2, MG-WORD-SIZE $-$ 1))

```
;; An integer literal is of the form (INT-MG n) where n is in the
;; range (minint..maxint).
```

DEFINITION:
int-literalp (*exp*)
= ((`'int-mg` = car (*exp*))
$\wedge$ length-plistp (*exp*, 2)
$\wedge$ small-integerp (cadr (*exp*), MG-WORD-SIZE))

DEFINITION:
boolean-literalp (*exp*)
= ((`'boolean-mg` = car (*exp*))
$\wedge$ length-plistp (*exp*, 2)
$\wedge$ (cadr (*exp*) $\in$ `'(true-mg false-mg)`))

DEFINITION:
character-literalp (*exp*)
= (('character-mg = car (*exp*))
∧ length-plistp (*exp*, 2)
∧ (cadr (*exp*) ∈ **N**)
∧ (cadr (*exp*) ≤ 127))

EVENT: Disable int-literalp.

EVENT: Disable boolean-literalp.

EVENT: Disable character-literalp.

DEFINITION:
simple-mg-type-refp (*typref*)
= (*typref* ∈ '(int-mg boolean-mg character-mg))

DEFINITION:
array-mg-type-refp (*typref*)
= (length-plistp (*typref*, 3)
∧ (car (*typref*) = 'array-mg)
∧ simple-mg-type-refp (array-elemtype (*typref*))
∧ (array-length (*typref*) ≠ 0))

EVENT: Disable array-mg-type-refp.

DEFINITION:
mg-type-refp (*typref*)
= (simple-mg-type-refp (*typref*) ∨ array-mg-type-refp (*typref*))

DEFINITION:
simple-typed-literalp (*lit*, *type*)
= **if** *type* = 'int-mg **then** int-literalp (*lit*)
elseif *type* = 'boolean-mg **then** boolean-literalp (*lit*)
elseif *type* = 'character-mg **then** character-literalp (*lit*)
else f endif

DEFINITION:
simple-typed-literal-plistp (*lst*, *type*)
= **if** *lst* ≃ **nil** **then** *lst* = **nil**
else simple-typed-literalp (car (*lst*), *type*)
∧ simple-typed-literal-plistp (cdr (*lst*), *type*) **endif**

DEFINITION:
array-literalp (*exp*, *length*, *elemtype*)
= (simple-typed-literal-plistp (*exp*, *elemtype*)
∧ (length (*exp*) = *length*))

EVENT: Disable array-literalp.

DEFINITION:
ok-mg-array-value (*exp*, *type*)
= array-literalp (*exp*, array-length (*type*), array-elemtype (*type*))

EVENT: Disable ok-mg-array-value.

DEFINITION:
ok-mg-valuep (*exp*, *type*)
= **if** simple-mg-type-refp (*type*) **then** simple-typed-literalp (*exp*, *type*)
elseif array-mg-type-refp (*type*) **then** ok-mg-array-value (*exp*, *type*)
else f endif

EVENT: Disable ok-mg-valuep.

```
;; The recognizer has a structure called the name-alist of the following
;; form:
;; (... (name type) ...)
;; This allows the identification of the types of variables.  It should be
;; the case that the values of the variables in the meaning alist correspond
;; to their types on the name-alist.
```

DEFINITION: m-type (*x*) = cadr (*x*)

DEFINITION: get-m-type (*name*, *alist*) = m-type (assoc (*name*, *alist*))

DEFINITION:
has-array-type (*name*, *alist*)
= (car (get-m-type (*name*, *alist*)) = `'array-mg`)

DEFINITION:
mg-name-alist-elementp (*x*)
= (ok-mg-namep (car (*x*)) ∧ mg-type-refp (m-type (*x*)))

DEFINITION:
mg-name-alistp (*alist*)
= **if** *alist* ≃ **nil** **then** *alist* = **nil**
else mg-name-alist-elementp (car (*alist*))
∧ mg-name-alistp (cdr (*alist*)) **endif**

DEFINITION: identifierp (*name*) = ok-mg-namep (*name*)

DEFINITION:
defined-identifierp (*name*, *alist*)
= (identifierp (*name*) ∧ definedp (*name*, *alist*))

```
;; I'm restricting the if-condition to be a boolean identifier rather
;; than a boolean expression just so that the compilation will take a
;; fixed number of steps.  That is, I want every expression within the
;; code to be a reference to a variable.  This can obviously be relaxed
;; but is general since I allow the initialization of locals.  Thus,
;; any literal values can be stored in local variables.

;; >>> These should probably be removed in favor of one function with a
;;     type argument.
```

DEFINITION:
boolean-identifierp (*name*, *alist*)
= (identifierp (*name*) ∧ (get-m-type (*name*, *alist*) = `'boolean-mg`))

DEFINITION:
int-identifierp (*name*, *alist*)
= (identifierp (*name*) ∧ (get-m-type (*name*, *alist*) = `'int-mg`))

DEFINITION:
character-identifierp (*name*, *alist*)
= (identifierp (*name*) ∧ (get-m-type (*name*, *alist*) = `'character-mg`))

DEFINITION:
array-identifierp (*name*, *alist*)
= (defined-identifierp (*name*, *alist*) ∧ has-array-type (*name*, *alist*))

EVENT: Disable boolean-identifierp.

EVENT: Disable int-identifierp.

EVENT: Disable character-identifierp.

EVENT: Disable array-identifierp.

DEFINITION:
simple-identifierp (*name*, *alist*)

= (boolean-identifierp (*name*, *alist*)
∨ int-identifierp (*name*, *alist*)
∨ character-identifierp (*name*, *alist*))

DEFINITION:
simple-typed-identifierp (*ident*, *type*, *alist*)
= **if** *type* = `'int-mg` **then** int-identifierp (*ident*, *alist*)
elseif *type* = `'boolean-mg` **then** boolean-identifierp (*ident*, *alist*)
elseif *type* = `'character-mg`
then character-identifierp (*ident*, *alist*)
else f endif

THEOREM: int-identifierp-simple
int-identifierp (*x*, *alist*) → simple-identifierp (*x*, *alist*)

THEOREM: boolean-identifierp-simple
boolean-identifierp (*x*, *alist*) → simple-identifierp (*x*, *alist*)

THEOREM: character-identifierp-simple
character-identifierp (*x*, *alist*) → simple-identifierp (*x*, *alist*)

EVENT: Disable simple-identifierp.

EVENT: Disable simple-typed-identifierp.

THEOREM: int-identifierp-implies-definedp
int-identifierp (*x*, *alist*) → definedp (*x*, *alist*)

EVENT: Disable int-identifierp-implies-definedp.

THEOREM: boolean-identifierp-implies-definedp
boolean-identifierp (*x*, *alist*) → definedp (*x*, *alist*)

EVENT: Disable boolean-identifierp-implies-definedp.

THEOREM: character-identifierp-implies-definedp
character-identifierp (*x*, *alist*) → definedp (*x*, *alist*)

EVENT: Disable character-identifierp-implies-definedp.

THEOREM: simple-identifierp-implies-definedp
simple-identifierp (*x*, *alist*) → definedp (*x*, *alist*)

EVENT: Disable simple-identifierp-implies-definedp.

THEOREM: simple-typed-identifierp-implies-definedp
simple-typed-identifierp $(x,\ type,\ alist) \rightarrow$ definedp $(x,\ alist)$

EVENT: Disable simple-typed-identifierp-implies-definedp.

THEOREM: array-identifierp-implies-definedp
array-identifierp $(x,\ alist) \rightarrow$ definedp $(x,\ alist)$

EVENT: Disable array-identifierp-implies-definedp.

```
;; I have decided to make references to individual array elements impossible.
;; That is, a user program can only pass a whole array to a subroutine.  Arrays
;; are viewed as abstract data types and the only accesses to them are via
;; the predefined functions which select elements and write elements.  This will
;; guarantee that the param list of a routine is unchanged from the non-structured
;; situation.
```

DEFINITION:
identifier-plistp (lst)
$=$ **if** $lst \simeq$ **nil** **then** $lst =$ **nil**
else identifierp (car (lst)) $\wedge$ identifier-plistp (cdr (lst)) **endif**

THEOREM: identifier-plistp-plistp
identifier-plistp $(x) \rightarrow$ plistp (x)

EVENT: Disable identifier-plistp-plistp.

THEOREM: identifier-plistp-distributes
(identifier-plistp $(x) \wedge$ identifier-plistp (y))
$\rightarrow$ identifier-plistp (append $(x,\ y)$)

THEOREM: leave-not-in-identifier-plistp
identifier-plistp $(lst) \rightarrow$ (`'leave` $\notin lst$)

EVENT: Disable leave-not-in-identifier-plistp.

DEFINITION:
cond-identifierp $(x,\ cond\text{-}list)$
$=$ $((x =$ `'routineerror`$) \vee$ (identifierp $(x) \wedge (x \in cond\text{-}list)))$

EVENT: Disable cond-identifierp.

DEFINITION:
cond-identifier-plistp (*lst*, *cond-list*)
= **if** *lst* ≃ **nil** **then** *lst* = **nil**
else cond-identifierp (car (*lst*), *cond-list*)
∧ cond-identifier-plistp (cdr (*lst*), *cond-list*) **endif**

EVENT: Disable cond-identifier-plistp.

THEOREM: cond-identifier-plistp-cond-subsetp
cond-identifier-plistp (*lst*, *cond-list*) → cond-subsetp (*lst*, *cond-list*)

EVENT: Disable cond-identifier-plistp-cond-subsetp.

THEOREM: normal-not-in-cond-identifier-plistp
cond-identifier-plistp (*lst*, *cond-list*) → (`'normal` ∉ *lst*)

EVENT: Disable normal-not-in-cond-identifier-plistp.

THEOREM: adding-element-preserves-cond-identifier-plistp
cond-identifier-plistp (*y*, *cond-list*)
→ cond-identifier-plistp (*y*, cons (*x*, *cond-list*))

EVENT: Disable adding-element-preserves-cond-identifier-plistp.

THEOREM: superset-preserves-cond-identifier-plistp
(subset (*x*, *y*) ∧ cond-identifier-plistp (*lst*, *x*))
→ cond-identifier-plistp (*lst*, *y*)

EVENT: Disable superset-preserves-cond-identifier-plistp.

DEFINITION:
nonempty-cond-identifier-plistp (*lst*, *cond-list*)
= (cond-identifier-plistp (*lst*, *cond-list*) ∧ (*lst* ≠ **nil**))

DEFINITION:
ok-condition (*exp*, *cond-list*)
= ((*exp* = `'routineerror`)
∨ ((ok-mg-namep (*exp*) ∨ (*exp* = `'leave`))
∧ (*exp* ∈ *cond-list*)))

THEOREM: ok-condition-litatom
ok-condition (*exp*, *cond-list*) → litatom (*exp*)

EVENT: Disable ok-condition-litatom.

```
;; Notice that I'm not allowing references to individual array elements, so
;; this doesn't change from the previous version.
```

DEFINITION:
ok-actual-params-list (*lst*, *alist*)
= **if** *lst* ≃ **nil** **then** *lst* = **nil**
else defined-identifierp (car (*lst*), *alist*)
∧ ok-actual-params-list (cdr (*lst*), *alist*) **endif**

```
;; The formal is of the form (name kind typeref);  kind is in {var-mg const-mg}.

;; There is a match between actual and formal if the actual is a literal of the type of the
;; formal or is an identifier of exactly the same type as the formal.
;; This will have to be relaxed for subtypes when I add them.
```

DEFINITION:
ok-identifier-actual (*actual*, *formal*, *alist*)
= (identifierp (*actual*)
∧ (get-m-type (*actual*, *alist*) = formal-type (*formal*)))

DEFINITION:
data-params-match (*actual*, *formal*, *alist*)
= ok-identifier-actual (*actual*, *formal*, *alist*)

DEFINITION:
data-param-lists-match (*actuals*, *formals*, *alist*)
= **if** (*actuals* ≃ **nil**) ∨ (*formals* ≃ **nil**)
then (*formals* = **nil**) ∧ (*actuals* = **nil**)
else data-params-match (car (*actuals*), car (*formals*), *alist*)
∧ data-param-lists-match (cdr (*actuals*),
cdr (*formals*),
alist) **endif**

THEOREM: data-param-lists-match-in-length
data-param-lists-match (*actuals*, *formals*, *alist*)
→ (length (*formals*) = length (*actuals*))

EVENT: Disable data-param-lists-match-in-length.

