

# Upper and Lower Loop Bound Estimation by Symbolic Execution and Loop Acceleration

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## Loop Bound Analysis

What is the maximal number of loop iterations?

### Usage

- worst case execution time
- complexity analysis
- resource consumption
- schedulers
- ...

## Reachability Bounds

What is the maximal number of executions of a specific program part?

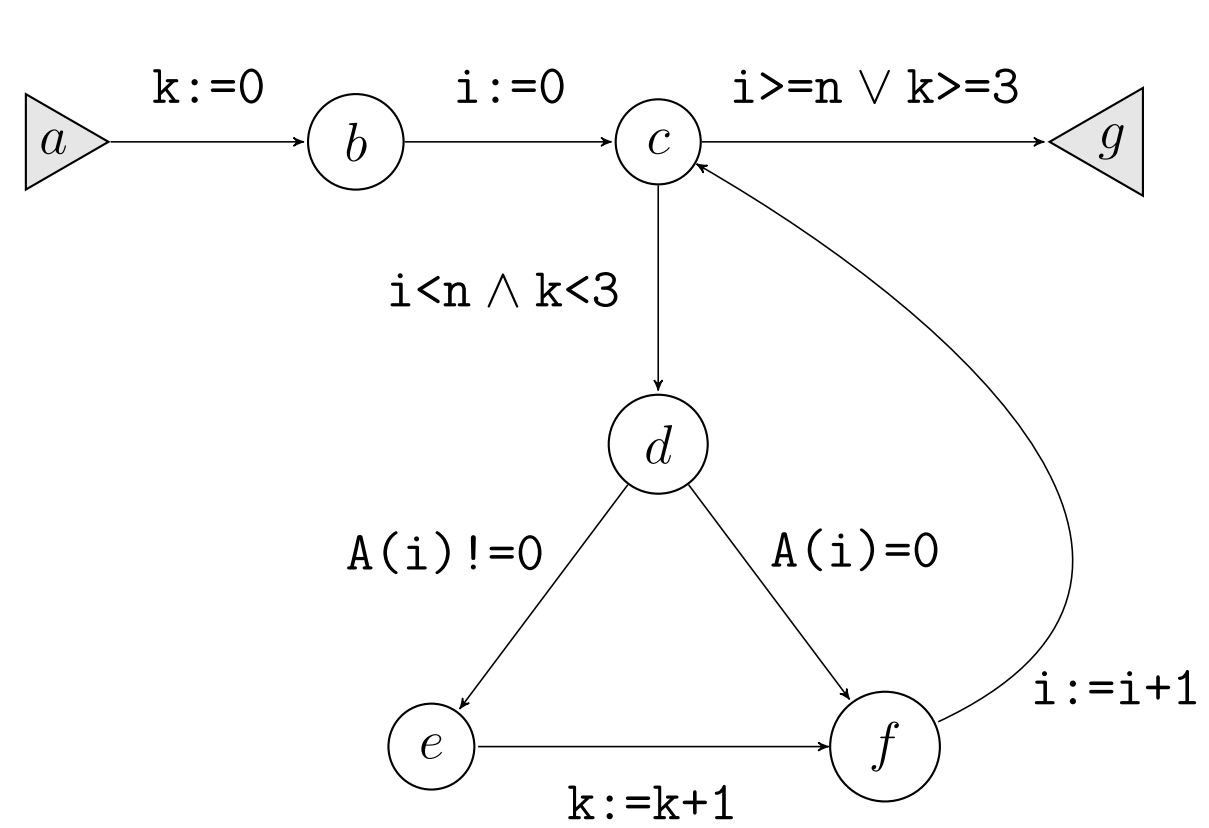
## Lower Bounds

What is the minimal number of executions of a specific program part?

