Systems I

Machine-Level Programming II: Introduction

Topics

- Addresses and Pointers
- Memory address modes
- Arithmetic operations
- RISC vs. CISC

Addresses and Pointers in C

C programming model is close to machine language

- Machine language manipulates memory addresses
  - Address computation
  - Store addresses in registers or memory
- C employs pointers, which are just addresses of primitive data elements or data structures

Examples of operators * and &

- int a, b; /* declaration of a and b as an integers */
- int *a_ptr; /* a is a pointer to an integer (address of memory)
- a_ptr = a; /* illegal as the types don’t match */
- a_ptr = &a; /* a_ptr holds address of “a” */
- b = *a_ptr; /* dereference (lookup) value at address a_ptr and assign value to b */

Using Simple Addressing Modes

```
void swap(int *xp, int *yp)
{
  int t0 = *xp;
  int t1 = *yp;
  *xp = t1;
  *yp = t0;
}
```

```
swap:
    pushl %ebp
    movl %esp,%ebp
    pushl %ebx
```

Set Up

```
    movl 12(%ebp),%ecx
    movl 8(%ebp),%edx
    movl (%ecx),%eax
    movl (%edx),%ebx
    movl %eax,(%edx)
    movl %ebx,(%ecx)
```

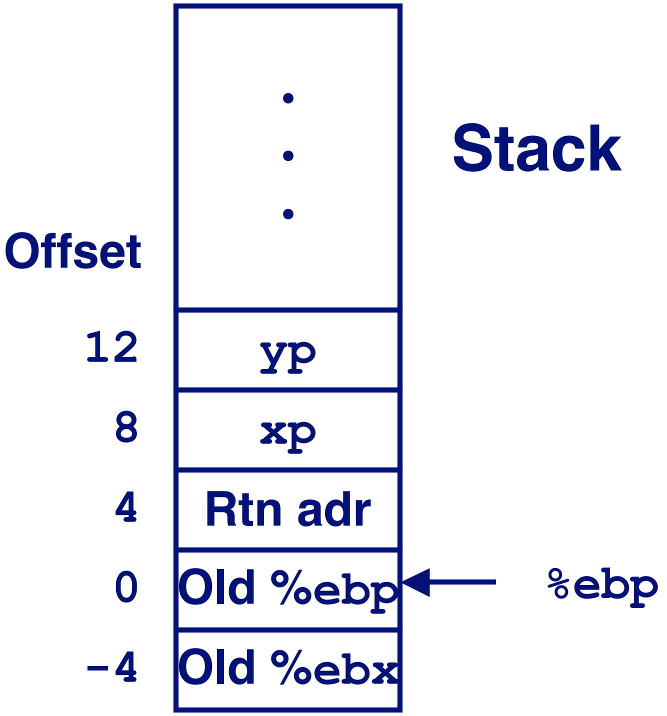
Body

```
    movl -4(%ebp),%ebx
    movl %ebp,%esp
    popl %ebp
    ret
```

Finish

Understanding Swap

```
void swap(int *xp, int *yp)
{
  int t0 = *xp;
  int t1 = *yp;
  *xp = t1;
  *yp = t0;
}
```