DEFINITION:
cond-params-match (*cond-actuals*, *conds*)
= (length (*cond-actuals*) = length (*conds*))

THEOREM: list-count-decreases
$((x = \text{car}(stmt)) \land (x \neq 0))$
$\rightarrow \; (((\text{count}(\text{cadr}(stmt)) < \text{count}(stmt)) = \mathbf{t})$
$\land \; ((\text{count}(\text{caddr}(stmt)) < \text{count}(stmt)) = \mathbf{t})$
$\land \; ((\text{count}(\text{cadddr}(stmt)) < \text{count}(stmt)) = \mathbf{t})$
$\land \; ((\text{count}(\text{caddddr}(stmt)) < \text{count}(stmt)) = \mathbf{t}))$

```
;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;
;;                                                                      ;;
;;                         PREDEFINED PROCEDURES                        ;;
;;                                                                      ;;
;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;

;; Notice that I define the semantics, etc. of each of the predefined routines
;; individually.  This allows me to loosen the restrictions on user-defined
;; procedures.  It also, allows me to dispense with much of the overhead of
;; user defined routines.  The only condition that any of the predefined's can
;; return is 'routineerror.  Consequently, I can do away with the cond translation
;; mechanism required of user-defined routines.  These simply set 'routineerror
;; themselves.   Also, I can eliminate the aliasing requirement for these since
;; I guarantee by coding the aliasing does not cause any problem.

;; This approach does seem to add an additional burden with respect to the
;; amount of things which must be proved.  However, since the predefined's
;; are not defined recursively, this is not particularly bad.

;; THE 'GENERIC' OPERATIONS
;; The operations of assignment and EQ work on any of the simple types.  For
;; each of them, we allow the variant where one of the args is a literal
;; of the appropriate type.  This is not strictly necessary except insofar
;; as it makes the subset more realistic and usable.

;; (mg-simple-variable-assignment x y)
;; x := y  where source must be a variable of the same type as the destination.
```

DEFINITION:
ok-mg-simple-variable-assignment-args (*args*, *alist*)
= (length-plistp (*args*, 2)
∧ simple-identifierp (car (*args*), *alist*)
∧ simple-identifierp (cadr (*args*), *alist*)

∧ (cadr (assoc (car (*args*), *alist*))
= cadr (assoc (cadr (*args*), *alist*))))

```
;; (mg-simple-constant-assignment x c)
;; x := c where c must be a literal of the same type as the destination
```

DEFINITION:
ok-mg-simple-constant-assignment-args (*args*, *alist*)
= (length-plistp (*args*, 2)
∧ simple-identifierp (car (*args*), *alist*)
∧ simple-typed-literalp (cadr (*args*), cadr (assoc (car (*args*), *alist*))))

```
;; (mg-simple-variable-eq b x y)
;; b := (x = y) where both x and y are variables of the same type.
```

DEFINITION:
ok-mg-simple-variable-eq-args (*args*, *alist*)
= (length-plistp (*args*, 3)
∧ boolean-identifierp (car (*args*), *alist*)
∧ simple-identifierp (cadr (*args*), *alist*)
∧ simple-identifierp (caddr (*args*), *alist*)
∧ (cadr (assoc (cadr (*args*), *alist*))
= cadr (assoc (caddr (*args*), *alist*))))

```
;; (mg-simple-constant-eq b x c)
;; b := (x = c) where x is a variable and c a literal of compatible type.
```

DEFINITION:
ok-mg-simple-constant-eq-args (*args*, *alist*)
= (length-plistp (*args*, 3)
∧ boolean-identifierp (car (*args*), *alist*)
∧ simple-identifierp (cadr (*args*), *alist*)
∧ simple-typed-literalp (caddr (*args*),
cadr (assoc (cadr (*args*), *alist*))))

```
;; THE INTEGER OPERATIONS

;; (mg-integer-le b x y)
;; b := (x le y) where both x and y are integer variables
```

DEFINITION:

ok-mg-integer-le-args (*args*, *alist*)
= (length-plistp (*args*, 3)
∧ boolean-identifierp (car (*args*), *alist*)
∧ int-identifierp (cadr (*args*), *alist*)
∧ int-identifierp (caddr (*args*), *alist*))

```
;; (mg-integer-unary-minus (args alist))
;; z := -x
```

Definition:
ok-mg-integer-unary-minus-args (*args*, *alist*)
= (length-plistp (*args*, 2)
∧ int-identifierp (car (*args*), *alist*)
∧ int-identifierp (cadr (*args*), *alist*))

```
;; (mg-integer-add (x y z))
;; x := y + z  >> Notice that this is a change from the previous version.
```

Definition:
ok-mg-integer-add-args (*args*, *alist*)
= (length-plistp (*args*, 3)
∧ int-identifierp (car (*args*), *alist*)
∧ int-identifierp (cadr (*args*), *alist*)
∧ int-identifierp (caddr (*args*), *alist*))

```
;; (mg-integer-subtract (x y z))
;; x := y - z
```

Definition:
ok-mg-integer-subtract-args (*args*, *alist*)
= (length-plistp (*args*, 3)
∧ int-identifierp (car (*args*), *alist*)
∧ int-identifierp (cadr (*args*), *alist*)
∧ int-identifierp (caddr (*args*), *alist*))

```
;; BOOLEAN OPERATIONS

;; (mg-boolean-or (b c d))
;; b := c or d   -- both disjuncts must be boolean identifiers
```

Definition:

ok-mg-boolean-or-args (*args*, *alist*)
= (length-plistp (*args*, 3)
∧ boolean-identifierp (car (*args*), *alist*)
∧ boolean-identifierp (cadr (*args*), *alist*)
∧ boolean-identifierp (caddr (*args*), *alist*))

```
;; (mg-boolean-and (b c d))
;; b := c and d
```

Definition:
ok-mg-boolean-and-args (*args*, *alist*)
= (length-plistp (*args*, 3)
∧ boolean-identifierp (car (*args*), *alist*)
∧ boolean-identifierp (cadr (*args*), *alist*)
∧ boolean-identifierp (caddr (*args*), *alist*))

```
;; (mg-boolean-not (b c))
;; b := not c
```

Definition:
ok-mg-boolean-not-args (*args*, *alist*)
= (length-plistp (*args*, 2)
∧ boolean-identifierp (car (*args*), *alist*)
∧ boolean-identifierp (cadr (*args*), *alist*))

```
;; ARRAY OPERATIONS

;; (mg-index-array (z A i size))  >> Notice the change in the order of the args
;; z := A[i]
;; It is necessary for that last to be passed to do bounds-checking.  That information
;; is not available to the translator otherwise.  It is a special argument which is a
;; numberp rather than an MG literal.  It is expected that the preprocessor
;; will actually supply this argument, not
;; the mg programmer.  Thus the prefix form might have args (z A i 24) where 24 is the
;; size of A.
```

Definition:
ok-mg-index-array-args (*args*, *alist*)
= (length-plistp (*args*, 4)
∧ array-identifierp (cadr (*args*), *alist*)
∧ int-identifierp (caddr (*args*), *alist*)
∧ simple-typed-identifierp (car (*args*),

array-elemtype (cadr (assoc (cadr (*args*), *alist*))), *alist*)
∧ (cadddr (*args*) = array-length (cadr (assoc (cadr (*args*), *alist*))))
∧ (cadddr (*args*) < MAXINT))

```
;; (mg-array-element-assignment A i value size)
;; A[i] := value
;; Here i must be an integer variable and value a variable of
;; the appropriate element-type.
```

DEFINITION:
ok-mg-array-element-assignment-args (*args*, *alist*)
= (length-plistp (*args*, 4)
∧ array-identifierp (car (*args*), *alist*)
∧ int-identifierp (cadr (*args*), *alist*)
∧ (cadddr (*args*) = array-length (cadr (assoc (car (*args*), *alist*))))
∧ (cadddr (*args*) < MAXINT)
∧ simple-typed-identifierp (caddr (*args*),
array-elemtype (cadr (assoc (car (*args*), *alist*))), *alist*))

DEFINITION:
ok-predefined-proc-args (*name*, *args*, *alist*)
= **case on** *name*:
case = *mg-simple-variable-assignment*
then ok-mg-simple-variable-assignment-args (*args*, *alist*)
case = *mg-simple-constant-assignment*
then ok-mg-simple-constant-assignment-args (*args*, *alist*)
case = *mg-simple-variable-eq*
then ok-mg-simple-variable-eq-args (*args*, *alist*)
case = *mg-simple-constant-eq*
then ok-mg-simple-constant-eq-args (*args*, *alist*)
case = *mg-integer-le*
then ok-mg-integer-le-args (*args*, *alist*)
case = *mg-integer-unary-minus*
then ok-mg-integer-unary-minus-args (*args*, *alist*)
case = *mg-integer-add*
then ok-mg-integer-add-args (*args*, *alist*)
case = *mg-integer-subtract*
then ok-mg-integer-subtract-args (*args*, *alist*)
case = *mg-boolean-or*
then ok-mg-boolean-or-args (*args*, *alist*)

case = *mg-boolean-and*
then ok-mg-boolean-and-args (*args*, *alist*)
case = *mg-boolean-not*
then ok-mg-boolean-not-args (*args*, *alist*)
case = *mg-index-array*
then ok-mg-index-array-args (*args*, *alist*)
case = *mg-array-element-assignment*
then ok-mg-array-element-assignment-args (*args*, *alist*)
otherwise f endcase

EVENT: Disable ok-predefined-proc-args.

```
;; A predefined proc-call is of the form (predefined-proc-call-mg name actuals)
;; where name is one of the legal predefineds and the actuals are legitimate arguments
;; for that predefined procedure according to the definitions above.
```

DEFINITION:
ok-predefined-proc-call (*stmt*, *alist*)
= (length-plistp (*stmt*, 3)
∧ predefined-procp (call-name (*stmt*))
∧ ok-predefined-proc-args (call-name (*stmt*),
call-actuals (*stmt*),
alist))

EVENT: Disable ok-predefined-proc-call.

DEFINITION:
ok-proc-call (*stmt*, *r-cond-list*, *alist*, *proc-list*)
= (length-plistp (*stmt*, 4)
∧ identifierp (call-name (*stmt*))
∧ user-defined-procp (call-name (*stmt*), *proc-list*)
∧ ok-actual-params-list (call-actuals (*stmt*), *alist*)
∧ no-duplicates (call-actuals (*stmt*))
∧ data-param-lists-match (call-actuals (*stmt*),
def-formals (fetch-called-def (*stmt*,
proc-list)),
alist)
∧ cond-identifier-plistp (call-conds (*stmt*), *r-cond-list*)
∧ cond-params-match (call-conds (*stmt*),
def-conds (fetch-called-def (*stmt*, *proc-list*))))

EVENT: Disable ok-proc-call.

DEFINITION:
ok-mg-statement (*stmt*, *r-cond-list*, *alist*, *proc-list*)
= **case on** car (*stmt*):
case = *no-op-mg*
then cdr (*stmt*) = **nil**
case = *signal-mg*
then length-plistp (*stmt*, 2)
∧ ok-condition (signalled-condition (*stmt*), *r-cond-list*)
case = *prog2-mg*
then length-plistp (*stmt*, 3)
∧ ok-mg-statement (prog2-left-branch (*stmt*),
r-cond-list,
alist,
proc-list)
∧ ok-mg-statement (prog2-right-branch (*stmt*),
r-cond-list,
alist,
proc-list)
case = *loop-mg*
then length-plistp (*stmt*, 2)
∧ ok-mg-statement (loop-body (*stmt*),
cons (`'leave`, *r-cond-list*),
alist,
proc-list)
case = *if-mg*
then length-plistp (*stmt*, 4)
∧ boolean-identifierp (if-condition (*stmt*), *alist*)
∧ ok-mg-statement (if-true-branch (*stmt*),
r-cond-list,
alist,
proc-list)
∧ ok-mg-statement (if-false-branch (*stmt*),
r-cond-list,
alist,
proc-list)
case = *begin-mg*
then length-plistp (*stmt*, 4)
∧ ok-mg-statement (begin-body (*stmt*),
append (when-labels (*stmt*), *r-cond-list*),
alist,
proc-list)
∧ nonempty-cond-identifier-plistp (when-labels (*stmt*),
r-cond-list)
∧ ok-mg-statement (when-handler (*stmt*),

r-cond-list,
alist,
proc-list)
case = *proc-call-mg*
then ok-proc-call (*stmt*, *r-cond-list*, *alist*, *proc-list*)
case = *predefined-proc-call-mg*
then ok-predefined-proc-call (*stmt*, *alist*)
otherwise f endcase

EVENT: Disable signalled-condition.

EVENT: Disable prog2-left-branch.

EVENT: Disable prog2-right-branch.

EVENT: Disable loop-body.

EVENT: Disable if-condition.

EVENT: Disable if-true-branch.

EVENT: Disable if-false-branch.

EVENT: Disable begin-body.

EVENT: Disable when-handler.

EVENT: Disable when-labels.

EVENT: Disable call-name.

EVENT: Disable call-actuals.

EVENT: Disable call-conds.

THEOREM: ok-signal-expansion
ok-mg-statement (cons (`'signal-mg`, *args*), *r-cond-list*, *alist*, *proc-list*)

```
=    (length-plistp (cons ('signal-mg, args), 2)
      ∧   ok-condition (signalled-condition (cons ('signal-mg, args)),
                                            r-cond-list))
```

THEOREM: ok-prog2-statement

```
((car (stmt) = 'prog2-mg)
 ∧   ok-mg-statement (stmt, r-cond-list, alist, proc-list))
→    (ok-mg-statement (prog2-left-branch (stmt), r-cond-list, alist, proc-list)
      ∧   ok-mg-statement (prog2-right-branch (stmt),
                                            r-cond-list,
                                            alist,
                                            proc-list))
```

THEOREM: ok-loop-statement

```
((car (stmt) = 'loop-mg)
 ∧   ok-mg-statement (stmt, r-cond-list, alist, proc-list))
→    ok-mg-statement (loop-body (stmt),
                                 cons ('leave, r-cond-list),
                                 alist,
                                 proc-list)
```

THEOREM: ok-if-statement

```
((car (stmt) = 'if-mg)
 ∧   ok-mg-statement (stmt, r-cond-list, alist, proc-list))
→    (ok-mg-statement (if-true-branch (stmt), r-cond-list, alist, proc-list)
      ∧   ok-mg-statement (if-false-branch (stmt),
                                            r-cond-list,
                                            alist,
                                            proc-list))
```

THEOREM: ok-begin-expansion

```
ok-mg-statement (cons ('begin-mg, args), r-cond-list, alist, proc-list)
=    (length-plistp (cons ('begin-mg, args), 4)
      ∧   ok-mg-statement (begin-body (cons ('begin-mg, args)),
                                      append (when-labels (cons ('begin-mg, args)),
                                                          r-cond-list),
                                      alist,
                                      proc-list)
      ∧   nonempty-cond-identifier-plistp (when-labels (cons ('begin-mg,
                                                              args)),
                                                        r-cond-list)
      ∧   ok-mg-statement (when-handler (cons ('begin-mg, args)),
                                      r-cond-list,
                                      alist,
                                      proc-list))
```

THEOREM: ok-begin-statement
((car (*stmt*) = 'begin-mg)
 ∧ ok-mg-statement (*stmt*, *r-cond-list*, *alist*, *proc-list*))
→ (ok-mg-statement (begin-body (*stmt*),
append (when-labels (*stmt*), *r-cond-list*),
alist,
proc-list)
∧ ok-mg-statement (when-handler (*stmt*),
r-cond-list,
alist,
proc-list))

THEOREM: ok-proc-call-expansion
ok-mg-statement (cons ('proc-call-mg, *args*), *r-cond-list*, *alist*, *proc-list*)
= ok-proc-call (cons ('proc-call-mg, *args*), *r-cond-list*, *alist*, *proc-list*)

EVENT: Disable ok-mg-statement.

