Notes are taken from Tibshirani's lecture notes: https://www.stat.cmu.edu/~ryantibs/convexoptF13/scribes/

## Convex Function

Definition 4.28 A function $f: \mathbb{R}^{n} \rightarrow \mathbb{R}$ is convex if $\operatorname{dom} f$ is a convex set and if for all $x, y \in \operatorname{dom} f$, and $\theta$ with $0 \leq \theta \leq 1$, we have

$$
f(\theta x+(1-\theta) y) \leq \theta f(x)+(1-\theta) f(y) .
$$

Definition 4.29 A function $f$ is strictly convex if whenever $x \neq y$, and $0<\theta<1$, strict inequality holds, that is, we have

$$
f(\theta x+(1-\theta) y)<\theta f(x)+(1-\theta) f(y) .
$$



$$
f(y)>=f(x)+\backslash \operatorname{grad} f(x)^{\wedge} T(y-x)
$$

## Strong Convexity

Definition 4.32 $A$ differentiable function $f$ is called $m$-strongly convex if $m>0$ and

$$
(\nabla f(x)-\nabla f(y))^{T}(x-y) \geq m\|x-y\|_{2}^{2}, \forall x, y \in \operatorname{dom} f
$$

An equivalent condition is

$$
f(y) \geq f(x)+\nabla f(x)^{T}(y-x)+\frac{m}{2}\|y-x\|_{2}^{2}, \forall x, y \in \operatorname{dom} f
$$

It is not necessary for a function to be differentiable. We could have the definition without gradient.

Definition 4.33 $A$ function $f$ is called $m$-strongly convex if $m>0$ and for $0 \leq t \leq 1$

$$
f(t x+(1-t) y) \leq t f(x)+(1-t) f(y)-\frac{1}{2} m t(1-t)\|x-y\|_{2}^{2}, \forall x, y \in \operatorname{dom} f
$$

If the function is twice continuously differentiable, we could have the definition with Hessian matrix.

Definition $4.34 f$ is called $m$-stronly convex if $m>0$ and

$$
\nabla^{2} f(x) \geq m I, \forall x, y \in \operatorname{dom} f
$$

A strongly convex function is also strictly convex, but not vice-versa.

## Extended-Value Extension

Definition $4.37 \tilde{f}: \mathbb{R}^{n} \rightarrow \mathbb{R} \bigcup\{\infty\}$ is extended-value extension of $f$ :
$\tilde{f}(x)=\left\{\begin{array}{cc}f(x) & x \in \operatorname{dom} f \\ \infty & x \notin \operatorname{dom} f\end{array}\right.$
The extension $\tilde{f}$ is defined on all $\mathbb{R}^{n}$, and takes values in $\mathbb{R} \bigcup\{\infty\}$. This does not change its convexity
Theorem $4.38 f$ is convex
$\Leftrightarrow \tilde{f}$ is convex
$\Leftrightarrow \tilde{f}(\theta x+(1-\theta) y) \leq \theta \tilde{f}(x)+(1-\theta) \tilde{f}(y), 0 \leq \theta \leq 1$

## Properties of Convex Functions

Let $f$ be a differentiable function, $\operatorname{dom} f$ is open and convex, then we have
$f$ is convex $\Leftrightarrow f(y) \geq f(x)+\nabla f(x)^{T}(y-x)$
The inequality states that for a convex function, the first-order Taylor
approximation is a global underestimator of the function. Conversely, if the first-order Taylor approximation
of a function is always a global underestimator of the function, then the function is convex.

Let $f$ be twice differentiable, $\operatorname{dom} f$ is open, then we have
$f$ is convex $\Leftrightarrow \nabla^{2} f(x) \geq 0, \forall x \in \operatorname{dom} f$
If $\nabla^{2} f(x)>0, \forall x \in \operatorname{dom} f, f$ is strictly convex. The converse is not true.
For example, the function $f: \mathbb{R} \rightarrow \mathbb{R}$ given by $f(x)=x^{4}$ is strictly convex but has zero second derivative at $x=0$

## Gradient Descent

Recall that we have $f: \mathbb{R}^{n} \rightarrow \mathbb{R}$, convex and differentiable. We want to solve

$$
\min _{x \in \mathbb{R}^{n}} f(x)
$$

i.e, to find $x^{\star}$ such that $f\left(x^{\star}\right)=\min f(x)$.