### Usage

- improving precision of upper bounds
- best case computational complexity
- estimating tightness of upper bounds
- invariant generation

## Example



## Loop Acceleration

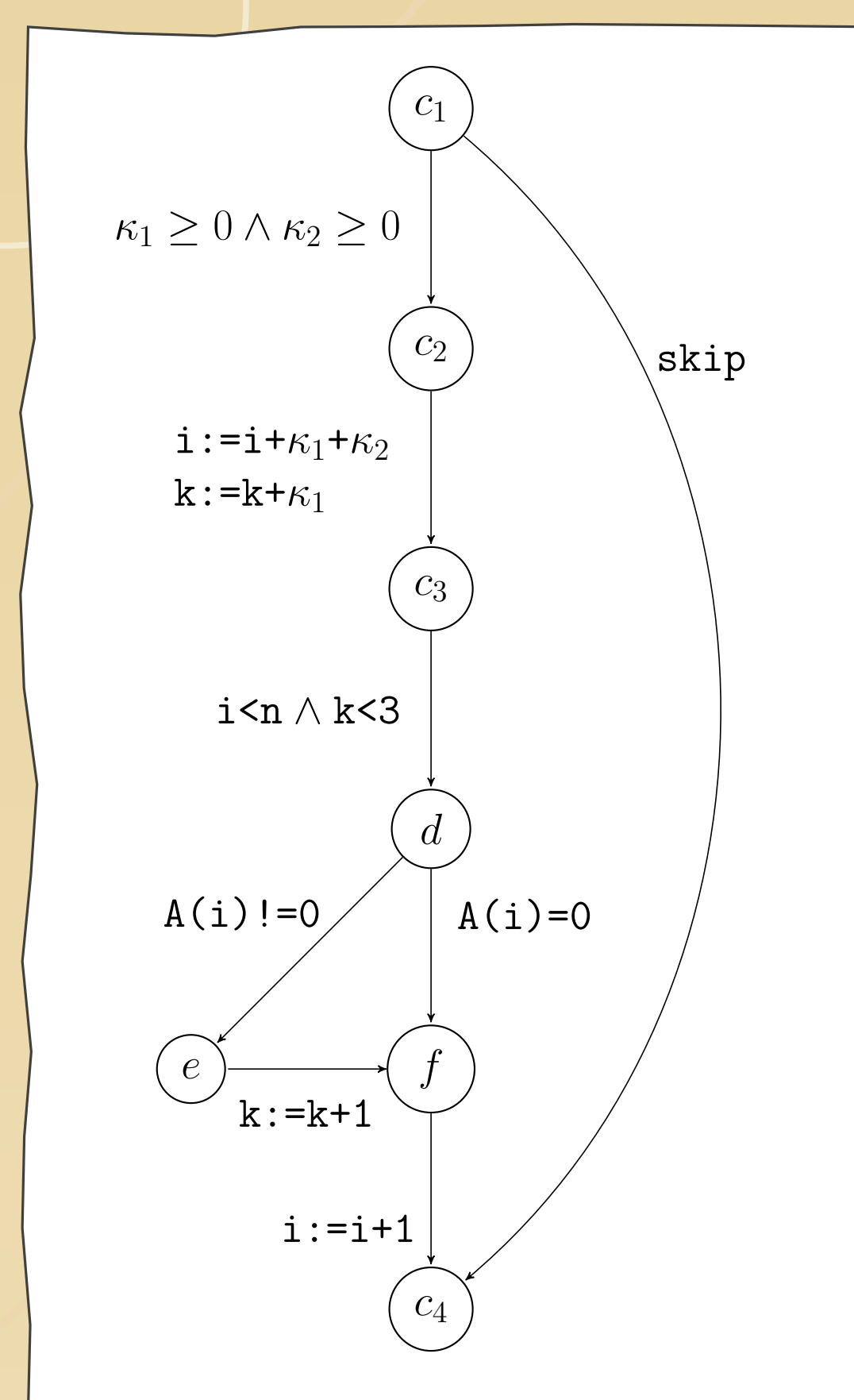