THEOREM: signalled-condition-not-normal
(('signal-mg = car (*stmt*))
 ∧ ok-mg-statement (*stmt*, *r-cond-list*, *alist*, *proc-list*))
→ (signalled-condition (*stmt*) ≠ 'normal)

EVENT: Disable signalled-condition-not-normal.

```
;; member of the data formal list is of the form
;;   (name typeref)
```

DEFINITION:
ok-mg-formal-data-param (*exp*)
= (length-plistp (*exp*, 2)
∧ ok-mg-namep (car (*exp*))
∧ mg-type-refp (formal-type (*exp*)))

EVENT: Disable ok-mg-formal-data-param.

DEFINITION:
ok-mg-formal-data-params-plistp (*lst*)
= **if** *lst* ≃ **nil** **then** *lst* = **nil**
else ok-mg-formal-data-param (car (*lst*))
∧ ok-mg-formal-data-params-plistp (cdr (*lst*)) **endif**

```
;; A local decl is of the form (name typeref initial-value).
;; Notice that this obviates the need to compute initial values.
```

DEFINITION:
ok-mg-local-data-decl (*exp*)
= (length-plistp (*exp*, 3)
  ∧ ok-mg-namep (car (*exp*))
  ∧ mg-type-refp (formal-type (*exp*))
  ∧ ok-mg-valuep (formal-initial-value (*exp*), formal-type (*exp*)))

DEFINITION:
ok-mg-local-data-plistp (*lst*)
= **if** *lst* ≃ **nil** **then** *lst* = **nil**
  else ok-mg-local-data-decl (car (*lst*))
    ∧ ok-mg-local-data-plistp (cdr (*lst*)) **endif**

```
;; The legal conditions which can signalled are those which appear in the
;; the formal or in the local conds.
```

DEFINITION:
make-cond-list (*def*) = append (def-conds (*def*), def-cond-locals (*def*))

```
;; This takes a list of formal data params or local var decls and
;; makes a name-alist.

;; >>> This is a useless definition.  It does nothing
```

DEFINITION:
make-alist-from-formals (*lst*)
= **if** *lst* ≃ **nil** **then** **nil**
  else cons (list (name (car (*lst*)), formal-type (car (*lst*))),
    make-alist-from-formals (cdr (*lst*))) **endif**

```
;; This takes the formals and locals and makes a name-alist.

;; >>> Why not just concatenate them.
```

DEFINITION:
make-name-alist (*def*)
= append (make-alist-from-formals (def-formals (*def*)),
    make-alist-from-formals (def-locals (*def*)))

```
;; Given a list of formal-data-params or local-data-decls this lists
;; off the names.  This is necessary to check that all local names are
;; unique.
```

DEFINITION:
collect-local-names (*def*)
= append (listcars (def-formals (*def*)), listcars (def-locals (*def*)))

DEFINITION:
ok-mg-def (*def*, *proc-list*)
= (length-plistp (*def*, 6)
∧ ok-mg-namep (def-name (*def*))
∧ ok-mg-formal-data-params-plistp (def-formals (*def*))
∧ identifier-plistp (def-conds (*def*))
∧ ok-mg-local-data-plistp (def-locals (*def*))
∧ identifier-plistp (def-cond-locals (*def*))
∧ no-duplicates (collect-local-names (*def*))
∧ ((length (def-conds (*def*)) + length (def-cond-locals (*def*)))
< (((exp (2, MG-WORD-SIZE) − 1) − 1) − 1))
∧ ok-mg-statement (def-body (*def*),
make-cond-list (*def*),
make-name-alist (*def*),
proc-list))

THEOREM: make-cond-list-legal-length
ok-mg-def (*def*, *proc-list*)
→ ((length (make-cond-list (*def*))
< (((exp (2, MG-WORD-SIZE) − 1) − 1) − 1))
= **t**)

```
;; (name formals locals conds local-conds body)
```

DEFINITION:
ok-mg-def-plistp1 (*lst1*, *lst2*)
= **if** *lst1* ≃ **nil** **then** *lst1* = **nil**
else ok-mg-def (car (*lst1*), *lst2*)
∧ ok-mg-def-plistp1 (cdr (*lst1*), *lst2*) **endif**

DEFINITION:
ok-mg-def-plistp (*proc-list*) = ok-mg-def-plistp1 (*proc-list*, *proc-list*)

EVENT: Disable ok-mg-def.

THEOREM: assoc-def-ok1
(ok-mg-def-plistp1 (*proc-list1*, *proc-list2*) ∧ definedp (*name*, *proc-list1*))
→ ok-mg-def (assoc (*name*, *proc-list1*), *proc-list2*)

EVENT: Disable assoc-def-ok1.

THEOREM: called-def-ok
((`'proc-call-mg` = car (*stmt*))
∧ ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*)
∧ ok-mg-def-plistp (*proc-list*))
→ ok-mg-def (fetch-called-def (*stmt*, *proc-list*), *proc-list*)

EVENT: Disable fetch-called-def.

THEOREM: called-def-formals-ok
((`'proc-call-mg` = car (*stmt*))
∧ ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*)
∧ ok-mg-def-plistp (*proc-list*))
→ (ok-mg-formal-data-params-plistp (def-formals (fetch-called-def (*stmt*, *proc-list*)))
∧ ok-mg-local-data-plistp (def-locals (fetch-called-def (*stmt*, *proc-list*)))
∧ data-param-lists-match (call-actuals (*stmt*), def-formals (fetch-called-def (*stmt*, *proc-list*)), *name-alist*)
∧ no-duplicates (listcars (def-formals (fetch-called-def (*stmt*, *proc-list*))))
∧ plistp (make-cond-list (fetch-called-def (*stmt*, *proc-list*)))
∧ no-duplicates (append (listcars (def-formals (fetch-called-def (*stmt*, *proc-list*))), listcars (def-locals (fetch-called-def (*stmt*, *proc-list*)))))
∧ all-cars-unique (def-formals (fetch-called-def (*stmt*, *proc-list*)))
∧ all-cars-unique (def-locals (fetch-called-def (*stmt*, *proc-list*)))))

THEOREM: ok-locals-plistp
ok-mg-local-data-plistp (*x*) → plistp (*x*)

EVENT: Disable ok-locals-plistp.

THEOREM: locals-plistp
((`'proc-call-mg` = car (*stmt*))

$\wedge$ ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*)
$\wedge$ ok-mg-def-plistp (*proc-list*))
$\rightarrow$ plistp (def-locals (fetch-called-def (*stmt*, *proc-list*)))

EVENT: Disable mg-name-alist-elementp.

```
;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;
;;                                                                            ;;
;;                        MG INTERPRETER FUNCTIONS                            ;;
;;                                                                            ;;
;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;

;; The following set of functions defines the interpreter for Micro-Gypsy.
;; A state is defined to be an ordered pair consisting of a variable a-list and a
;; global condition indicator.
;; The meaning of a statement in an environment (state) is the environment
;; which results from executing that statement.  Thus, the semantics is very
;; much an operational semantics.

;; The condition component of the state is simply a litatom which is 'normal
;; in the initial state.  The variable a-list is of the form
;; ((v1 (value1 type1)) ... (vn (valuen typen))) for each of the entities known in
;; the current scope.  Notice that the recognizer alist is not required though I
;; will need to record on the var-alist whether a variable is a const or var
;; param when I add procedure calls.
```

EVENT: Add the shell *mg-state*, with recognizer function symbol *mg-statep* and 3 accessors: *cc*, with type restriction (none-of) and default value false; *mg-alist*, with type restriction (none-of) and default value false; *mg-psw*, with type restriction (none-of) and default value false.

DEFINITION:
resource-errorp (*mg-state*) = (mg-psw (*mg-state*) $\neq$ 'run)

DEFINITION:
signal-system-error (*mg-state*, *error*)
= mg-state (cc (*mg-state*), mg-alist (*mg-state*), *error*)

DEFINITION: normal (*mg-state*) = (cc (*mg-state*) = 'normal)

DEFINITION: m-value (*x*) = caddr (*x*)

DEFINITION:
get-m-value (*name*, *alist*) = m-value (assoc (*name*, *alist*))

DEFINITION:
mg-alist-elementp (x)
= (length-plistp $(x,$ 3)
$\wedge$ ok-mg-namep (car (x))
$\wedge$ mg-type-refp (m-type (x))
$\wedge$ ok-mg-valuep (m-value (x), m-type (x)))

THEOREM: new-value-mg-alist-elementp
(mg-alist-elementp $(x) \wedge$ ok-mg-valuep $(value,$ cadr $(x)))$
$\rightarrow$ mg-alist-elementp (cons (car (x), cons (cadr (x), cons $(value,$ cdddr $(x)))))$

EVENT: Disable mg-alist-elementp.

DEFINITION:
mg-alistp (lst)
= **if** $lst \simeq$ **nil then** $lst =$ **nil**
else mg-alist-elementp (car $(lst)) \wedge$ mg-alistp (cdr $(lst))$ **endif**

THEOREM: mg-alistp-cdr
(listp $(x) \wedge$ mg-alistp $(x)) \rightarrow$ mg-alistp (cdr $(x))$

THEOREM: mg-alistp-cons
mg-alistp (cons $(x,$ cons $(y,\ z))) \rightarrow$ mg-alistp (cons $(x,\ z))$

THEOREM: mg-alist-member-mg-alist-elementp
(mg-alistp $(mg\text{-}alist) \wedge (x \in mg\text{-}alist)) \rightarrow$ mg-alist-elementp (x)

THEOREM: mg-alistp-distributes
mg-alistp (append $(x,\ y)) \rightarrow$ mg-alistp (y)

THEOREM: mg-alistp-distributes2
(mg-alistp (append $(x,\ y)) \wedge$ plistp $(x)) \rightarrow$ mg-alistp (x)

EVENT: Disable mg-alistp-distributes2.

THEOREM: mg-alist-mg-name-alistp
mg-alistp $(lst) \rightarrow$ mg-name-alistp (lst)

THEOREM: mg-alistp-plistp
mg-alistp $(alist) \rightarrow$ plistp $(alist)$

EVENT: Disable mg-alistp-plistp.

THEOREM: mg-alist-elements-have-ok-values
(mg-alistp $(alist) \wedge$ definedp $(x,\ alist))$
$\rightarrow$ ok-mg-valuep (caddr (assoc $(x,\ alist)),$ cadr (assoc $(x,\ alist)))$

THEOREM: restrict-preserves-mg-alistp
mg-alistp (*alist*) → mg-alistp (restrict (*alist*, *names*))

DEFINITION:
ok-cc (*c*, *cond-list*)
= (litatom (*c*)
∧ ((*c* ∈ '(normal routineerror)) ∨ (*c* ∈ *cond-list*)))

THEOREM: mg-alistp-implies-mg-statep
mg-alistp (mg-alist (*mg-state*)) → mg-statep (*mg-state*)

DEFINITION:
ok-mg-statep (*mg-state*, *cond-list*)
= (ok-cc (cc (*mg-state*), *cond-list*) ∧ mg-alistp (mg-alist (*mg-state*)))

THEOREM: ok-mg-statep-alist-plistp
ok-mg-statep (*mg-state*, *cond-list*) → plistp (mg-alist (*mg-state*))

THEOREM: cons-preserves-ok-mg-statep
ok-mg-statep (*mg-state*, *cond-list*)
→ ok-mg-statep (*mg-state*, cons (*x*, *cond-list*))

EVENT: Disable cons-preserves-ok-mg-statep.

DEFINITION:
set-condition (*mg-state*, *condition-name*)
= mg-state (*condition-name*, mg-alist (*mg-state*), mg-psw (*mg-state*))

THEOREM: cc-set-condition
cc (set-condition (*mg-state*, *cond*)) = *cond*

THEOREM: mg-alist-set-condition
mg-alist (set-condition (*mg-state*, *cond*)) = mg-alist (*mg-state*)

THEOREM: ok-mg-statep-mg-alist-mg-alistp
ok-mg-statep (*mg-state*, *r-cond-list*) → mg-alistp (mg-alist (*mg-state*))

DEFINITION:
remove-leave (*mg-state*)
= **if** cc (*mg-state*) = 'leave **then** set-condition (*mg-state*, 'normal)
else *mg-state* **endif**

DEFINITION:
mg-expression-falsep (*exp*, *mg-state*)
= (get-m-value (*exp*, mg-alist (*mg-state*))
= '(boolean-mg false-mg))

```
;; This assumes that the formals and actuals are both given as lists of
;; bare cond names.  This will have to change when the identifier lists
;; are reinstated.
;; Given condition lists (formal1 formal2 ... formaln) and
;; (actual1 actual2 ... actualn), if cond is formali then return actuali,
;; else 'routineerror.
```

DEFINITION:
convert-condition1 (*cond*, *formals*, *actuals*)
= **if** *formals* $\simeq$ **nil** **then** `'routineerror`
elseif *cond* = car (*formals*) **then** car (*actuals*)
else convert-condition1 (*cond*, cdr (*formals*), cdr (*actuals*)) **endif**

DEFINITION:
convert-condition (*cond*, *formals*, *actuals*)
= **if** *cond* $\in$ `'(normal routineerror)` **then** *cond*
else convert-condition1 (*cond*, *formals*, *actuals*) **endif**

THEOREM: convert-condition-non-member
(*cond* $\notin$ *formals*)
$\rightarrow$ (convert-condition1 (*cond*, *formals*, *actuals*) = `'routineerror`)

DEFINITION:
set-alist-value (*name*, *val*, *alist*)
= **if** *alist* $\simeq$ **nil** **then nil**
elseif car (car (*alist*)) = *name*
then cons (cons (*name*,
cons (m-type (car (*alist*)), cons (*val*, cdddr (car (*alist*))))),
cdr (*alist*))
else cons (car (*alist*), set-alist-value (*name*, *val*, cdr (*alist*))) **endif**

THEOREM: set-alist-value-preserves-definedp
definedp (*v*, *alist*) $\rightarrow$ definedp (*v*, set-alist-value (*x*, *y*, *alist*))

EVENT: Disable set-alist-value-preserves-definedp.

THEOREM: set-alist-value-preserves-ok-actual-params-list
ok-actual-params-list (*actuals*, *alist*)
$\rightarrow$ ok-actual-params-list (*actuals*, set-alist-value (*x*, *y*, *alist*))

EVENT: Disable set-alist-value-preserves-ok-actual-params-list.