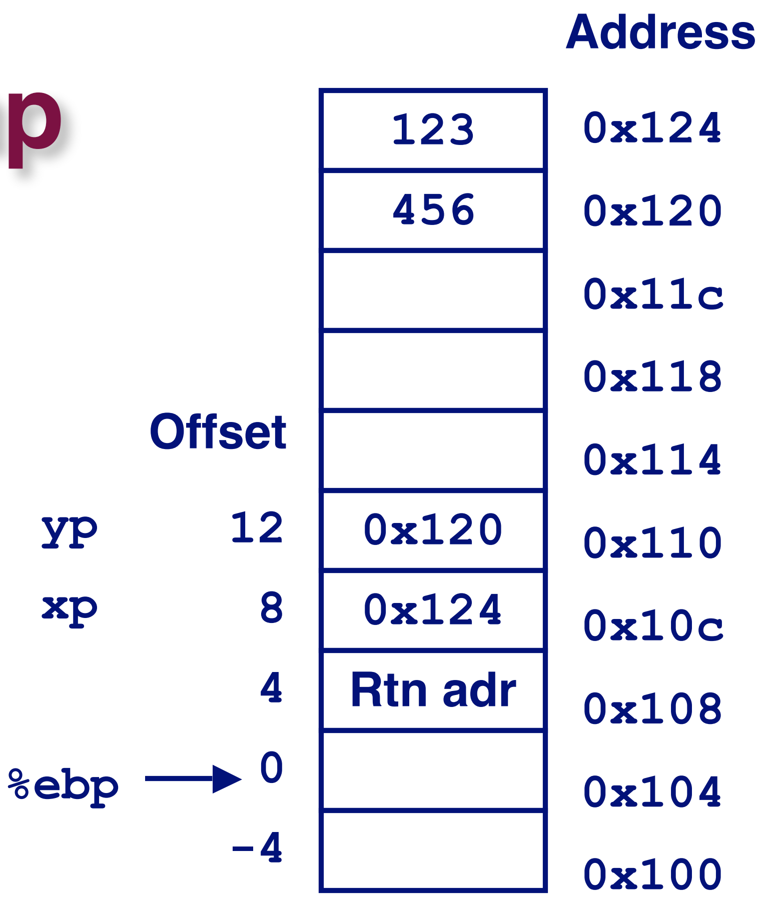


Register	Variable
%ecx	yp
%edx	xp
%eax	t1
%ebx	t0

```
movl 12(%ebp),%ecx   # ecx = yp
movl 8(%ebp),%edx    # edx = xp
movl (%ecx),%eax     # eax = *yp (t1)
movl (%edx),%ebx     # ebx = *xp (t0)
movl %eax,(%edx)     # *xp = eax
movl %ebx,(%ecx)     # *yp = ebx
```

Understanding Swap

%eax	
%edx	
%ecx	
%ebx	
%esi	
%edi	
%esp	
%ebp	0x104



	Offset		Address
		123	0x124
		456	0x120
			0x11c
			0x118
			0x114
yp	12	0x120	0x110
xp	8	0x124	0x10c
	4	Rtn adr	0x108
%ebp →	0		0x104
	-4		0x100

```
movl 12(%ebp),%ecx   # ecx = yp
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Understanding Swap

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Understanding Swap

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Understanding Swap

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Understanding Swap

%eax	456
%edx	0x124
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%ebx	123
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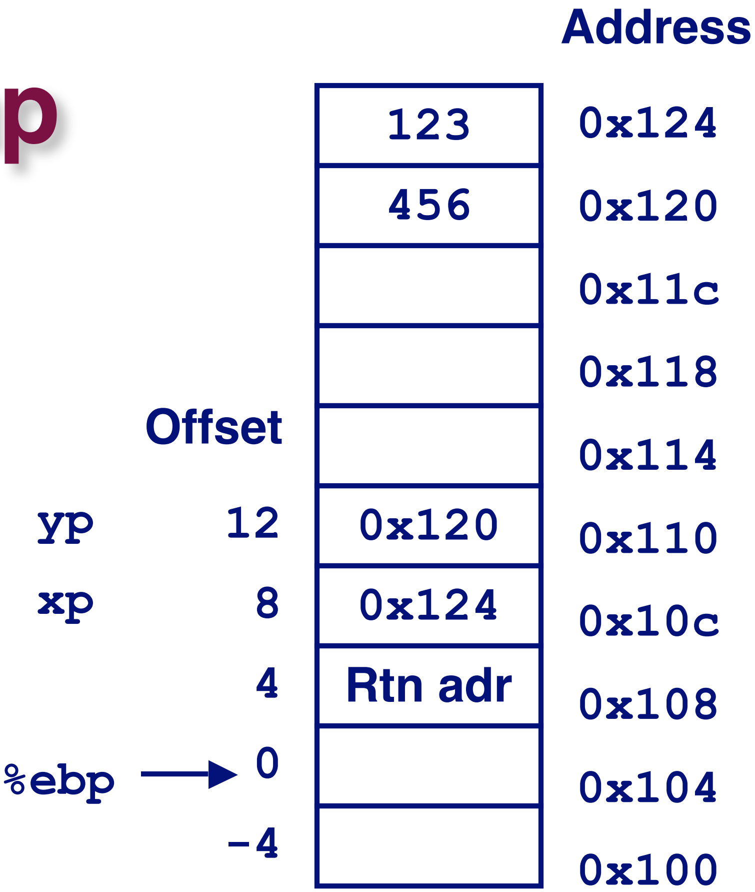

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Understanding Swap

%eax	456
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%edi	
%esp	
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movl (%ecx),%eax     # eax = *yp (t1)
movl (%edx),%ebx     # ebx = *xp (t0)
movl %eax,(%edx)     # *xp = eax
movl %ebx,(%ecx)     # *yp = ebx
```