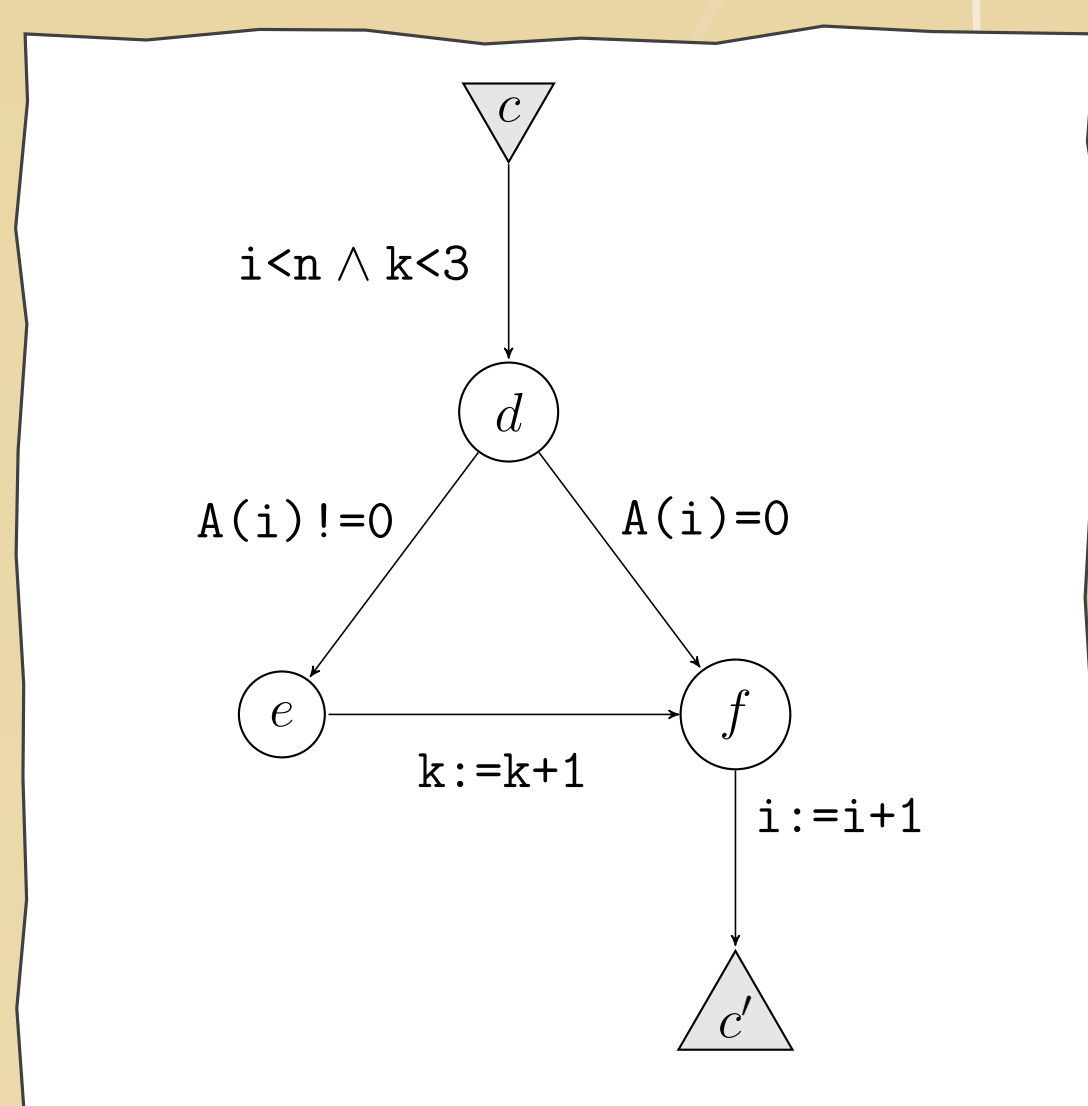
- representing several loop iterations by one transition
- solving path explosion for bounded model checking and symbolic execution