THEOREM: set-alist-value-preserves-cadr-assoc
mg-alistp (*alist*)
$\rightarrow$ (cadr (assoc (*v*, set-alist-value (*x*, *y*, *alist*))) = cadr (assoc (*v*, *alist*)))

EVENT: Disable set-alist-value-preserves-cadr-assoc.

THEOREM: set-alist-value-preserves-data-param-lists-match
(mg-alistp (*alist*) ∧ data-param-lists-match (*actuals*, *formals*, *alist*))
→ data-param-lists-match (*actuals*, *formals*, set-alist-value (*x*, *y*, *alist*))

EVENT: Disable set-alist-value-preserves-data-param-lists-match.

THEOREM: set-alist-value-preserves-listcars
listcars (set-alist-value (*x*, *y*, *alist*)) = listcars (*alist*)

THEOREM: set-alist-value-preserves-all-cars-unique
all-cars-unique (*alist*) → all-cars-unique (set-alist-value (*x*, *y*, *alist*))

EVENT: Disable set-alist-value-preserves-all-cars-unique.

THEOREM: set-alist-value-preserves-signatures-match
plistp (*alist*) → signatures-match (*alist*, set-alist-value (*x*, *y*, *alist*))

EVENT: Disable set-alist-value-preserves-signatures-match.

DEFINITION:
copy-out-params (*formals*, *actuals*, *new-var-alist*, *old-var-alist*)
= **if** *formals* ≃ **nil** **then** *old-var-alist*
else copy-out-params (cdr (*formals*),
cdr (*actuals*),
new-var-alist,
set-alist-value (car (*actuals*),
caddr (assoc (caar (*formals*),
new-var-alist)),
old-var-alist)) **endif**

DEFINITION:
map-call-effects (*new-state*, *def*, *stmt*, *old-state*)
= mg-state (convert-condition (cc (*new-state*),
def-conds (*def*),
call-conds (*stmt*)),
copy-out-params (def-formals (*def*),
call-actuals (*stmt*),
mg-alist (*new-state*),
mg-alist (*old-state*)),
mg-psw (*new-state*))

EVENT: Disable map-call-effects.

THEOREM: map-call-effects-preserves-normal
normal (*new-state*)
→ (cc (map-call-effects (*new-state*, *def*, *stmt*, *old-state*)) = 'normal)

THEOREM: map-call-effects-preserves-routineerror
(cc (*new-state*) = 'routineerror)
→ (cc (map-call-effects (*new-state*, *def*, *stmt*, *old-state*))
= 'routineerror)

```
;; This creates the part of the var alist for the call corresponding to the
;; formals.  The formal has form (name kind type) and the actual is either
;; a literal or an identifierp.
```

DEFINITION:
make-call-param-alist (*formals*, *actuals*, *mg-alist*)
= **if** *formals* ≃ **nil then nil**
else cons (list (caar (*formals*),
cadar (*formals*),
caddr (assoc (car (*actuals*), *mg-alist*))),
make-call-param-alist (cdr (*formals*),
cdr (*actuals*),
mg-alist)) **endif**

THEOREM: make-call-param-alist-plistp
plistp (make-call-param-alist (*formals*, *actuals*, *alist*))

THEOREM: make-call-param-alist-preserves-listcars
listcars (make-call-param-alist (*formals*, *actuals*, *mg-alist*))
= listcars (*formals*)

DEFINITION:
make-call-var-alist (*mg-alist*, *stmt*, *def*)
= append (make-call-param-alist (def-formals (*def*),
call-actuals (*stmt*),
mg-alist),
def-locals (*def*))

```
;; This doesn't really need the hypothesis.
```

THEOREM: plistp-make-call-var-alist
plistp (def-locals (*def*)) → plistp (make-call-var-alist (*state*, *stmt*, *def*))

DEFINITION:
make-call-environment (*mg-state*, *stmt*, *def*)
= mg-state ('normal,
make-call-var-alist (mg-alist (*mg-state*), *stmt*, *def*),
mg-psw (*mg-state*))

EVENT: Disable make-call-environment.

THEOREM: make-call-environment-decomposition
(cc (make-call-environment (*mg-state*, *stmt*, *def*)) = 'normal)
∧ (mg-alist (make-call-environment (*mg-state*, *stmt*, *def*))
= make-call-var-alist (mg-alist (*mg-state*), *stmt*, *def*))
∧ (mg-psw (make-call-environment (*mg-state*, *stmt*, *def*))
= mg-psw (*mg-state*))

```
;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;
;;                                                                              ;;
;;            SEMANTICS FOR THE PREDEFINED PROCEDURES                           ;;
;;                                                                              ;;
;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;

;; These are structurally identical to the Piton version merely
;; altered to change 't and 'f to 'true-mg and 'false-mg.
```

DEFINITION:
mg-bool (*x*)
= tag ('boolean-mg,
if *x* **then** 'true-mg
else 'false-mg **endif**)

DEFINITION:
mg-or-bool (*x*, *y*)
= **if** *x* = 'false-mg **then** *y*
else 'true-mg **endif**

DEFINITION:
mg-and-bool (*x*, *y*)
= **if** *x* = 'false-mg **then** 'false-mg
else *y* **endif**

DEFINITION:
mg-not-bool (*x*)
= **if** *x* = 'false-mg **then** 'true-mg
else 'false-mg **endif**

DEFINITION:
fetch-array-element (*a*, *i*, *alist*) = get (*i*, caddr (assoc (*a*, *alist*)))

```
;; This returns the array with the substitution, not the resulting
;; alist.
```

DEFINITION:
put-array-element (*a*, *i*, *val*, *alist*) = put (*val*, *i*, caddr (assoc (*a*, *alist*)))

```
;; x := y  -- y is a variable
```

DEFINITION:
mg-meaning-mg-simple-variable-assignment (*stmt*, *mg-state*)
= mg-state (`'normal`,
set-alist-value (car (call-actuals (*stmt*)),
get-m-value (cadr (call-actuals (*stmt*)),
mg-alist (*mg-state*)),
mg-alist (*mg-state*)),
mg-psw (*mg-state*))

```
;; x := c  -- c is a constant
```

DEFINITION:
mg-meaning-mg-simple-constant-assignment (*stmt*, *mg-state*)
= mg-state (`'normal`,
set-alist-value (car (call-actuals (*stmt*)),
cadr (call-actuals (*stmt*)),
mg-alist (*mg-state*)),
mg-psw (*mg-state*))

```
;; b := (x = y)
```

DEFINITION:
mg-meaning-mg-simple-variable-eq (*stmt*, *mg-state*)
= mg-state (`'normal`,
set-alist-value (car (call-actuals (*stmt*)),
mg-bool (untag (get-m-value (cadr (call-actuals (*stmt*)),
mg-alist (*mg-state*)))
= untag (get-m-value (caddr (call-actuals (*stmt*)),
mg-alist (*mg-state*)))),
mg-alist (*mg-state*)),
mg-psw (*mg-state*))

```
;; b := (x = c)
```

DEFINITION:
mg-meaning-mg-simple-constant-eq (*stmt*, *mg-state*)
= mg-state ('normal,
set-alist-value (car (call-actuals (*stmt*)),
mg-bool (untag (get-m-value (cadr (call-actuals (*stmt*)),
mg-alist (*mg-state*)))
= untag (caddr (call-actuals (*stmt*)))),
mg-alist (*mg-state*)),
mg-psw (*mg-state*))

```
;; b := (x le y)  -- Here x and y are integer variables
```

DEFINITION:
mg-meaning-mg-integer-le (*stmt*, *mg-state*)
= mg-state ('normal,
set-alist-value (car (call-actuals (*stmt*)),
mg-bool (ileq (untag (get-m-value (cadr (call-actuals (*stmt*)),
mg-alist (*mg-state*))),
untag (get-m-value (caddr (call-actuals (*stmt*)),
mg-alist (*mg-state*))))),
mg-alist (*mg-state*)),
mg-psw (*mg-state*))

```
;; x := -y
```

DEFINITION:
mg-meaning-mg-integer-unary-minus (*stmt*, *mg-state*)
= **let** *value* **be** inegate (untag (get-m-value (cadr (call-actuals (*stmt*)),
mg-alist (*mg-state*))))
in
if small-integerp (*value*, MG-WORD-SIZE)
then mg-state ('normal,
set-alist-value (car (call-actuals (*stmt*)),
tag ('int-mg, *value*),
mg-alist (*mg-state*)),
mg-psw (*mg-state*))
else set-condition (*mg-state*, 'routineerror) **endif endlet**

```
;; x := y + z
```

```
DEFINITION:
mg-meaning-mg-integer-add (stmt, mg-state)
=  let sum  be  iplus (untag (get-m-value (cadr (call-actuals (stmt)),
                                           mg-alist (mg-state))),
                       untag (get-m-value (caddr (call-actuals (stmt)),
                                           mg-alist (mg-state))))
   in
   if small-integerp (sum, MG-WORD-SIZE)
   then mg-state ('normal,
                  set-alist-value (car (call-actuals (stmt)),
                                   tag ('int-mg, sum),
                                   mg-alist (mg-state)),
                  mg-psw (mg-state))
   else set-condition (mg-state, 'routineerror) endif endlet

;; x := y - z


DEFINITION:
mg-meaning-mg-integer-subtract (stmt, mg-state)
=  let diff  be  idifference (untag (get-m-value (cadr (call-actuals (stmt)),
                                                  mg-alist (mg-state))),
                              untag (get-m-value (caddr (call-actuals (stmt)),
                                                  mg-alist (mg-state))))
   in
   if small-integerp (diff, MG-WORD-SIZE)
   then mg-state ('normal,
                  set-alist-value (car (call-actuals (stmt)),
                                   tag ('int-mg, diff),
                                   mg-alist (mg-state)),
                  mg-psw (mg-state))
   else set-condition (mg-state, 'routineerror) endif endlet

;; b := b1 or b2


DEFINITION:
mg-meaning-mg-boolean-or (stmt, mg-state)
=  mg-state ('normal,
             set-alist-value (car (call-actuals (stmt)),
                              tag ('boolean-mg,
                                   mg-or-bool (untag (get-m-value (cadr (call-actuals (stmt)),
                                                                   mg-alist (mg-state))),
                                               untag (get-m-value (caddr (call-actuals (stmt)),
                                                                   mg-alist (mg-state)))))),
```

```
                    mg-alist (mg-state)),
          mg-psw (mg-state))

;; b := b1 and b2
```

DEFINITION:
mg-meaning-mg-boolean-and (*stmt*, *mg-state*)
= mg-state ('normal,
          set-alist-value (car (call-actuals (*stmt*)),
                    tag ('boolean-mg,
                         mg-and-bool (untag (get-m-value (cadr (call-actuals (*stmt*)),
                                                          mg-alist (*mg-state*))),
                                      untag (get-m-value (caddr (call-actuals (*stmt*)),
                                                          mg-alist (*mg-state*))))),
                    mg-alist (*mg-state*)),
          mg-psw (*mg-state*))

```
;; b := not b1
```

DEFINITION:
mg-meaning-mg-boolean-not (*stmt*, *mg-state*)
= mg-state ('normal,
          set-alist-value (car (call-actuals (*stmt*)),
                    tag ('boolean-mg,
                         mg-not-bool (untag (get-m-value (cadr (call-actuals (*stmt*)),
                                                          mg-alist (*mg-state*)))),
                    mg-alist (*mg-state*)),
          mg-psw (*mg-state*))

```
;; THE ARRAY OPERATIONS

;; z := A[i] for A of size
;; The call is (predefined-proc-call-mg mg-index-array (z A i size))
```

DEFINITION:
mg-meaning-mg-index-array (*stmt*, *mg-state*)
= **let** *index* **be** untag (get-m-value (caddr (call-actuals (*stmt*)),
                                        mg-alist (*mg-state*)))
  in
  if (*index* ∈ **N**)
    ∧ (*index* < array-length (get-m-type (cadr (call-actuals (*stmt*)),
                                           mg-alist (*mg-state*))))

then mg-state (`'normal`,
set-alist-value (car (call-actuals (*stmt*)),
fetch-array-element (cadr (call-actuals (*stmt*)),
index,
mg-alist (*mg-state*)),
mg-alist (*mg-state*)),
mg-psw (*mg-state*))
else set-condition (*mg-state*, `'routineerror`) **endif endlet**

```
;; A[i] := v   -- The actual argument list is (A i v size) where size is the
;;      array-length of the type of A.  Here i and v are both variables.
```

DEFINITION:
mg-meaning-mg-array-element-assignment (*stmt*, *mg-state*)
= **let** *index* **be** untag (get-m-value (cadr (call-actuals (*stmt*)),
mg-alist (*mg-state*))),
val **be** get-m-value (caddr (call-actuals (*stmt*)),
mg-alist (*mg-state*))
in
if (*index* $\in$ **N**)
$\wedge$ (*index* $<$ array-length (get-m-type (car (call-actuals (*stmt*)),
mg-alist (*mg-state*))))
then mg-state (`'normal`,
set-alist-value (car (call-actuals (*stmt*)),
put-array-element (car (call-actuals (*stmt*)),
index,
val,
mg-alist (*mg-state*)),
mg-alist (*mg-state*)),
mg-psw (*mg-state*))
else set-condition (*mg-state*, `'routineerror`) **endif endlet**

DEFINITION:
mg-meaning-predefined-proc-call (*stmt*, *mg-state*)
= **case on** call-name (*stmt*):
case = *mg-simple-variable-assignment*
then mg-meaning-mg-simple-variable-assignment (*stmt*, *mg-state*)
case = *mg-simple-constant-assignment*
then mg-meaning-mg-simple-constant-assignment (*stmt*, *mg-state*)
case = *mg-simple-variable-eq*
then mg-meaning-mg-simple-variable-eq (*stmt*, *mg-state*)
case = *mg-simple-constant-eq*
then mg-meaning-mg-simple-constant-eq (*stmt*, *mg-state*)
case = *mg-integer-le*

```
    then mg-meaning-mg-integer-le (stmt, mg-state)
  case = mg-integer-unary-minus
    then mg-meaning-mg-integer-unary-minus (stmt, mg-state)
  case = mg-integer-add
    then mg-meaning-mg-integer-add (stmt, mg-state)
  case = mg-integer-subtract
    then mg-meaning-mg-integer-subtract (stmt, mg-state)
  case = mg-boolean-or
    then mg-meaning-mg-boolean-or (stmt, mg-state)
  case = mg-boolean-and
    then mg-meaning-mg-boolean-and (stmt, mg-state)
  case = mg-boolean-not
    then mg-meaning-mg-boolean-not (stmt, mg-state)
  case = mg-index-array
    then mg-meaning-mg-index-array (stmt, mg-state)
  case = mg-array-element-assignment
    then mg-meaning-mg-array-element-assignment (stmt, mg-state)
  otherwise mg-state endcase
```