Gradient descent: choose initial $x^{(0)} \in \mathbb{R}^{n}$, repeat :

$$
x^{(k)}=x^{(k-1)}-t_{k} \cdot \nabla f\left(x^{(k-1)}\right), k=1,2,3, \ldots
$$

Stop at some point(When to stop is quite dependent on what problems you are looking at).

## Coordinate Descent

Similar but coordinate-by-coordinate by picking the coordinate with maximum gradient

## Step Size

Fixed

## Backtracking Line Search

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Starting with a maximum candidate step size value 的>0, using search control parameters }\tau\in(0,1)\mathrm{ and c}\in(0,1)\mathrm{ , the backtracking line search algorithm can be expressed as follows
    1. Set }t=-cm\mathrm{ and iteration counter }j=0
    2. Until the condition is satisfied that f(\mathbf{x})-f(\mathbf{x}+\mp@subsup{\alpha}{j}{}\mathbf{p})\geq\mp@subsup{\alpha}{j}{}t\mathrm{ , repeatedly increment j and set }\mp@subsup{\alpha}{j}{}=\tau\mp@subsup{\alpha}{j-1}{}
    3. Return }\mp@subsup{\alpha}{j}{}\mathrm{ as the solution.
In other words, reduce }\mp@subsup{\alpha}{0}{}\mathrm{ by a factor of }\tau\mathrm{ in each iteration until the Armijo-Goldstein condition is fulfilled.
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Define the local slope of the function of $\alpha$ along the search direction $\mathbf{p}$ as $m=\nabla f(\mathbf{x})^{\mathrm{T}} \mathbf{p}$. It is assumed that $\mathbf{p}$ is a unit vector in a direction in which some local decrease is possible, i.e., it is assumed that $m<0$.
 objective function. The condition is fulfilled if $f(\mathbf{x}+\alpha \mathbf{p}) \leq f(\mathbf{x})+\alpha c m$.


## Exact Line Search

At each iteration, do the best we can along the direction of the gradient,

$$
t=\underset{s \geq 0}{\operatorname{argmin}} f(x-s \nabla f(x)) .
$$

Usually, it is not possible to do this minimization exactly.
Approximations to exact line search are often not much more efficient than backtracking, and it's not worth it.

## Proof (first inequality is Lagrange form of Taylor's theorem)

Theorem 6.1 Suppose the function $f: \mathbb{R}^{n} \rightarrow \mathbb{R}$ is convex and differentiable, and that its gradient is Lipschitz continuous with constant $L>0$, i.e. we have that $\|\nabla f(x)-\nabla f(y)\|_{2} \leq L\|x-y\|_{2}$ for any $x, y$. Then if we run gradient descent for $k$ iterations with a fixed step size $t \leq 1 / L$, it will yield a solution $f(k)$ which satisfies

$$
\begin{equation*}
f\left(x^{(k)}\right)-f\left(x^{*}\right) \leq \frac{\left\|x^{(0)}-x^{*}\right\|_{2}^{2}}{2 t k}, \tag{6.1}
\end{equation*}
$$

where $f\left(x^{*}\right)$ is the optimal value. Intuitively, this means that gradient descent is guaranteed to converge and that it converges with rate $O(1 / k)$.

Proof: Our assumption that $\nabla f$ is Lipschitz continuous with constant $L$ implies that $\nabla^{2} f(x) \preceq L I$, or equivalently that $\nabla^{2} f(x)-L I$ is a negative semidefinite matrix. Using this fact, we can perform a quadratic expansion of $f$ around $f(x)$ and obtain the following inequality:

$$
\begin{aligned}
f(y) & \leq f(x)+\nabla f(x)^{T}(y-x)+\frac{1}{2} \nabla^{2} f(x)\|y-x\|_{2}^{2} \\
& \leq f(x)+\nabla f(x)^{T}(y-x)+\frac{1}{2} L\|y-x\|_{2}^{2}
\end{aligned}
$$

