

Asymptotic Dominance Theory

- **Definition 1:** Given the functions $f: \mathbb{N} \rightarrow \mathbb{R}$ and $g: \mathbb{N} \rightarrow \mathbb{R}$, f is *asymptotically dominated* by g if there exist non-negative constants M and N such that for all $n \geq N$, $|f(n)| \leq M|g(n)|$. This is denoted by $f = O(g)$.
- **Definition 2:** Given the functions $f: \mathbb{N} \rightarrow \mathbb{R}$ and $g: \mathbb{N} \rightarrow \mathbb{R}$, $f = o(g)$ if for every positive ε , there exists a non-negative constant N such that for all $n \geq N$, $|f(n)| \leq \varepsilon|g(n)|$.

Theorem 1: If $f = O(g)$, then for any constant s , $sf = O(g)$.

Proof: By definition, there exist non-negative constants M and N such that for all $n \geq N$, $|f(n)| \leq M|g(n)|$. Thus for all $n \geq N$, $|sf(n)| \leq |s|M|g(n)|$. Therefore, $sf = O(g)$. \square

Theorem 2: If $f_1 = O(g_1)$ and $f_2 = O(g_2)$, then $f_1 + f_2 = O(|g_1| + |g_2|)$.

Proof: By definition, there exist non-negative constants M_1 and N_1 such that for all $n \geq N_1$, $|f_1(n)| \leq M_1|g_1(n)|$ and there exist non-negative constants M_2 and N_2 such that for all $n \geq N_2$, $|f_2(n)| \leq M_2|g_2(n)|$. For $n \geq \max\{N_1, N_2\}$ both inequalities hold so $|f_1(n) + f_2(n)| \leq |f_1(n)| + |f_2(n)| \leq M_1|g_1(n)| + M_2|g_2(n)| \leq \max\{M_1, M_2\}(|g_1(n)| + |g_2(n)|)$. Therefore, $f_1 + f_2 = O(|g_1| + |g_2|)$. \square

Corollary 2.1: If for $i = 1, 2, \dots, k$, $f_i = O(g_i)$, then $\sum_{i=1}^k f_i = O(\sum_{i=1}^k |g_i|)$.

Theorem 3: If $f_1 = O(g_1)$ and $f_2 = O(g_2)$, then $f_1 + f_2 = O(\max\{|g_1|, |g_2|\})$.

Proof: By definition, there exist non-negative constants M_1 and N_1 such that for all $n \geq N_1$, $|f_1(n)| \leq M_1|g_1(n)|$ and there exist non-negative constants M_2 and N_2 such that for all $n \geq N_2$, $|f_2(n)| \leq M_2|g_2(n)|$. For $n \geq \max\{N_1, N_2\}$ both inequalities hold so $|f_1(n) + f_2(n)| \leq |f_1(n)| + |f_2(n)| \leq M_1|g_1(n)| + M_2|g_2(n)| \leq (M_1 + M_2) \max\{|g_1(n)|, |g_2(n)|\}$. Therefore, $f_1 + f_2 = O(\max\{|g_1|, |g_2|\})$. \square

Corollary 3.1: If for $i = 1, 2, \dots, k$, $f_i = O(g_i)$, then $\sum_{i=1}^k f_i = O(\max_{i=1, \dots, k} |g_i|)$.

Corollary 3.2: If for $i = 1, 2, \dots, k$, $f_i = O(g)$, then $\sum_{i=1}^k f_i = O(g)$.

Theorem 4: If $f_1 = O(g_1)$ and $f_2 = O(g_2)$, then $f_1 \cdot f_2 = O(g_1 \cdot g_2)$.

Proof: By definition, there exist non-negative constants M_1 and N_1 such that for all $n \geq N_1$, $|f_1(n)| \leq M_1 |g_1(n)|$ and there exist non-negative constants M_2 and N_2 such that for all $n \geq N_2$, $|f_2(n)| \leq M_2 |g_2(n)|$. For $n \geq \max\{N_1, N_2\}$ both inequalities hold so $|f_1(n) \cdot f_2(n)| = |f_1(n)| \cdot |f_2(n)| \leq M_1 |g_1(n)| \cdot M_2 |g_2(n)| \leq (M_1 \cdot M_2) (|g_1(n)| \cdot |g_2(n)|)$. Therefore, $f_1 \cdot f_2 = O(g_1 \cdot g_2)$. \square

Corollary 4.1: If for $i = 1, 2, \dots, k$, $f_i = O(g_i)$, then $\prod_{i=1}^k f_i = O(\prod_{i=1}^k g_i)$.