EVENT: Disable mg-meaning-predefined-proc-call.

```
DEFINITION:
mg-meaning (stmt, proc-list, mg-state, n)
=  if n ≃ 0 then signal-system-error (mg-state, 'timed-out)
   elseif ¬ normal (mg-state) then mg-state
   else case on car (stmt):
        case = no-op-mg
        then mg-state
        case = signal-mg
          then set-condition (mg-state, signalled-condition (stmt))
        case = prog2-mg
          then mg-meaning (prog2-right-branch (stmt),
                           proc-list,
                           mg-meaning (prog2-left-branch (stmt),
                                       proc-list,
                                       mg-state,
                                       n − 1),
                           n − 1)
        case = loop-mg
          then remove-leave (mg-meaning (stmt,
                                         proc-list,
                                         mg-meaning (loop-body (stmt),
                                                     proc-list,
                                                     mg-state,
```

```
                                                    n − 1),
                                          n − 1))
case = if-mg
 then if mg-expression-falsep (if-condition (stmt), mg-state)
      then mg-meaning (if-false-branch (stmt),
                       proc-list,
                       mg-state,
                       n − 1)
      else mg-meaning (if-true-branch (stmt),
                       proc-list,
                       mg-state,
                       n − 1) endif
case = begin-mg
 then if cc (mg-meaning (begin-body (stmt),
                         proc-list,
                         mg-state,
                         n − 1))
         ∈  when-labels (stmt)
      then mg-meaning (when-handler (stmt),
                       proc-list,
                       set-condition (mg-meaning (begin-body (stmt),
                                                  proc-list,
                                                  mg-state,
                                                  n − 1),
                                      'normal),
                       n − 1)
      else mg-meaning (begin-body (stmt),
                       proc-list,
                       mg-state,
                       n − 1) endif
case = proc-call-mg
 then map-call-effects (mg-meaning (def-body (fetch-called-def (stmt,
                                                                proc-list)),
                                    proc-list,
                                    make-call-environment (mg-state,
                                                           stmt,
                                                           fetch-called-def (stmt,
                                                                             proc-list)),
                                    n − 1),
                        fetch-called-def (stmt, proc-list),
                        stmt,
                        mg-state)
case = predefined-proc-call-mg
 then mg-meaning-predefined-proc-call (stmt, mg-state)
```

otherwise *mg-state* **endcase endif**

```
;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;
;;                                                                            ;;
;;                          RESOURCE ERRORS                                   ;;
;;                                                                            ;;
;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;

;; This is the version of mg-meaning with resource errors.  It should be a
;; theorem that in the absence of resource-error, it behaves exactly as
;; mg-meaning.

;; The resource descriptor is a pair <temp-stk-size, ctrl-stk-size> where
;; the two components are numberps which characterize the number of free
;; slots on the top of the stacks.  We cause a resource error if there is
;; not enough space to continue without stack overflow.  If the resource
;; requirements exceed the available space, the cc is set to 'resource-error.

;; This computes the amount of space required for the storage of the locals.  It is the
;; number of simple variables plus the sum of the lengths of the arrays.
```

DEFINITION:
data-length (*locals*)
= **if** *locals* $\simeq$ **nil then** 0
elseif simple-mg-type-refp (cadr (car (*locals*)))
then 1 + data-length (cdr (*locals*))
else array-length (cadr (car (*locals*)))
\+ data-length (cdr (*locals*)) **endif**

THEOREM: data-length-not-zerop
(ok-mg-local-data-plistp (*locals*) $\wedge$ listp (*locals*))
$\rightarrow$ (data-length (*locals*) $\not\simeq$ 0)

EVENT: Disable data-length-not-zerop.

THEOREM: data-length-not-zerop2
(ok-mg-local-data-plistp (*locals*) $\wedge$ listp (*locals*))
$\rightarrow$ ((data-length (*locals*) $\in$ **N**) $\wedge$ (data-length (*locals*) $\neq$ 0))

DEFINITION:
predefined-proc-call-temp-stk-requirement (*name*)
= **case on** *name*:
case = *mg-simple-variable-assignment*

then 2
case = *mg-simple-constant-assignment*
then 2
case = *mg-simple-variable-eq*
then 3
case = *mg-simple-constant-eq*
then 3
case = *mg-integer-le*
then 3
case = *mg-integer-unary-minus*
then 2
case = *mg-integer-add*
then 3
case = *mg-integer-subtract*
then 3
case = *mg-boolean-or*
then 3
case = *mg-boolean-and*
then 3
case = *mg-boolean-not*
then 2
case = *mg-index-array*
then 4
case = *mg-array-element-assignment*
then 4
otherwise 0 **endcase**

```
;; The number associated with each predefined procedure
;; represents the number of formals plus locals in the
;; Piton implementation.  This is required because the
;; ctrl-stk requirements for the call p-frame is
;; (plus 2
;;       (length (formal-vars def))
;;       (length (temp-var-dcls def)))
```

DEFINITION:
predefined-proc-call-bindings-count (*name*)
= **case on** *name*:
case = *mg-simple-variable-assignment*
then 2
case = *mg-simple-constant-assignment*
then 2
case = *mg-simple-variable-eq*

then 3
case = *mg-simple-constant-eq*
then 3
case = *mg-integer-le*
then 3
case = *mg-integer-unary-minus*
then 4
case = *mg-integer-add*
then 4
case = *mg-integer-subtract*
then 4
case = *mg-boolean-or*
then 3
case = *mg-boolean-and*
then 3
case = *mg-boolean-not*
then 3
case = *mg-index-array*
then 5
case = *mg-array-element-assignment*
then 5
otherwise 0 **endcase**

DEFINITION:
predefined-proc-call-p-frame-size (*name*)
= (1 + (1 + predefined-proc-call-bindings-count (*name*)))

```
;; I'm implemententing the resources-available as a pair
;; <t-size c-size> of numberps.
```

DEFINITION: t-size (x) = car (x)

DEFINITION: c-size (x) = cadr (x)

```
;; An interesting fact is that the requirements for execution of a statement are not
;; dependent on the state.  (Except in the case of begin where a when-label is
;; signalled.)  This makes the computation of the resource requirements independent of
;; mg-meaning and may allow a much cleaner treatment.
```

DEFINITION:
temp-stk-requirements (*stmt*, *proc-list*)
= **case on** car (*stmt*):
case = *no-op-mg*

then 0
case = *signal-mg*
then 1
case = *prog2-mg*
then 0
case = *loop-mg*
then 1
case = *if-mg*
then 1
case = *begin-mg*
then 1
case = *proc-call-mg*
then max (data-length (def-locals (fetch-called-def (*stmt*, *proc-list*)))
+ length (def-locals (fetch-called-def (*stmt*, *proc-list*)))
+ length (call-actuals (*stmt*)),
1)
case = *predefined-proc-call-mg*
then predefined-proc-call-temp-stk-requirement (call-name (*stmt*))
otherwise 0 **endcase**

DEFINITION:
ctrl-stk-requirements (*stmt*, *proc-list*)
= **case on** car (*stmt*):
case = *no-op-mg*
then 0
case = *signal-mg*
then 0
case = *prog2-mg*
then 0
case = *loop-mg*
then 0
case = *if-mg*
then 0
case = *begin-mg*
then 0
case = *proc-call-mg*
then 2
+ length (def-locals (fetch-called-def (*stmt*, *proc-list*)))
+ length (def-formals (fetch-called-def (*stmt*, *proc-list*)))
case = *predefined-proc-call-mg*
then predefined-proc-call-p-frame-size (call-name (*stmt*))
otherwise 0 **endcase**

EVENT: Disable temp-stk-requirements.

EVENT: Disable ctrl-stk-requirements.

```
;; Resources are inadequate if I can't perform the current operation without
;; running out of space.  This can alternatively be phrased as follows.
;; In this version, the size-pair contains the
;; <current temp-stk length, current ctrl-stk length>
```

DEFINITION:
resources-inadequatep (*stmt*, *proc-list*, *size-pair*)
= ((temp-stk-requirements (*stmt*, *proc-list*)
$\not<$ (MG-MAX-TEMP-STK-SIZE − t-size (*size-pair*)))
$\vee$ (ctrl-stk-requirements (*stmt*, *proc-list*)
$\not<$ (MG-MAX-CTRL-STK-SIZE − c-size (*size-pair*))))

EVENT: Disable resources-inadequatep.

DEFINITION:
mg-meaning-r (*stmt*, *proc-list*, *mg-state*, *n*, *sizes*)
= **if** $n \simeq$ `0` **then** signal-system-error (*mg-state*, `'timed-out`)
elseif ¬ normal (*mg-state*) **then** *mg-state*
elseif resources-inadequatep (*stmt*, *proc-list*, *sizes*)
then signal-system-error (*mg-state*, `'resource-error`)
else case on car (*stmt*):
case = *no-op-mg*
then *mg-state*
case = *signal-mg*
then set-condition (*mg-state*, signalled-condition (*stmt*))
case = *prog2-mg*
then mg-meaning-r (prog2-right-branch (*stmt*),
proc-list,
mg-meaning-r (prog2-left-branch (*stmt*),
proc-list,
mg-state,
$n - 1$,
sizes),
$n - 1$,
sizes)
case = *loop-mg*
then remove-leave (mg-meaning-r (*stmt*,
proc-list,
mg-meaning-r (loop-body (*stmt*),
proc-list,

```
                                                    mg-state,
                                                    n − 1,
                                                    sizes),
                                          n − 1,
                                          sizes))
case = if-mg
 then if mg-expression-falsep (if-condition (stmt), mg-state)
      then mg-meaning-r (if-false-branch (stmt),
                         proc-list,
                         mg-state,
                         n − 1,
                         sizes)
      else mg-meaning-r (if-true-branch (stmt),
                         proc-list,
                         mg-state,
                         n − 1,
                         sizes) endif
case = begin-mg
 then if cc (mg-meaning-r (begin-body (stmt),
                           proc-list,
                           mg-state,
                           n − 1,
                           sizes))
        ∈  when-labels (stmt)
      then mg-meaning-r (when-handler (stmt),
                         proc-list,
                         set-condition (mg-meaning-r (begin-body (stmt),
                                                     proc-list,
                                                     mg-state,
                                                     n − 1,
                                                     sizes),
                                        'normal),
                         n − 1,
                         sizes)
      else mg-meaning-r (begin-body (stmt),
                         proc-list,
                         mg-state,
                         n − 1,
                         sizes) endif
case = proc-call-mg
 then map-call-effects (mg-meaning-r (def-body (fetch-called-def (stmt,
                                                                  proc-list)),
                                      proc-list,
                                      make-call-environment (mg-state,
```

stmt,
fetch-called-def (*stmt*,
proc-list)),
$n - 1$,
list (t-size (*sizes*)
+ data-length (def-locals (fetch-called-def (*stmt*,
proc-list))),
c-size (*sizes*)
+ (2
+ length (def-locals (fetch-called-def (*stmt*,
proc-list)))
+ length (def-formals (fetch-called-def (*stmt*,
proc-list)))
fetch-called-def (*stmt*, *proc-list*),
stmt,
mg-state)
case = *predefined-proc-call-mg*
then mg-meaning-predefined-proc-call (*stmt*, *mg-state*)
otherwise *mg-state* **endcase endif**

THEOREM: map-call-effects-preserves-resource-errorp
resource-errorp (*new-state*)
→ resource-errorp (map-call-effects (*new-state*, *def*, *stmt*, *old-state*))

THEOREM: map-call-effects-preserves-mg-psw
mg-psw (map-call-effects (*new-state*, *def*, *stmt*, *old-state*))
= mg-psw (*new-state*)

THEOREM: mg-meaning-predefined-proc-call-preserves-resource-error
mg-psw (mg-meaning-predefined-proc-call (*stmt*, *mg-state*)) = mg-psw (*mg-state*)

THEOREM: resource-errors-propogate
resource-errorp (*mg-state*)
→ resource-errorp (mg-meaning-r (*stmt*, *proc-list*, *mg-state*, *n*, *sizes*))

THEOREM: resource-errors-propogate2
(mg-psw (*mg-state*) ≠ 'run)
→ (mg-psw (mg-meaning-r (*stmt*, *proc-list*, *mg-state*, *n*, *sizes*)) ≠ 'run)

```
;; This lemma shows that in the absence of resource errors, the two interpreters
;; are equivalent.
```

THEOREM: mg-meaning-equivalence
(¬ resource-errorp (mg-meaning-r (*stmt*, *proc-list*, *mg-state*, *n*, *sizes*)))
→ (mg-meaning-r (*stmt*, *proc-list*, *mg-state*, *n*, *sizes*)
= mg-meaning (*stmt*, *proc-list*, *mg-state*, *n*))

THEOREM: lessp-preserves-difference-lessp
$((y < (a - \textit{t-size1})) \wedge (\textit{t-size1} \not< \textit{t-size2}))$
$\rightarrow \quad ((y < (a - \textit{t-size2})) = \mathbf{t})$

THEOREM: map-call-effects-preserves-resource-errorp2
resource-errorp (map-call-effects (*new-state*, *stmt*, *def*, *old-state*))
$=$ resource-errorp (*new-state*)

```
;; KEY POINT, t-size and c-size are the amount used, not the amount left
```

THEOREM: more-resources-preserves-resources-adequatep2
((t-size (*sizes1*) $\not<$ t-size (*sizes2*))
$\wedge$ (c-size (*sizes1*) $\not<$ c-size (*sizes2*))
$\wedge$ ($\neg$ resources-inadequatep (*stmt*, *proc-list*, *sizes1*)))
$\rightarrow$ (resources-inadequatep (*stmt*, *proc-list*, *sizes2*) = **f**)

EVENT: Disable mg-meaning-equivalence.