Understanding Swap

%eax	456
%edx	0x124
%ecx	0x120
%ebx	123
%esi	
%edi	
%esp	
%ebp	0x104

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```
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movl 8(%ebp),%edx    # edx = xp
movl (%ecx),%eax     # eax = *yp (t1)
movl (%edx),%ebx     # ebx = *xp (t0)
movl %eax,(%edx)     # *xp = eax
movl %ebx,(%ecx)     # *yp = ebx
```

Indexed Addressing Modes

Most General Form

D(Rb,Ri,S) Mem[Reg[Rb]+S*Reg[Ri]+ D]

- D: Constant “displacement” 1, 2, or 4 bytes
- Rb: Base register: Any of 8 integer registers
- Ri: Index register: Any, except for `%esp`
    - Unlikely you’d use `%ebp`, either
- S: Scale: 1, 2, 4, or 8

Special Cases

(Rb,Ri)	Mem[Reg[Rb]+Reg[Ri]]
D(Rb,Ri)	Mem[Reg[Rb]+Reg[Ri]+D]
(Rb,Ri,S)	Mem[Reg[Rb]+S*Reg[Ri]]

Address Computation Examples

`%edx`	`0xf000`
`%ecx`	`0x100`

Expression	Computation	Address
`0x8(%edx)`	`0xf000 + 0x8`	`0xf008`
`(%edx,%ecx)`	`0xf000 + 0x100`	`0xf100`
`(%edx,%ecx,4)`	`0xf000 + 4*0x100`	`0xf400`
`0x80(,%edx,2)`	`2*0xf000 + 0x80`	`0x1e080`

Address Computation Instruction

`leal` *Src*,*Dest*

- *Src* is address mode expression
- Set *Dest* to address denoted by expression

Uses

- Computing address without doing memory reference
  - E.g., translation of `p = &x[i];`
- Computing arithmetic expressions of the form x + k*y
  - k = 1, 2, 4, or 8.

Some Arithmetic Operations

Format	Computation	
Two Operand Instructions		
`addl` *Src,Dest*	*Dest* = *Dest* + *Src*	
`subl` *Src,Dest*	*Dest* = *Dest* - *Src*	
`imull` *Src,Dest*	*Dest* = *Dest* * *Src*	
`sall` *Src,Dest*	*Dest* = *Dest* << *Src*	Also called `shll`
`sarl` *Src,Dest*	*Dest* = *Dest* >> *Src*	Arithmetic
`shrl` *Src,Dest*	*Dest* = *Dest* >> *Src*	Logical
`xorl` *Src,Dest*	*Dest* = *Dest* ^ *Src*	
`andl` *Src,Dest*	*Dest* = *Dest* & *Src*	
`orl` *Src,Dest*	*Dest* = *Dest* | *Src*	

Some Arithmetic Operations

Format	Computation
One Operand Instructions	
`incl` ***Dest***	***Dest*** `= ` ***Dest*** ` + 1`
`decl` ***Dest***	***Dest*** `= ` ***Dest*** ` - 1`
`negl` ***Dest***	***Dest*** `= -` ***Dest***
`notl` ***Dest***	***Dest*** `= ~` ***Dest***

Using `leal` for Arithmetic Expressions