Now let's plug in the gradient descent update by letting $y=x^{+}=x-t \nabla f(x)$. We then get:

$$
\begin{align*}
f\left(x^{+}\right) & \leq f(x)+\nabla f(x)^{T}\left(x^{+}-x\right)+\frac{1}{2} L\left\|x^{+}-x\right\|_{2}^{2} \\
& =f(x)+\nabla f(x)^{T}(x-t \nabla f(x)-x)+\frac{1}{2} L\|x-t \nabla f(x)-x\|_{2}^{2} \\
& =f(x)-\nabla f(x)^{T} t \nabla f(x)+\frac{1}{2} L\|t \nabla f(x)\|_{2}^{2} \\
& =f(x)-t\|\nabla f(x)\|_{2}^{2}+\frac{1}{2} L t^{2}\|\nabla f(x)\|_{2}^{2} \\
& =f(x)-\left(1-\frac{1}{2} L t\right) t\|\nabla f(x)\|_{2}^{2} \tag{6.2}
\end{align*}
$$

Using $t \leq 1 / L$, we know that $-\left(1-\frac{1}{2} L t\right)=\frac{1}{2} L t-1 \leq \frac{1}{2} L(1 / L)-1=\frac{1}{2}-1=-\frac{1}{2}$. Plugging this in to ??6. 2
we can conclude the following:

$$
\begin{equation*}
f\left(x^{+}\right) \leq f(x)-\frac{1}{2} t\|\nabla f(x)\|_{2}^{2} \tag{6.3}
\end{equation*}
$$

Since $\frac{1}{2} t\|\nabla f(x)\|_{2}^{2}$ will always be positive unless $\nabla f(x)=0$, this inequality implies that the objective function value strictly decreases with each iteration of gradient descent until it reaches the optimal value $f(x)=f\left(x^{*}\right)$. Note that this convergence result only holds when we choose $t$ to be small enough, i.e. $t \leq 1 / L$. This explains why we observe in practice that gradient descent diverges when the step size is too large.
Next, we can bound $f\left(x^{+}\right)$, the objective value at the next iteration, in terms of $f\left(x^{*}\right)$, the optimal objective value. Since $f$ is convex, we can write

$$
\begin{aligned}
& f\left(x^{*}\right) \geq f(x)+\nabla f(x)^{T}\left(x^{*}-x\right) \\
& f(x) \leq f\left(x^{*}\right)+\nabla f(x)^{T}\left(x-x^{*}\right)
\end{aligned}
$$

where the first inequality yields the second through simple rearrangement of terms. Plugging this in to we obtain

$$
\begin{align*}
& f\left(x^{+}\right) \leq f\left(x^{*}\right)+\nabla f(x)^{T}\left(x-x^{*}\right)-\frac{t}{2}\|\nabla f(x)\|_{2}^{2} \\
& f\left(x^{+}\right)-f\left(x^{*}\right) \leq \frac{1}{2 t}\left(2 t \nabla f(x)^{T}\left(x-x^{*}\right)-t^{2}\|\nabla f(x)\|_{2}^{2}\right) \\
& f\left(x^{+}\right)-f\left(x^{*}\right) \leq \frac{1}{2 t}\left(2 t \nabla f(x)^{T}\left(x-x^{*}\right)-t^{2}\|\nabla f(x)\|_{2}^{2}-\left\|x-x^{*}\right\|_{2}^{2}+\left\|x-x^{*}\right\|_{2}^{2}\right) \\
& f\left(x^{+}\right)-f\left(x^{*}\right) \leq \frac{1}{2 t}\left(\left\|x-x^{*}\right\|_{2}^{2}-\left\|x-t \nabla f(x)-x^{*}\right\|_{2}^{2}\right) \tag{6.4}
\end{align*}
$$

where the final inequality is obtained by observing that expanding the square of $\left\|x-t \nabla f(x)-x^{*}\right\|_{2}^{2}$ yields $\left\|x-x^{*}\right\|_{2}^{2}-2 t \nabla f(x)^{T}\left(x-x^{*}\right)+t^{2}\|\nabla f(x)\|_{2}^{2}$. Notice that by definition we have $x^{+}=x-t \nabla f(x)$. Plugging this in to (?? 6.4 ) yiflds:

$$
\begin{equation*}
f\left(x^{+}\right)-f\left(x^{*}\right) \leq \frac{1}{2 t}\left(\left\|x-x^{*}\right\|_{2}^{2}-\left\|x^{+}-x^{*}\right\|_{2}^{2}\right) \tag{6.5}
\end{equation*}
$$