DEFINITION:
meaning-induction-hint0 (*stmt*, *proc-list*, *mg-state*, *n*, *sizes1*, *sizes2*)
$=$ **if** $n \simeq$ `0` **then t**
elseif resources-inadequatep (*stmt*, *proc-list*, *sizes1*) **then t**
elseif $\neg$ normal (*mg-state*) **then t**
elseif `'no-op-mg` $=$ car (*stmt*) **then t**
elseif `'signal-mg` $=$ car (*stmt*) **then t**
elseif `'prog2-mg` $=$ car (*stmt*)
then meaning-induction-hint0 (prog2-left-branch (*stmt*),
proc-list,
mg-state,
$n - 1$,
sizes1,
sizes2)
$\wedge$ meaning-induction-hint0 (prog2-right-branch (*stmt*),
proc-list,
mg-meaning-r (prog2-left-branch (*stmt*),
proc-list,
mg-state,
$n - 1$,
sizes1),
$n - 1$,
sizes1,
sizes2)
elseif `'loop-mg` $=$ car (*stmt*)

then meaning-induction-hint0 (loop-body (*stmt*),
proc-list,
mg-state,
$n - 1$,
sizes1,
sizes2)
$\wedge$ meaning-induction-hint0 (*stmt*,
proc-list,
mg-meaning-r (loop-body (*stmt*),
proc-list,
mg-state,
$n - 1$,
sizes1),
$n - 1$,
sizes2,
sizes2)
elseif `'if-mg` = car (*stmt*)
then meaning-induction-hint0 (if-false-branch (*stmt*),
proc-list,
mg-state,
$n - 1$,
sizes1,
sizes2)
$\wedge$ meaning-induction-hint0 (if-true-branch (*stmt*),
proc-list,
mg-state,
$n - 1$,
sizes1,
sizes2)
elseif `'begin-mg` = car (*stmt*)
then meaning-induction-hint0 (begin-body (*stmt*),
proc-list,
mg-state,
$n - 1$,
sizes1,
sizes2)
$\wedge$ meaning-induction-hint0 (when-handler (*stmt*),
proc-list,
set-condition (mg-meaning-r (begin-body (*stmt*),
proc-list,
mg-state,
$n - 1$,
sizes1),
`'normal`),

$n - 1$,
sizes1,
sizes2)
elseif 'proc-call-mg = car (*stmt*)
then meaning-induction-hint0 (def-body (fetch-called-def (*stmt*, *proc-list*)),
proc-list,
make-call-environment (*mg-state*,
stmt,
fetch-called-def (*stmt*,
proc-list)),
$n - 1$,
list (t-size (*sizes1*)
+ data-length (def-locals (fetch-called-def (*stmt*,
proc-list))),
c-size (*sizes1*)
+ (2
+ length (def-locals (fetch-called-def (*stmt*,
proc-list)))
+ length (def-formals (fetch-called-def (*stmt*,
proc-list))))),
list (t-size (*sizes2*)
+ data-length (def-locals (fetch-called-def (*stmt*,
proc-list))),
c-size (*sizes2*)
+ (2
+ length (def-locals (fetch-called-def (*stmt*,
proc-list)))
+ length (def-formals (fetch-called-def (*stmt*,
proc-list))))))
elseif 'predefined-proc-call-mg = car (*stmt*) **then t**
else f endif

THEOREM: mg-meaning-equivalence3
((¬ resource-errorp (mg-meaning-r (*stmt*, *proc-list*, *mg-state*, *n*, *sizes1*)))
∧ (t-size (*sizes1*) ≮ t-size (*sizes2*))
∧ (c-size (*sizes1*) ≮ c-size (*sizes2*)))
→ (mg-meaning-r (*stmt*, *proc-list*, *mg-state*, *n*, *sizes1*)
= mg-meaning-r (*stmt*, *proc-list*, *mg-state*, *n*, *sizes2*))

EVENT: Disable mg-meaning-equivalence3.

THEOREM: mg-meaning-equivalence4
((¬ resource-errorp (mg-meaning-r (*stmt*, *proc-list*, *mg-state*, *n*, *sizes1*)))
∧ (t-size (*sizes1*) ≮ t-size (*sizes2*))

$\wedge$ (c-size (*sizes1*) $\not<$ c-size (*sizes2*)))
$\rightarrow$ (mg-meaning-r (*stmt*, *proc-list*, *mg-state*, *n*, *sizes2*)
= mg-meaning-r (*stmt*, *proc-list*, *mg-state*, *n*, *sizes1*))

EVENT: Disable mg-meaning-equivalence4.

THEOREM: more-resources-preserves-not-resource-errorp0
((*t-size1* $\not<$ *t-size2*)
$\wedge$ (*c-size1* $\not<$ *c-size2*)
$\wedge$ resource-errorp (mg-meaning-r (*stmt*,
proc-list,
mg-state,
n,
list (*t-size2*, *c-size2*))))
$\rightarrow$ resource-errorp (mg-meaning-r (*stmt*,
proc-list,
mg-state,
n,
list (*t-size1*, *c-size1*))))

THEOREM: more-resources-preserves-not-resource-errorp
((*t-size1* $\not<$ *t-size2*)
$\wedge$ (*c-size1* $\not<$ *c-size2*)
$\wedge$ ($\neg$ resource-errorp (mg-meaning-r (*stmt*,
proc-list,
mg-state,
n,
list (*t-size1*, *c-size1*)))))
$\rightarrow$ ($\neg$ resource-errorp (mg-meaning-r (*stmt*,
proc-list,
mg-state,
n,
list (*t-size2*, *c-size2*))))

THEOREM: mg-meaning-equivalence2
(($\neg$ resource-errorp (mg-meaning-r (*stmt*,
proc-list,
mg-state,
n,
list (*t-size1*, *c-size1*))))
$\wedge$ (*t-size1* $\not<$ *t-size2*)
$\wedge$ (*c-size1* $\not<$ *c-size2*))
$\rightarrow$ (mg-meaning-r (*stmt*, *proc-list*, *mg-state*, *n*, list (*t-size2*, *c-size2*))
= mg-meaning (*stmt*, *proc-list*, *mg-state*, *n*))

THEOREM: mg-meaning-mg-meaning-r-resource-error-equivalence
(¬ resource-errorp (mg-meaning-r (*stmt*, *proc-list*, *mg-state*, *n*, *sizes*)))
→ (¬ resource-errorp (mg-meaning (*stmt*, *proc-list*, *mg-state*, *n*)))

THEOREM: zerop-n-mg-meaning
(*n* ≃ 0)
→ (mg-meaning (*stmt*, *proc-list*, *mg-state*, *n*)
= signal-system-error (*mg-state*, 'timed-out))

THEOREM: not-normal-mg-meaning
((*n* ≄ 0) ∧ (¬ normal (*mg-state*)))
→ (mg-meaning (*stmt*, *proc-list*, *mg-state*, *n*) = *mg-state*)

THEOREM: proc-call-meaning-2
(car (*stmt*) = 'proc-call-mg)
→ (mg-meaning (*stmt*, *proc-list*, *mg-state*, *n*)
= **if** *n* ≃ 0 **then** signal-system-error (*mg-state*, 'timed-out)
elseif ¬ normal (*mg-state*) **then** *mg-state*
else map-call-effects (mg-meaning (def-body (fetch-called-def (*stmt*,
proc-list)),
proc-list,
make-call-environment (*mg-state*,
stmt,
fetch-called-def (*stmt*,
proc-list)),
n − 1),
fetch-called-def (*stmt*, *proc-list*),
stmt,
mg-state) **endif**)

```
;; The versions for mg-meaning-r
```

THEOREM: zerop-n-mg-meaning-r
(*n* ≃ 0)
→ (mg-meaning-r (*stmt*, *proc-list*, *mg-state*, *n*, *sizes*)
= signal-system-error (*mg-state*, 'timed-out))

THEOREM: not-normal-mg-meaning-r
((*n* ≄ 0) ∧ (¬ normal (*mg-state*)))
→ (mg-meaning-r (*stmt*, *proc-list*, *mg-state*, *n*, *sizes*) = *mg-state*)

THEOREM: resources-inadequatep-mg-meaning-r
((*n* ≄ 0)
∧ normal (*mg-state*)

$\wedge$ resources-inadequatep (*stmt*, *proc-list*, *sizes*))
$\rightarrow$ (mg-meaning-r (*stmt*, *proc-list*, *mg-state*, *n*, *sizes*)
= signal-system-error (*mg-state*, `'resource-error`))

EVENT: Disable mg-meaning-r.

THEOREM: call-cond-lists-lengths-match
((`'proc-call-mg` = car (*stmt*))
$\wedge$ ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*))
$\rightarrow$ (length (def-conds (fetch-called-def (*stmt*, *proc-list*)))
= length (call-conds (*stmt*)))

EVENT: Disable call-cond-lists-lengths-match.

THEOREM: set-alist-value-preserves-mg-alistp
(mg-alistp (*alist*) $\wedge$ ok-mg-valuep (*val*, cadr (assoc (*name*, *alist*))))
$\rightarrow$ mg-alistp (set-alist-value (*name*, *val*, *alist*))

THEOREM: mg-alistps-append
(mg-alistp (*lst1*) $\wedge$ mg-alistp (*lst2*)) $\rightarrow$ mg-alistp (append (*lst1*, *lst2*))

```
;; The recognizer has a structure called the name-alist of the following
;; form:
;;        (... (namei typei other-stuff) ...)
;; This allows the identification of the types of variables.  It should be
;; the case that the values of the variables in the meaning alist correspond
;; to their types on the name-alist.

;; This says that a variable on the name-alist has the same name and type as
;; on the variable alist.  I need to know this to guarantee that the checks for
;; legality are visible in the execution world.

;; The only time a name-alist is ever created in the recognizer is from the formal
;; and locals lists.  There is really no reason why the same structure couldn't
;; be adhered to in the interpreter.  That is, the order of the variables could
;; be maintained.  >> Where would the initial values of the variables come from
;; in that case?
;; I could have an initial alist which serves both as the name-alist and var-alist
;; for the execution of the stmt.  That is, the initial alist only has to be an
;; assignment of values to the vars which is consistent with the types.

;; Notice that this checks a very strong correspondence between the two alists.
;; Each element agrees in name and type.  The idea here is that the name and var
;; alists are really identical except that some values have been added on the
```

```
;; var alists.

;; Notice that this correspondence really defines an equivalence relation.  I
;; don't know that I'll need the full power of this.
```

THEOREM: signatures-match-preserves-get-m-type
signatures-match (*alist1*, *alist2*)
→ (cadr (assoc (*x*, *alist1*)) = cadr (assoc (*x*, *alist2*)))

EVENT: Disable signatures-match-preserves-get-m-type.

THEOREM: signatures-match-preserves-definedp
signatures-match (*alist1*, *alist2*)
→ (definedp (*x*, *alist1*) = definedp (*x*, *alist2*))

THEOREM: signatures-match-preserves-ok-actual-params-list
(signatures-match (*alist1*, *alist2*) ∧ ok-actual-params-list (*lst*, *alist1*))
→ ok-actual-params-list (*lst*, *alist2*)

THEOREM: set-alist-value-preserves-plistp
plistp (*lst*) → plistp (set-alist-value (*name*, *val*, *lst*))

THEOREM: signatures-match-preserves-boolean-identifierp
(signatures-match (*alist1*, *alist2*) ∧ boolean-identifierp (*b*, *alist1*))
→ boolean-identifierp (*b*, *alist2*)

EVENT: Disable signatures-match-preserves-boolean-identifierp.

THEOREM: signatures-match-preserves-int-identifierp
(signatures-match (*alist1*, *alist2*) ∧ int-identifierp (*x*, *alist1*))
→ int-identifierp (*x*, *alist2*)

EVENT: Disable signatures-match-preserves-int-identifierp.

THEOREM: signatures-match-preserves-character-identifierp
(signatures-match (*alist1*, *alist2*) ∧ character-identifierp (*x*, *alist1*))
→ character-identifierp (*x*, *alist2*)

EVENT: Disable signatures-match-preserves-character-identifierp.

THEOREM: signatures-match-preserves-array-identifierp
(signatures-match (*alist1*, *alist2*) ∧ array-identifierp (*x*, *alist1*))
→ array-identifierp (*x*, *alist2*)

EVENT: Disable signatures-match-preserves-array-identifierp.

THEOREM: signatures-match-preserves-simple-identifierp
(signatures-match (*alist1*, *alist2*) ∧ simple-identifierp (*x*, *alist1*))
→ simple-identifierp (*x*, *alist2*)

THEOREM: signatures-match-preserves-simple-typed-identifierp
(signatures-match (*alist1*, *alist2*) ∧ simple-typed-identifierp (*x*, *type*, *alist1*))
→ simple-typed-identifierp (*x*, *type*, *alist2*)

THEOREM: signatures-match-preserves-data-param-lists-match
(signatures-match (*alist1*, *alist2*)
∧ data-param-lists-match (*actuals*, *formals*, *alist1*))
→ data-param-lists-match (*actuals*, *formals*, *alist2*)

THEOREM: signatures-match-preserves-ok-predefined-proc-args
(signatures-match (*alist1*, *alist2*)
∧ ok-predefined-proc-args (*name*, *actuals*, *alist1*))
→ ok-predefined-proc-args (*name*, *actuals*, *alist2*)

THEOREM: signatures-match-preserves-ok-predefined-proc-call
(signatures-match (*alist1*, *alist2*) ∧ ok-predefined-proc-call (*stmt*, *alist1*))
→ ok-predefined-proc-call (*stmt*, *alist2*)

THEOREM: signatures-match-preserves-ok-mg-statement
(signatures-match (*alist1*, *alist2*)
∧ ok-mg-statement (*stmt*, *r-cond-list*, *alist1*, *proc-list*))
→ ok-mg-statement (*stmt*, *r-cond-list*, *alist2*, *proc-list*)

THEOREM: mg-meaning-predefined-proc-call-preserves-signatures-match
plistp (mg-alist (*mg-state*))
→ signatures-match (mg-alist (*mg-state*),
mg-alist (mg-meaning-predefined-proc-call (*stmt*,
mg-state)))

THEOREM: copy-out-params-preserves-signatures-match
plistp (*v*) → signatures-match (*v*, copy-out-params (*x*, *y*, *z*, *v*))

THEOREM: signatures-match-preserves-plistp
(plistp (*x*) ∧ signatures-match (*x*, *y*)) → plistp (*y*)

THEOREM: mg-meaning-preserves-signatures-match
plistp (mg-alist (*mg-state*))
→ signatures-match (mg-alist (*mg-state*),
mg-alist (mg-meaning (*stmt*, *proc-list*, *mg-state*, *n*)))

THEOREM: cadr-litatom-implies-definedp
(cadr (assoc (*x*, *alist*)) ≠ 0) → definedp (*x*, *alist*)

EVENT: Disable cadr-litatom-implies-definedp.