### Usage

- Loop Underapproximation
  - an accelerated path added, others remain
  - bug finding (reducing depth of a bug)
- Loop Overapproximation
  - all paths replaced
  - resulting flowgraph acyclic
  - proving correctness

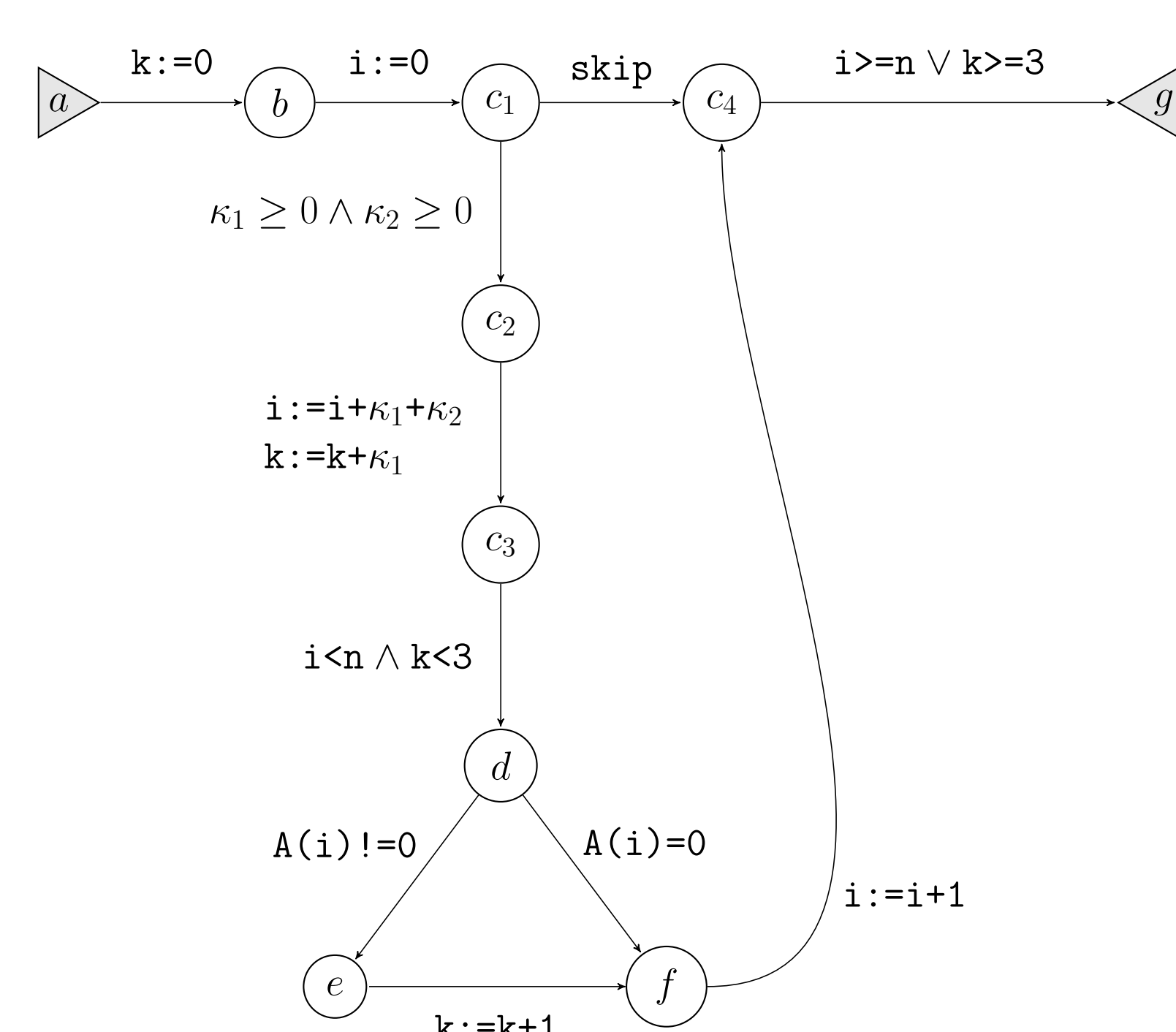
## Our Acceleration Approach

1. make a flowgraph induced by the loop, which represents one loop iteration
2. symbolically execute the induced flowgraph
3. introduce a counter  $\kappa_i$  for each loop path  $i$
4. summarize the effect of  $\kappa_1 + \kappa_2 + \dots$  iterations
5. make 2 branches
  - 0 iterations
  - summary of  $\kappa_1 + \kappa_2 + \dots$  iterations and a copy of the last iteration