This inequality holds for $x^{+}$on every iteration of gradient descent. Summing over iterations, we get:

$$
\begin{align*}
\sum_{i=1}^{k} f\left(x^{(i)}-f\left(x^{*}\right)\right. & \leq \sum_{i=1}^{k} \frac{1}{2 t}\left(\left\|x^{(i-1)}-x^{*}\right\|_{2}^{2}-\left\|x^{(i)}-x^{*}\right\|_{2}^{2}\right) \\
& =\frac{1}{2 t}\left(\left\|x^{(0)}-x^{*}\right\|_{2}^{2}-\left\|x^{(k)}-x^{*}\right\|_{2}^{2}\right) \\
& \leq \frac{1}{2 t}\left(\left\|x^{(0)}-x^{*}\right\|_{2}^{2}\right) \tag{6.6}
\end{align*}
$$

where the summation on the right-hand side disappears because it is a telescoping sum. Finally, using the fact that $f$ decreasing on every iteration, we can conclude that

$$
\begin{align*}
f\left(x^{(k)}\right)-f\left(x^{*}\right) & \leq \frac{1}{k} \sum_{i=1}^{k} f\left(x^{(i)}\right)-f\left(x^{*}\right) \\
& \leq \frac{\left\|x^{(0)}-x^{*}\right\|_{2}^{2}}{2 t k} \tag{6.7}
\end{align*}
$$

where in the final step, we plug in (??) to get the inequality from (??) that we were trying to prove.

Theorem 6.2 Suppose the function $f: \mathbb{R}^{n} \rightarrow \mathbb{R}$ is convex and differentiable, and that its gradient is Lipschitz continuous with constant $L>0$, i.e. we have that $\|\nabla f(x)-\nabla f(y)\|_{2} \leq L\|x-y\|_{2}$ for any $x, y$. Then if we run gradient descent for $k$ iterations with step size $t_{i}$ chosen using backtracking line search on each iteration $i$, it will yield a solution $f^{(k)}$ which satisfies

$$
\begin{equation*}
f\left(x^{(k)}\right)-f\left(x^{*}\right) \leq \frac{\left\|x^{(0)}-x^{*}\right\|_{2}^{2}}{2 t_{\min } k} \tag{6.8}
\end{equation*}
$$

where $t_{\min }=\min \{1, \beta / L\}$

Convex $f$. From Theorem? ${ }^{6.1}$ ? we know that the convergence rate of gradient descent with convex $f$ is $O(1 / k)$, where $k$ is the number of iterations. This implies that in order to achieve a bound of $f\left(x^{(k)}\right)-f\left(x^{*}\right) \leq$ $\epsilon$, we must run $O(1 / \epsilon)$ iterations of gradient descent. This rate is referred to as "sub-linear convergence."
Strongly convex $\boldsymbol{f}$. In contrast, if we assume that $f$ is strongly convex, we can show that gradient descent converges with rate $O\left(c^{k}\right)$ for $0<c<1$. This means that a bound of $f\left(x^{(k)}\right)-f\left(x^{*}\right) \leq \epsilon$ can be achieved using only $O(\log (1 / \epsilon))$ iterations. This rate is typically called "linear convergence."