THEOREM: call-param-alist-mg-alistp
(mg-alistp (*mg-alist*)
 ∧ ok-mg-formal-data-params-plistp (*formals*)
 ∧ data-param-lists-match (*actuals*, *formals*, *name-alist*)
 ∧ signatures-match (*mg-alist*, *name-alist*))
→ mg-alistp (make-call-param-alist (*formals*, *actuals*, *mg-alist*))

THEOREM: call-locals-alist-mg-alistp
ok-mg-local-data-plistp (*locals-list*) → mg-alistp (*locals-list*)

THEOREM: make-call-var-alist-mg-alistp
((`'proc-call-mg` = car (*stmt*))
 ∧ ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*)
 ∧ ok-mg-def-plistp (*proc-list*)
 ∧ mg-alistp (*mg-alist*)
 ∧ signatures-match (*mg-alist*, *name-alist*))
→ mg-alistp (make-call-var-alist (*mg-alist*,
 stmt,
 fetch-called-def (*stmt*, *proc-list*)))

THEOREM: ok-mg-statep-preserved-call-case
((`'proc-call-mg` = car (*stmt*))
 ∧ ok-mg-statep (*mg-state*, *r-cond-list*)
 ∧ ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*)
 ∧ signatures-match (mg-alist (*mg-state*), *name-alist*)
 ∧ ok-mg-def-plistp (*proc-list*))
→ ok-mg-statep (make-call-environment (*mg-state*,
 stmt,
 fetch-called-def (*stmt*, *proc-list*)),
 make-cond-list (fetch-called-def (*stmt*, *proc-list*)))

THEOREM: call-formal-signatures-match
signatures-match (make-alist-from-formals (*formals*),
 make-call-param-alist (*formals*, *actuals*, *mg-alist*))

THEOREM: call-formal-signatures-match2
signatures-match (make-call-param-alist (*formals*, *actuals*, *mg-alist*),
 make-alist-from-formals (*formals*))

THEOREM: call-local-signatures-match
plistp (*locals*) → signatures-match (make-alist-from-formals (*locals*), *locals*)

THEOREM: call-local-signatures-match2
signatures-match (*locals*, make-alist-from-formals (*locals*))

THEOREM: call-signatures-match1
plistp (def-locals (*def*))
→ signatures-match (make-name-alist (*def*),
make-call-var-alist (*mg-alist*, *stmt*, *def*))

THEOREM: call-signatures-match2
signatures-match (make-call-var-alist (*mg-alist*, *stmt*, *def*), make-name-alist (*def*))

THEOREM: call-signatures-match3
signatures-match (mg-alist (make-call-environment (*mg-state*,
stmt,
fetch-called-def (*stmt*,
proc-list))),
make-name-alist (fetch-called-def (*stmt*, *proc-list*)))

DEFINITION:
formal-types-preserved (*formals*, *alist*)
= **if** *formals* ≃ **nil then t**
else definedp (caar (*formals*), *alist*)
∧ (cadar (*formals*)
= cadr (assoc (caar (*formals*), *alist*)))
∧ formal-types-preserved (cdr (*formals*), *alist*) **endif**

THEOREM: formal-types-preserved-append
formal-types-preserved (*formals*, *lst1*)
→ formal-types-preserved (*formals*, append (*lst1*, *lst2*))

THEOREM: formal-types-unaffected-by-extra-binding
(car (*x*) ∉ listcars (*y*))
→ (formal-types-preserved (*y*, cons (*x*, *z*)) = formal-types-preserved (*y*, *z*))

THEOREM: formal-types-preserved-in-call-param-alist
(ok-mg-formal-data-params-plistp (*formals*)
∧ no-duplicates (listcars (*formals*)))
→ formal-types-preserved (*formals*,
make-call-param-alist (*formals*, *actuals*, *mg-alist*))

THEOREM: formal-types-preserved-in-call-environment
((`'proc-call-mg` = car (*stmt*))
∧ ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*)
∧ ok-mg-def-plistp (*proc-list*)
∧ mg-alistp (mg-alist (*mg-state*)))

→ formal-types-preserved (def-formals (fetch-called-def (*stmt*, *proc-list*)),
mg-alist (make-call-environment (*mg-state*,
stmt,
fetch-called-def (*stmt*,
proc-list))))

THEOREM: copy-out-params-preserves-mg-alistp
(mg-alistp (*old-alist*)
∧ mg-alistp (*new-alist*)
∧ formal-types-preserved (*formals*, *new-alist*)
∧ data-param-lists-match (*actuals*, *formals*, *name-alist*)
∧ signatures-match (*old-alist*, *name-alist*)
∧ ok-mg-formal-data-params-plistp (*formals*))
→ mg-alistp (copy-out-params (*formals*, *actuals*, *new-alist*, *old-alist*))

THEOREM: formal-types-preserved-in-matching-signatures
(mg-name-alistp (*old-alist*)
∧ formal-types-preserved (*formals*, *old-alist*)
∧ signatures-match (*old-alist*, *new-alist*))
→ formal-types-preserved (*formals*, *new-alist*)

```
;; This is the case needed for map-call-effects-preserves-ok-state.
```

THEOREM: map-call-effects-preserves-mg-alistp
((`'proc-call-mg` = car (*stmt*))
∧ mg-alistp (mg-alist (*mg-state*))
∧ ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*)
∧ signatures-match (mg-alist (*mg-state*), *name-alist*)
∧ ok-mg-def-plistp (*proc-list*)
∧ ok-mg-def (fetch-called-def (*stmt*, *proc-list*), *proc-list*)
∧ mg-alistp (mg-alist (mg-meaning (def-body (fetch-called-def (*stmt*,
proc-list)),
proc-list,
make-call-environment (*mg-state*,
stmt,
fetch-called-def (*stmt*,
proc-list)),
$n - 1$))))
→ mg-alistp (copy-out-params (def-formals (fetch-called-def (*stmt*, *proc-list*)),
call-actuals (*stmt*),
mg-alist (mg-meaning (def-body (fetch-called-def (*stmt*,
proc-list)),
proc-list,
make-call-environment (*mg-state*,

stmt,
fetch-called-def (*stmt*,
proc-list)),
$n - 1$)),
mg-alist (*mg-state*)))

THEOREM: convert-condition1-membership
(length (*def-conds*) = length (*call-conds*))
→ (convert-condition1 (*cc*, *def-conds*, *call-conds*)
∈ cons (`'routineerror`, *call-conds*))

THEOREM: cond-identifier-plistp-preserves-membership
((*cc* ∈ *lst1*)
∧ cond-identifier-plistp (*lst1*, *lst2*)
∧ (*cc* ≠ `'routineerror`))
→ (*cc* ∈ *lst2*)

EVENT: Disable cond-identifier-plistp-preserves-membership.

THEOREM: cons-preserves-membership
((*x* ∈ cons (*y*, *z*)) ∧ (*x* ≠ *y*)) → (*x* ∈ *z*)

EVENT: Disable cons-preserves-membership.

THEOREM: cond-identifier-conversion-litatom
((length (*def-conds*) = length (*call-conds*))
∧ cond-identifier-plistp (*call-conds*, *cond-list*))
→ litatom (convert-condition1 (*cc*, *def-conds*, *call-conds*))

EVENT: Disable cond-identifier-conversion-litatom.

THEOREM: map-call-effects-preserves-ok-state
((`'proc-call-mg` = car (*stmt*))
∧ ok-mg-statep (*mg-state*, *r-cond-list*)
∧ ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*)
∧ signatures-match (mg-alist (*mg-state*), *name-alist*)
∧ ok-mg-def-plistp (*proc-list*)
∧ ok-mg-statep (mg-meaning (def-body (fetch-called-def (*stmt*, *proc-list*)),
proc-list,
make-call-environment (*mg-state*,
stmt,
fetch-called-def (*stmt*,
proc-list)),

$n - 1$),
make-cond-list (fetch-called-def (*stmt*, *proc-list*))))
→ ok-mg-statep (map-call-effects (mg-meaning (def-body (fetch-called-def (*stmt*,
proc-list)),
proc-list,
make-call-environment (*mg-state*,
stmt,
fetch-called-def (*stmt*,
proc-list)),
$n - 1$),
fetch-called-def (*stmt*, *proc-list*),
stmt,
mg-state),
r-cond-list)

THEOREM: simple-typed-literalp-ok-valuep
simple-typed-literalp (*exp*, *type*) → ok-mg-valuep (*exp*, *type*)

EVENT: Enable ok-predefined-proc-call.

EVENT: Enable ok-predefined-proc-args.

EVENT: Enable ok-mg-statement.

EVENT: Enable simple-identifierp-implies-definedp.

EVENT: Enable int-identifierp-implies-definedp.

EVENT: Enable boolean-identifierp-implies-definedp.

EVENT: Enable character-identifierp-implies-definedp.

EVENT: Enable boolean-identifierp.

EVENT: Enable character-identifierp.

EVENT: Enable int-identifierp.

EVENT: Disable mg-bool.

EVENT: Disable mg-or-bool.

EVENT: Disable mg-and-bool.

EVENT: Disable mg-not-bool.

THEOREM: tag-length-plistp
length-plistp (tag (*x*, *y*), 2)

THEOREM: car-tag
(car (tag (*x*, *y*)) = *x*) ∧ (cdr (tag (*x*, *y*)) = list (*y*))

THEOREM: simple-typed-literalp-boolean-literals
simple-typed-literalp (tag (`'boolean-mg`, `'true-mg`), `'boolean-mg`)
∧ simple-typed-literalp (tag (`'boolean-mg`, `'false-mg`),
`'boolean-mg`)

THEOREM: simple-typed-literalp-boolean-mg-bool
simple-typed-literalp (mg-bool (*x*), `'boolean-mg`)

THEOREM: simple-typed-literalp-boolean-mg-bool-not
simple-typed-literalp (tag (`'boolean-mg`, mg-not-bool (*x*)), `'boolean-mg`)

THEOREM: ok-mg-valuep-int-mg
ok-mg-valuep (tag (`'int-mg`, *x*), `'int-mg`)
= small-integerp (*x*, MG-WORD-SIZE)

THEOREM: boolean-literalp-tag-untag
boolean-literalp (*x*) → boolean-literalp (tag (`'boolean-mg`, untag (*x*)))

THEOREM: boolean-identifier-boolean-literal-value
(mg-alistp (*mg-alist*) ∧ (cadr (assoc (*x*, *mg-alist*)) = `'boolean-mg`))
→ boolean-literalp (caddr (assoc (*x*, *mg-alist*)))

THEOREM: boolean-identifier-boolean-literalp
(boolean-identifierp (*x*, *name-alist*)
∧ mg-alistp (*mg-alist*)
∧ signatures-match (*mg-alist*, *name-alist*))
→ boolean-literalp (caddr (assoc (*x*, *mg-alist*)))

THEOREM: mg-meaning-mg-simple-variable-assignment-preserves-ok-mg-statep
((car (*stmt*) = `'predefined-proc-call-mg`)
∧ (call-name (*stmt*) = `'mg-simple-variable-assignment`)
∧ ok-mg-statep (*mg-state*, *r-cond-list*)

∧ ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*)
∧ signatures-match (mg-alist (*mg-state*), *name-alist*))
→ ok-mg-statep (mg-meaning-predefined-proc-call (*stmt*, *mg-state*),
r-cond-list)

THEOREM: mg-meaning-mg-simple-constant-assignment-preserves-ok-mg-statep
((car (*stmt*) = `'predefined-proc-call-mg`)
∧ (call-name (*stmt*) = `'mg-simple-constant-assignment`)
∧ ok-mg-statep (*mg-state*, *r-cond-list*)
∧ ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*)
∧ signatures-match (mg-alist (*mg-state*), *name-alist*))
→ ok-mg-statep (mg-meaning-predefined-proc-call (*stmt*, *mg-state*),
r-cond-list)

THEOREM: mg-meaning-mg-simple-variable-eq-preserves-ok-mg-statep
((car (*stmt*) = `'predefined-proc-call-mg`)
∧ (call-name (*stmt*) = `'mg-simple-variable-eq`)
∧ ok-mg-statep (*mg-state*, *r-cond-list*)
∧ ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*)
∧ signatures-match (mg-alist (*mg-state*), *name-alist*))
→ ok-mg-statep (mg-meaning-predefined-proc-call (*stmt*, *mg-state*),
r-cond-list)

THEOREM: mg-meaning-mg-simple-constant-eq-preserves-ok-mg-statep
((car (*stmt*) = `'predefined-proc-call-mg`)
∧ (call-name (*stmt*) = `'mg-simple-constant-eq`)
∧ ok-mg-statep (*mg-state*, *r-cond-list*)
∧ ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*)
∧ signatures-match (mg-alist (*mg-state*), *name-alist*))
→ ok-mg-statep (mg-meaning-predefined-proc-call (*stmt*, *mg-state*),
r-cond-list)

THEOREM: mg-meaning-mg-integer-le-preserves-ok-mg-statep
((car (*stmt*) = `'predefined-proc-call-mg`)
∧ (call-name (*stmt*) = `'mg-integer-le`)
∧ ok-mg-statep (*mg-state*, *r-cond-list*)
∧ ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*)
∧ signatures-match (mg-alist (*mg-state*), *name-alist*))
→ ok-mg-statep (mg-meaning-predefined-proc-call (*stmt*, *mg-state*),
r-cond-list)