$$\begin{aligned} \theta_1(i) &= \underline{i} + 1 \\ \theta_1(k) &= \underline{k} + 1 \\ \theta_2(i) &= \underline{i} + 1 \\ \theta_2(k) &= \underline{k} \\ \theta^{\vec{\kappa}}(i) &= \underline{i} + \kappa_1 + \kappa_2 \\ \theta^{\vec{\kappa}}(k) &= \underline{k} + \kappa_1 \end{aligned}$$

## Replacing the Loop with its Acyclic Overapproximation



## Symbolic Execution

Instead of normal program inputs one supplies symbols representing arbitrary values. The execution splits into two at branching statements.

### Symbolic Memory ( $\theta$ )

- changes after each assignment
- values of variables as symbolic expressions

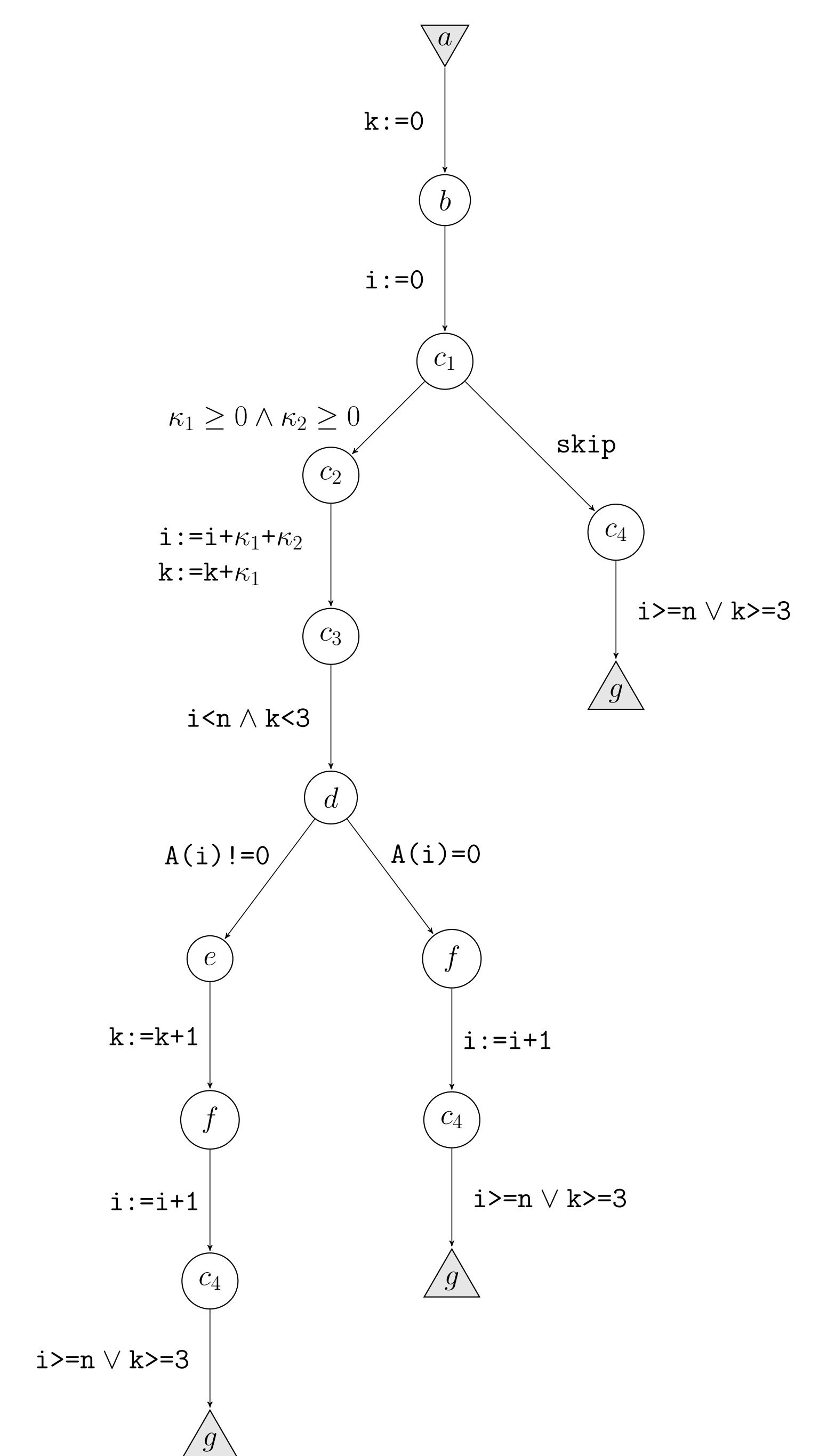
### Path Condition ( $\varphi$ )