### 6.1.4 Pros and cons of gradient descent

The principal advantages and disadvantages of gradient descent are:

- Simple algorithm that is easy to implement and each iteration is cheap; just need to compute a gradient
- Can be very fast for smooth objective functions, i.e. well-conditioned and strongly convex
- However, it's often slow because many interesting problems are not strongly convex
- Cannot handle non-differentiable functions (biggest downside)


## Subgradients

Definition 6.3 A subgradient of a convex function $f: \mathbb{R}^{n} \rightarrow \mathbb{R}$ at some point $x$ is any vector $g \in \mathbb{R}^{n}$ that achieves the same lower bound as the tangent line to $f$ at $x$, i.e. we have

$$
f(y) \geq f(x)+g^{T}(y-x) \quad \forall x, y
$$

The subgradient $g$ always exists for convex functions on the relative interior of their domain. Furthermore, if $f$ is differentiable at $x$, then there is a unique subgradient $g=\nabla f(x)$. Note that subgradients need not exist for nonconvex functions (for example, cubic functions do not have subgradients at their inflection points).

### 6.2.1 Examples of subgradients

absolute value. $f(x)=|x|$. Where $f$ is differentiable, the subgradient is identical to the gradient, sign $(x)$. At the point $x=0$, the subgradient is any point in the range $[-1,1]$ because any line passing through $x=0$ with a slope in this range will lower bound the function.
$\ell_{\mathbf{2}}$ norm. $f(x)=\|x\|_{2}$. For $x \neq 0, f$ is differentiable and the unique subgradient is given by $g=x /\|x\|_{2}$. For $x=0$, the subgradient is any vector whose $\ell_{2}$ norm is at most 1 . This holds because, by definition, in order for $g$ to be a subgradient of $f$ we must have that

$$
f(y)=\|y\|_{2} \geq f(x)+g^{T}(y-x)=g^{T} y \quad \forall y
$$

In order for $\|y\|_{2} \geq g^{T} y$ to hold, $g$ must have $\|g\|_{2} \leq 1$.
$\boldsymbol{\ell}_{1}$ norm. $f(x)=\|x\|_{1}$. Since $\|x\|_{1}=\sum_{i=1}^{n}\left|x_{i}\right|$, we can consider each element $g_{i}$ of the subgradient separately. The result is very analogous to the subgradient of the absolute value function. For $x_{i} \neq 0$, $g_{i}=\operatorname{sign}\left(g_{i}\right)$. For $x_{i}=0, g_{i}$ isanypointin $[-1,1]$.
maximum of two functions. $f(x)=\max \left\{f_{1}(x), f_{2}(x)\right\}$, where $f_{1}$ and $f_{2}$ are convex and differentiable. Here we must consider three cases. First, if $f_{1}(x)>f_{2}(x)$, then $f(x)=f_{1}(x)$ and therefore there is a unique subgradient $g=\nabla f_{1}(x)$. Likewise, if $f_{2}(x)>f_{1}(x)$, then $f(x)=f_{2}(x)$ and $g=\nabla f_{2}(x)$. Finally, if $f_{1}(x)=f_{2}(x)$, then $f$ may not be differentiable at $x$ and the subgradient will be any point on the line segment that joints $\nabla f_{1}(x)$ and $\nabla f_{2}(x)$.

### 6.2.2 Subdifferential

Definition 6.4 The subdifferential of a convex function $f: \mathbb{R}^{n} \rightarrow \mathbb{R}$ at some point $x$ is the set of all subgradients of $f$ at $x$, i.e. we say

$$
\partial f(x)=\left\{g \in \mathbb{R}^{n}: g \text { is a subgradient of } f \text { at } x\right\}
$$

An important property of the subdifferential $\partial f(x)$ is that it is a closed and convex set, which holds even for nonconvex $f$. To verify this, suppose we have two subgradients $g_{1}, g_{2} \in \partial f(x)$. We need to show that $g_{0}=\alpha g_{1}+(1-\alpha) g_{2}$ is also in $\partial f(x)$ for arbitrary $\alpha$. If we write the following inequalities,

$$
\begin{gathered}
\alpha\left(f(y) \geq f(x)+g_{1}^{T}(y-x)\right) \alpha \\
(1-\alpha)\left(f(y) \geq f(x)+g_{2}^{T}(y-x)\right)(1-\alpha)
\end{gathered}
$$

which follow from the definition of subgradient applied to $g_{1}$ and $g_{2}$, we can add them together to yield $f(y) \geq f(x)+\alpha g_{1}^{T}(y-x)+(1-\alpha) g_{2}^{T}(y-x)=g_{0}^{T}(y-x)$.