THEOREM: mg-meaning-mg-integer-unary-minus-preserves-ok-mg-statep
((car (*stmt*) = `'predefined-proc-call-mg`)
∧ (call-name (*stmt*) = `'mg-integer-unary-minus`)
∧ ok-mg-statep (*mg-state*, *r-cond-list*)

$\wedge$ ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*)
$\wedge$ signatures-match (mg-alist (*mg-state*), *name-alist*))
$\rightarrow$ ok-mg-statep (mg-meaning-predefined-proc-call (*stmt*, *mg-state*),
r-cond-list)

THEOREM: mg-meaning-mg-integer-add-preserves-ok-mg-statep
((car (*stmt*) = `'predefined-proc-call-mg`)
$\wedge$ (call-name (*stmt*) = `'mg-integer-add`)
$\wedge$ ok-mg-statep (*mg-state*, *r-cond-list*)
$\wedge$ ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*)
$\wedge$ signatures-match (mg-alist (*mg-state*), *name-alist*))
$\rightarrow$ ok-mg-statep (mg-meaning-predefined-proc-call (*stmt*, *mg-state*),
r-cond-list)

THEOREM: mg-meaning-mg-integer-subtract-preserves-ok-mg-statep
((car (*stmt*) = `'predefined-proc-call-mg`)
$\wedge$ (call-name (*stmt*) = `'mg-integer-subtract`)
$\wedge$ ok-mg-statep (*mg-state*, *r-cond-list*)
$\wedge$ ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*)
$\wedge$ signatures-match (mg-alist (*mg-state*), *name-alist*))
$\rightarrow$ ok-mg-statep (mg-meaning-predefined-proc-call (*stmt*, *mg-state*),
r-cond-list)

THEOREM: mg-meaning-mg-boolean-or-preserves-ok-mg-statep
((car (*stmt*) = `'predefined-proc-call-mg`)
$\wedge$ (call-name (*stmt*) = `'mg-boolean-or`)
$\wedge$ ok-mg-statep (*mg-state*, *r-cond-list*)
$\wedge$ ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*)
$\wedge$ signatures-match (mg-alist (*mg-state*), *name-alist*))
$\rightarrow$ ok-mg-statep (mg-meaning-predefined-proc-call (*stmt*, *mg-state*),
r-cond-list)

THEOREM: mg-meaning-mg-boolean-and-preserves-ok-mg-statep
((car (*stmt*) = `'predefined-proc-call-mg`)
$\wedge$ (call-name (*stmt*) = `'mg-boolean-and`)
$\wedge$ ok-mg-statep (*mg-state*, *r-cond-list*)
$\wedge$ ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*)
$\wedge$ signatures-match (mg-alist (*mg-state*), *name-alist*))
$\rightarrow$ ok-mg-statep (mg-meaning-predefined-proc-call (*stmt*, *mg-state*),
r-cond-list)

THEOREM: mg-meaning-mg-boolean-not-preserves-ok-mg-statep
((car (*stmt*) = `'predefined-proc-call-mg`)
$\wedge$ (call-name (*stmt*) = `'mg-boolean-not`)
$\wedge$ ok-mg-statep (*mg-state*, *r-cond-list*)

$\wedge$ ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*)
$\wedge$ signatures-match (mg-alist (*mg-state*), *name-alist*))
$\rightarrow$ ok-mg-statep (mg-meaning-predefined-proc-call (*stmt*, *mg-state*),
r-cond-list)

THEOREM: simple-typed-literal-list-elements
(simple-typed-literal-plistp (*lst*, *type*) $\wedge$ ($i <$ length (*lst*)))
$\rightarrow$ simple-typed-literalp (get (*i*, *lst*), *type*)

THEOREM: index-array-mg-alist-elementp
(mg-alistp (*mg-alist*)
$\wedge$ array-identifierp (*a*, *mg-alist*)
$\wedge$ ($i <$ array-length (cadr (assoc (*a*, *mg-alist*)))))
$\rightarrow$ simple-typed-literalp (get (*i*, caddr (assoc (*a*, *mg-alist*))),
array-elemtype (cadr (assoc (*a*, *mg-alist*))))

THEOREM: put-preserves-simple-typed-literal-plistp
(simple-typed-literal-plistp (*lst*, *type*)
$\wedge$ ($i <$ length (*lst*))
$\wedge$ simple-typed-literalp (*val*, *type*))
$\rightarrow$ simple-typed-literal-plistp (put (*val*, *i*, *lst*), *type*)

THEOREM: simple-type-literal-plistp-ok-valuep
(array-mg-type-refp (*type*)
$\wedge$ simple-typed-literal-plistp (*exp*, array-elemtype (*type*))
$\wedge$ (length (*exp*) = array-length (*type*)))
$\rightarrow$ ok-mg-valuep (*exp*, *type*)

THEOREM: mg-meaning-mg-index-array-preserves-ok-mg-statep
((car (*stmt*) = `'predefined-proc-call-mg`)
$\wedge$ (call-name (*stmt*) = `'mg-index-array`)
$\wedge$ ok-mg-statep (*mg-state*, *r-cond-list*)
$\wedge$ ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*)
$\wedge$ signatures-match (mg-alist (*mg-state*), *name-alist*))
$\rightarrow$ ok-mg-statep (mg-meaning-predefined-proc-call (*stmt*, *mg-state*),
r-cond-list)

THEOREM: array-identifiers-have-array-types
(array-identifierp (*x*, *mg-alist*) $\wedge$ mg-alistp (*mg-alist*))
$\rightarrow$ array-mg-type-refp (cadr (assoc (*x*, *mg-alist*)))

THEOREM: array-identifiers-have-array-types2
(array-identifierp (*a*, *mg-alist*) $\wedge$ mg-alistp (*mg-alist*))
$\rightarrow$ simple-typed-literal-plistp (caddr (assoc (*a*, *mg-alist*)),
array-elemtype (cadr (assoc (*a*, *mg-alist*))))

THEOREM: array-identifier-lengths-match
(array-identifierp (*a*, *mg-alist*) ∧ mg-alistp (*mg-alist*))
→ (length (caddr (assoc (*a*, *mg-alist*)))
= array-length (cadr (assoc (*a*, *mg-alist*))))

THEOREM: simple-typed-identifier-has-simple-typed-literal-value
(mg-alistp (*mg-alist*) ∧ simple-typed-identifierp (*x*, *type*, *mg-alist*))
→ simple-typed-literalp (caddr (assoc (*x*, *mg-alist*)), *type*)

THEOREM: mg-meaning-mg-array-element-assignment-preserves-ok-mg-statep
((car (*stmt*) = `'predefined-proc-call-mg`)
∧ (call-name (*stmt*) = `'mg-array-element-assignment`)
∧ ok-mg-statep (*mg-state*, *r-cond-list*)
∧ ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*)
∧ signatures-match (mg-alist (*mg-state*), *name-alist*))
→ ok-mg-statep (mg-meaning-predefined-proc-call (*stmt*, *mg-state*),
r-cond-list)

THEOREM: mg-meaning-predefined-proc-call-preserves-ok-mg-statep
((car (*stmt*) = `'predefined-proc-call-mg`)
∧ ok-mg-statep (*mg-state*, *r-cond-list*)
∧ ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*)
∧ signatures-match (mg-alist (*mg-state*), *name-alist*))
→ ok-mg-statep (mg-meaning-predefined-proc-call (*stmt*, *mg-state*),
r-cond-list)

THEOREM: set-condition-normal-preserves-ok-mg-statep
ok-mg-statep (*state*, *cond-list1*)
→ ok-mg-statep (set-condition (*state*, `'normal`), *cond-list2*)

THEOREM: append-conditions-preserves-ok-mg-statep
(ok-mg-statep (*state*, append (*lst*, *lst2*)) ∧ (cc (*state*) ∉ *lst*))
→ ok-mg-statep (*state*, *lst2*)

DEFINITION:
meaning-induction-hint (*stmt*, *proc-list*, *mg-state*, *n*, *name-alist*, *r-cond-list*)
= **if** *n* ≃ `0` **then t**
elseif ¬ normal (*mg-state*) **then t**
elseif `'no-op-mg` = car (*stmt*) **then t**
elseif `'signal-mg` = car (*stmt*) **then t**
elseif `'prog2-mg` = car (*stmt*)
then meaning-induction-hint (prog2-left-branch (*stmt*),
proc-list,
mg-state,
n − 1,

name-alist,
r-cond-list)
∧ meaning-induction-hint (prog2-right-branch (*stmt*),
proc-list,
mg-meaning (prog2-left-branch (*stmt*),
proc-list,
mg-state,
$n - 1$),
$n - 1$,
name-alist,
r-cond-list)
elseif `'loop-mg` = car (*stmt*)
then meaning-induction-hint (loop-body (*stmt*),
proc-list,
mg-state,
$n - 1$,
name-alist,
cons (`'leave`, *r-cond-list*))
∧ meaning-induction-hint (*stmt*,
proc-list,
mg-meaning (loop-body (*stmt*),
proc-list,
mg-state,
$n - 1$),
$n - 1$,
name-alist,
r-cond-list)
elseif `'if-mg` = car (*stmt*)
then meaning-induction-hint (if-false-branch (*stmt*),
proc-list,
mg-state,
$n - 1$,
name-alist,
r-cond-list)
∧ meaning-induction-hint (if-true-branch (*stmt*),
proc-list,
mg-state,
$n - 1$,
name-alist,
r-cond-list)
elseif `'begin-mg` = car (*stmt*)
then meaning-induction-hint (begin-body (*stmt*),
proc-list,
mg-state,

$n - 1$,
name-alist,
append (when-labels (*stmt*), *r-cond-list*))
$\wedge$ meaning-induction-hint (when-handler (*stmt*),
proc-list,
set-condition (mg-meaning (begin-body (*stmt*),
proc-list,
mg-state,
$n - 1$),
`'normal`),
$n - 1$,
name-alist,
r-cond-list)
elseif `'proc-call-mg` = car (*stmt*)
then meaning-induction-hint (def-body (fetch-called-def (*stmt*, *proc-list*)),
proc-list,
make-call-environment (*mg-state*,
stmt,
fetch-called-def (*stmt*,
proc-list)),
$n - 1$,
make-name-alist (fetch-called-def (*stmt*,
proc-list)),
make-cond-list (fetch-called-def (*stmt*,
proc-list)))
elseif `'predefined-proc-call-mg` = car (*stmt*) **then t**
else f endif

EVENT: Disable ok-predefined-proc-call.

EVENT: Disable ok-predefined-proc-args.

EVENT: Disable ok-mg-statement.

EVENT: Disable simple-identifierp-implies-definedp.

EVENT: Disable int-identifierp-implies-definedp.

EVENT: Disable boolean-identifierp-implies-definedp.

EVENT: Disable character-identifierp-implies-definedp.

EVENT: Disable boolean-identifierp.

EVENT: Disable character-identifierp.

EVENT: Disable int-identifierp.

EVENT: Disable signatures-match-preserves-ok-mg-statement.

EVENT: Disable signatures-match-preserves-ok-predefined-proc-call.

EVENT: Disable signatures-match-preserves-ok-predefined-proc-args.

EVENT: Disable signatures-match-preserves-simple-identifierp.

EVENT: Disable not-member-listcars-not-assoc.

EVENT: Disable ok-mg-statep.

EVENT: Disable call-locals-alist-mg-alistp.

THEOREM: removing-condition-preserves-ok-mg-statep
(ok-mg-statep (*state*, cons (*x*, *lst*)) $\wedge$ (cc (*state*) $\neq$ *x*))
$\rightarrow$ ok-mg-statep (*state*, *lst*)

EVENT: Disable removing-condition-preserves-ok-mg-statep.

THEOREM: adding-condition-preserves-ok-mg-statep
ok-mg-statep (*state*, *lst2*) $\rightarrow$ ok-mg-statep (*state*, append (*lst*, *lst2*))

EVENT: Disable adding-condition-preserves-ok-mg-statep.

THEOREM: mg-meaning-preserves-ok-mg-statep
(ok-mg-statep (*mg-state*, *r-cond-list*)
$\wedge$ ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*)
$\wedge$ signatures-match (mg-alist (*mg-state*), *name-alist*)
$\wedge$ ok-mg-def-plistp (*proc-list*))
$\rightarrow$ ok-mg-statep (mg-meaning (*stmt*, *proc-list*, *mg-state*, *n*), *r-cond-list*)

EVENT: Disable mg-meaning-preserves-ok-mg-statep.

THEOREM: mg-meaning-preserves-ok-cc
(ok-mg-statep (*mg-state*, *r-cond-list*)
 ∧ ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*)
 ∧ signatures-match (mg-alist (*mg-state*), *name-alist*)
 ∧ ok-mg-def-plistp (*proc-list*))
 → ok-cc (cc (mg-meaning (*stmt*, *proc-list*, *mg-state*, *n*)), *r-cond-list*)

EVENT: Disable mg-meaning-preserves-ok-cc.

THEOREM: mg-meaning-preserves-mg-alistp
(ok-mg-statep (*mg-state*, *r-cond-list*)
 ∧ ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*)
 ∧ signatures-match (mg-alist (*mg-state*), *name-alist*)
 ∧ ok-mg-def-plistp (*proc-list*))
 → mg-alistp (mg-alist (mg-meaning (*stmt*, *proc-list*, *mg-state*, *n*)))

EVENT: Disable mg-meaning-preserves-mg-alistp.

THEOREM: mg-meaning-condition-member-cond-list1
(ok-mg-statep (*mg-state*, *r-cond-list*)
 ∧ ok-mg-statement (*stmt*, *r-cond-list*, *name-alist*, *proc-list*)
 ∧ ok-mg-def-plistp (*proc-list*)
 ∧ signatures-match (mg-alist (*mg-state*), *name-alist*)
 ∧ (¬ resource-errorp (mg-meaning (*stmt*, *proc-list*, *mg-state*, *n*))))
 → (cc (mg-meaning (*stmt*, *proc-list*, *mg-state*, *n*))
  ∈ cons (`'normal`, cons (`'routineerror`, *r-cond-list*)))

EVENT: Disable mg-meaning-condition-member-cond-list1.

EVENT: Make the library `"c3"`.

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