## Problem Set 6

## CS 331H

## Due Thursday, April 1

1. Consider minimum spanning tree algorithms for the following graph:



- (a) In what order would Prim's algorithm, starting at s, add edges to the minimum spanning tree? Give the sequence of edge weights, in order.
- (b) In what order would Kruskal's algorithm add edges to the minimum spanning tree? Give the sequence of edge weights, in order.
- (c) In what order would Boruvka's algorithm add edges to the minimum spanning tree? Give the set of edge weights added in the first round, the second round, etc.
- 2. You are building out internet for a collection of rural houses. For each house, you need to either purchase satellite internet at that house, or connect it via a series of fiber links to a house that has purchased satellite internet.

There are *n* houses, and buying satellite internet costs *P* dollars at any house. There are *m* pairs of houses that can be directly connected by fiber; this is given as a list of triples  $(u_i, v_i, c_i)$ , denoting that houses  $u_i$  and  $v_i$  can be connected at a cost of  $c_i$  dollars.

Give an  $O(m \log n)$  time algorithm to determine the minimum cost of hooking everyone up to internet.

- 3. Consider a weighted, directed graph where all distances lie in [1, 2). We would like to find an O(E) time algorithm for single-source shortest paths on this graph.
  - (a) Consider a variant of Dijkstra's algorithm that does not always visit the unvisited node of smallest c(u), but instead arbitrarily picks one of the unvisited nodes of smallest  $\lfloor c(u) \rfloor$ . Show that such an algorithm still yields the correct answer.
  - (b) Now give a data structure that allows this Dijkstra variant to run in O(E) time. Hint: at any point during the execution, the set of [c(u)] for unvisited nodes u can only have a small number of options.
  - (c) Extend your result to O(EC) time and O(E+C) space for distances in [1, C) for any  $C \ge 1$ .