### 7.2.1 Subgradient method

For convex $f$, not necessarily differentiable, subgradient method finds the lowest value of the criterion by:

$$
x^{(k)}=x^{(k-1)}-t_{k} g^{(k-1)}, \quad, k=1,2,3, \cdots
$$

where $g^{(k-1)}$ is any subgradient of $f$ at $x^{(k-1)}$. Note that it is not a decent method, that the next iterative doesn't always find the lower criterion. So we need to keep the best lowest criterion value at every iteration, i.e., $f\left(x_{\text {best }}^{(k)}\right)=\min _{i} f\left(x^{(i)}\right)$.

### 7.2.2 Choosing the step size

i) Fixed step size: $t_{k}=t \forall k$.

However, for subgradient method, we do not typically chose fixed step size.
ii) Diminishing step size (Standard): choose $t_{k}$ that is square summable but not summable.

$$
\sum_{k=1}^{\infty} t_{k}^{2}<\infty, \quad \sum_{k=1}^{\infty}=\infty
$$

Note that step sizes are all pre-defined, not adaptively computed during the optimization iteration.

### 7.2.3 Convergence analysis

i) Fixed step size: Suboptimal Convergence.

For convex, not differentiable function $f$, if the function itself is Lipschitz with constant $G$ such as,

$$
|f(x)-f(y)| \leq G\|x-y\|_{2} \quad \forall x, y
$$

subgradient method using fixed step size $t$ would give a point that is suboptimal such as,

$$
\lim _{k \rightarrow \infty} f\left(x_{\text {best }}^{(k)}\right) \leq f\left(x^{*}\right)+G^{2} \frac{t}{2}
$$

In other words, the smaller the step size, the smaller the difference would be between the optimal and suboptimal convergence.
ii) Diminishing step size that is square summable: Optimal Convergence.

$$
\lim _{k \rightarrow \infty} f\left(x_{\text {best }}^{(k)}\right)=f\left(x^{*}\right) .
$$

Note that subgradient method is applicable to functions that may not look like Lipschitz, since the over the bounded set the function can be Lipschitz.

## Projection Method

Projected subgradient method can be used to minimize a convex function over a convex set C:

$$
\min _{x \in C} f(x)
$$

It is same as usual subgradient update except we project the solution back on to C every time so that at every iteration we move in the direction of the subgradient but still lies in the set C.

$$
x^{(k)}=P_{C}\left(x^{(k-1)}-t_{k} g^{(k-1)}\right), \quad k=1,2,3, \cdots
$$

Alternative method:

$$
\min _{x \in C} f(x)=\min _{x \in \mathbb{R}^{n}} f(x)+I_{C}(x)
$$

## Examples for projection onto solution set C:

i) $C=\left\{y: y_{i} \geq \forall i\right\} \Rightarrow\left[P_{C}(x)\right]_{i}=\max \left\{x_{i}, 0\right\}$.

### 7.2.7 Basic Pursuit Problem

We can use projected subgradient method to solve the basic pursuit problem:

$$
\min _{\beta \in \mathbb{R}^{p}}\|\beta\|_{1} \text { s.t. } X \beta=y \text {. }
$$

In this case, the solution set is $C=\{\beta: X \beta=y\}$.
The projection on to solution set C is $P_{C}(\beta)=\beta+X^{T}\left(X X^{T}\right)^{-1}(y-X \beta)$ as shown in example 2 above. Projected subgradient method performs step

$$
\begin{aligned}
\beta^{(k)} & =P_{C}\left(\beta^{(k-1)}-t_{k} g^{(k-1)}\right) \\
& \left.=\beta^{(k-1)}-t_{k} g^{(k-1)}+X^{( } X X^{T}\right)^{-1}\left(y-X \beta^{(k-1)}+X t_{k} g^{(k-1)}\right) \\
& =\beta^{(k-1)}-\left(I-X^{T}\left(X X^{T}\right)^{-1} X\right) t_{k} g^{(k-1)}
\end{aligned}
$$

Where, $g^{(k-1)} \in \partial\left\|\beta^{(k-1)}\right\|_{1}$.