Theorem 5: If $f_1 = O(g_1)$, $g_2 = O(f_2)$, and g_2 has no zeros. then $f_1 / f_2 = O(g_1 / g_2)$.

Proof: By definition, there exist non-negative constants M_1 and N_1 such that for all $n \geq N_1$, $|f_1(n)| \leq M_1 |g_1(n)|$ and there exist non-negative constants M_2 and N_2 such that for all $n \geq N_2$, $|g_2(n)| \leq M_2 |f_2(n)|$. Notice that since g_2 has no zeros, then neither does f_2 . Inverting this inequality, we obtain that for all $n \geq N_2$, $|1 / f_2(n)| \leq M_2 |1 / g_2(n)|$. For $n \geq \max\{N_1, N_2\}$ both inequalities hold so $|f_1(n) / f_2(n)| = |f_1(n)| \cdot |1 / f_2(n)| \leq M_1 |g_1(n)| \cdot M_2 |1 / g_2(n)| \leq (M_1 \cdot M_2) (|g_1(n)| / |g_2(n)|)$. Therefore, $f_1 / f_2 = O(g_1 / g_2)$. \square

Theorem 6: If $a \leq b$, then $n^a = O(n^b)$

Proof: For $n \geq 0$, $n^{-(b-a)} \leq n^0 = 1$, and $|n^a| = |n^{-(b-a)} n^b| = |n^{-(b-a)}| |n^b| \leq 1 \cdot |n^b|$. Therefore, $n^a = O(n^b)$. \square

Theorem 7: If $a < b$, then $n^a = o(n^b)$

Proof: Given any $\varepsilon > 0$, let $N = (1 / \varepsilon)^{1/(b-a)}$. Notice then for $n \geq N = (1 / \varepsilon)^{1/(b-a)}$, $n^{b-a} \geq 1 / \varepsilon$, and $n^{-(b-a)} \leq \varepsilon$. So $|n^a| = |n^{-(b-a)} n^b| = |n^{-(b-a)}| |n^b| \leq \varepsilon |n^b|$. Therefore, $n^a = o(n^b)$. \square

Example 1: If $f_1 = O(g_1)$ and $f_2 = O(g_2)$, then f_1 / f_2 may not be $O(g_1 / g_2)$.

Proof: Let $f_1(n) = f_2(n) = 1$, for all $n \geq 0$. Then $f_1 = O(1)$ and $f_2 = O(n)$ but $f_1 / f_2 = 1 \neq O(1/n)$. To see this, consider any $N \geq 0$ and $M \geq 0$. Choose any $n > \max\{N, M\}$. Notice that then $|1| = 1 > M / |n|$, so $1 \neq O(1/n)$. \square

Example 2: If $a < b$, then $n^b \neq O(n^a)$.

Proof: Suppose $n^b = O(n^a)$, then there exist $N \geq 0$ and $M \geq 0$ so that for all $n \geq N$, $|n^b| \leq M |n^a|$. Choose any $n > \max\{N, M^{1/(b-a)}\}$. Notice that then $n^{b-a} > M$, so $|n^b| = n^b > M n^a = M |n^a|$, and $n^b \neq O(n^a)$. \square

Example 3: If f is any polynomial of degree k then $f = O(n^k)$.

Proof-1: Without loss of generality, assume $f(n) = \sum_{i=0}^k a_i n^i$. For all $n \geq 0$ and $0 \leq i \leq k$,

$|n^i| \leq |n^k|$ and $|a_i n^i| \leq |a_i| |n^k|$. So for $N = 0$ and $M = \sum_{i=0}^k |a_i|$, we have

$n \geq N$ implies

$$|f(n)| = \left| \sum_{i=0}^k a_i n^i \right| \leq \sum_{i=0}^k |a_i n^i| \leq \sum_{i=0}^k |a_i| |n^k| \leq \left(\sum_{i=0}^k |a_i| \right) \cdot |n^k| = M |n^k| \quad \square$$

Proof- 2: Without loss of generality, assume $f(n) = \sum_{i=0}^k a_i n^i$. By Theorem 6, $n^i = O(n^k)$,

for $0 \leq i \leq k$. By Theorem 1, $a_i n^i = O(n^k)$ for $0 \leq i \leq k$. Finally, from Corollary 3.2,

$$f(n) = \sum_{i=0}^k a_i n^i = O(n^k). \quad \square$$