```
int arith
  (int x, int y, int z)
{
  int t1 = x+y;
  int t2 = z+t1;
  int t3 = x+4;
  int t4 = y * 48;
  int t5 = t3 + t4;
  int rval = t2 * t5;
  return rval;
}
```

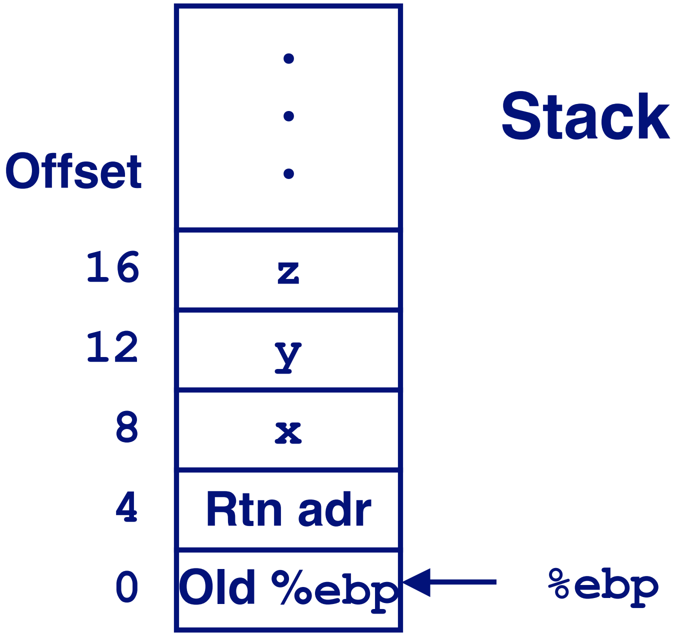
```
arith:
   pushl %ebp                           Set
   movl %esp,%ebp                       Up

   movl 8(%ebp),%eax
   movl 12(%ebp),%edx
   leal (%edx,%eax),%ecx
   leal (%edx,%edx,2),%edx
   sall $4,%edx                         Body
   addl 16(%ebp),%ecx
   leal 4(%edx,%eax),%eax
   imull %ecx,%eax

   movl %ebp,%esp
   popl %ebp                            Finish
   ret
```

Understanding `arith`