- changes after each branching condition
- a necessary and sufficient condition for an execution to follow the particular path

### Usage

- summarizing effect of an acyclic piece of code
- bug finding

## Symbolic Execution of the Acyclic Program



Lower reachability bound:

$$\begin{aligned} (b, c) &: 1 \\ (c, d) &: \max\{\min\{n, 3\}, 0\} \\ (d, e) &: 0 \\ (d, f) &: \max\{n - 3, 0\} \end{aligned}$$

Upper reachability bound:

$$\begin{aligned} (b, c) &: 1 \\ (c, d) &: \max\{n, 0\} \\ (d, e) &: \max\{\min\{n, 3\}, 0\} \\ (d, f) &: \max\{n, 0\} \end{aligned}$$

## References

- [1] J. Strejček and M. Trtík. "Abstracting Path Conditions". In: ISSA. 2012, pp. 155–165.
- [2] Pavel Čadek, Jan Strejček, and Marek Trtík. "Tighter Loop Bound Analysis". In: ATVA. 2016, pp. 512–527.

$$\begin{aligned} \varphi_1 &\equiv 0 \geq \underline{n} \vee 0 \geq 3 \\ \varphi_2 &\equiv \kappa_1 \geq 0 \wedge \kappa_2 \geq 0 \wedge \kappa_1 + \kappa_2 < \underline{n} \wedge \kappa_1 < 3 \wedge \underline{A}(\kappa_1 + \kappa_2) \neq 0 \wedge (\kappa_1 + \kappa_2 + 1 \geq \underline{n} \vee \kappa_1 + 1 \geq 3) \\ &\quad \kappa_1 + \kappa_2 < \underline{n} \\ &\quad \kappa_1 + \kappa_2 + 1 \geq \max\{\underline{n}, 3\} \\ &\Rightarrow \kappa_1 < \min\{\underline{n}, 3\} \\ &\quad \kappa_1 \geq 0 \\ &\quad \kappa_2 \geq \max\{\underline{n} - 3, 0\} \\ &\quad \kappa_2 < \underline{n} \\ \varphi_3 &\equiv \kappa_1 \geq 0 \wedge \kappa_2 \geq 0 \wedge \kappa_1 + \kappa_2 < \underline{n} \wedge \kappa_1 < 3 \wedge \underline{A}(\kappa_1 + \kappa_2) = 0 \wedge (\kappa_1 + \kappa_2 + 1 \geq \underline{n} \vee \kappa_1 \geq 3) \\ &\Rightarrow \dots \end{aligned}$$