```
int arith
  (int x, int y, int z)
{
  int t1 = x+y;
  int t2 = z+t1;
  int t3 = x+4;
  int t4 = y * 48;
  int t5 = t3 + t4;
  int rval = t2 * t5;
  return rval;
}
```

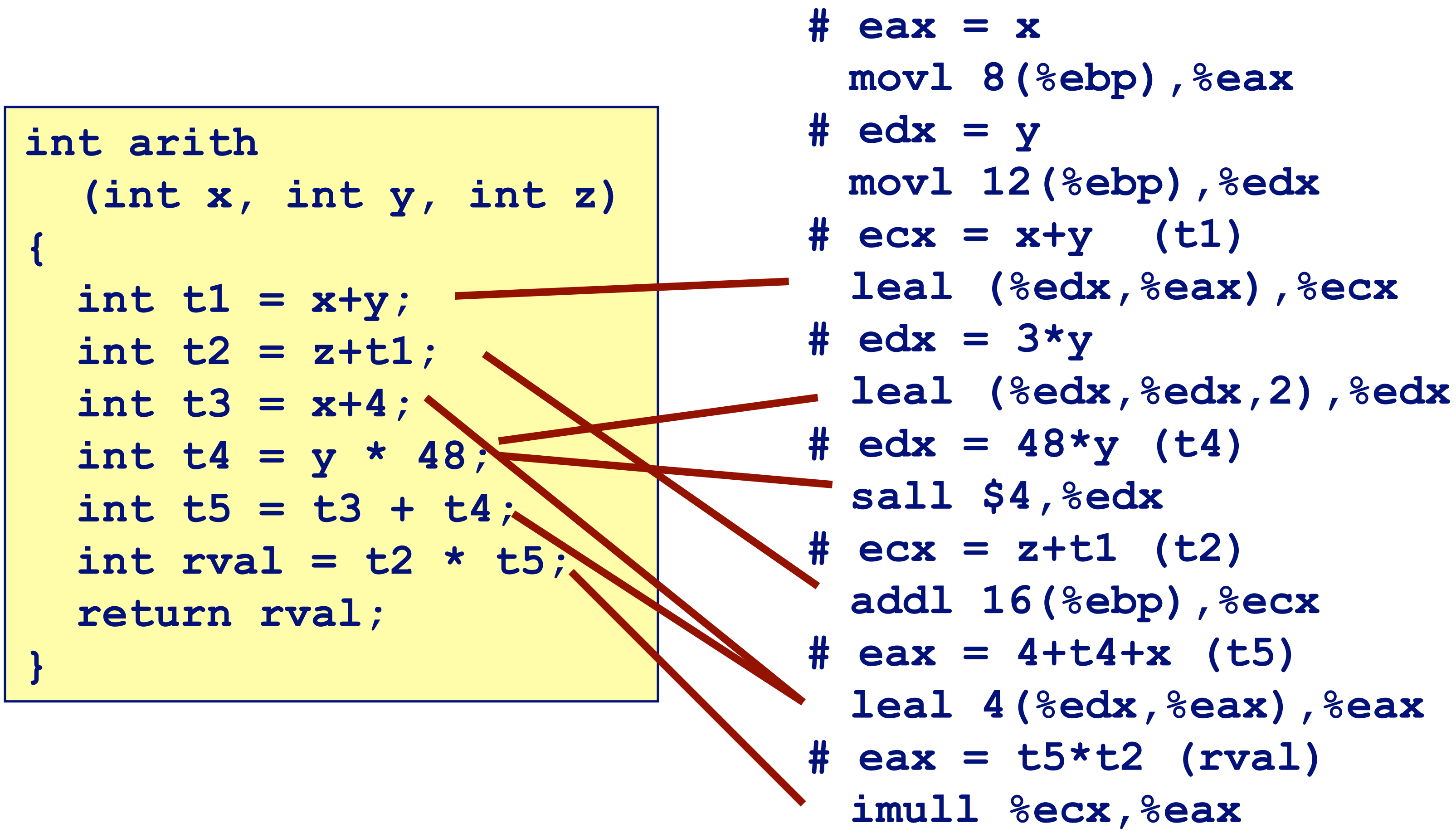


Stack

Offset	
	• • •
16	z
12	y
8	x
4	Rtn adr
0	Old %ebp ← %ebp

```
movl 8(%ebp),%eax        # eax = x
movl 12(%ebp),%edx       # edx = y
leal (%edx,%eax),%ecx    # ecx = x+y   (t1)
leal (%edx,%edx,2),%edx  # edx = 3*y
sall $4,%edx             # edx = 48*y (t4)
addl 16(%ebp),%ecx       # ecx = z+t1 (t2)
leal 4(%edx,%eax),%eax   # eax = 4+t4+x (t5)
imull %ecx,%eax          # eax = t5*t2 (rval)
```

Understanding `arith`



```
int arith
  (int x, int y, int z)
{
  int t1 = x+y;
  int t2 = z+t1;
  int t3 = x+4;
  int t4 = y * 48;
  int t5 = t3 + t4;
  int rval = t2 * t5;
  return rval;
}
```

```
# eax = x
  movl 8(%ebp),%eax
# edx = y
  movl 12(%ebp),%edx
# ecx = x+y   (t1)
  leal (%edx,%eax),%ecx
# edx = 3*y
  leal (%edx,%edx,2),%edx
# edx = 48*y (t4)
  sall $4,%edx
# ecx = z+t1 (t2)
  addl 16(%ebp),%ecx
# eax = 4+t4+x (t5)
  leal 4(%edx,%eax),%eax
# eax = t5*t2 (rval)
  imull %ecx,%eax
```

Another Example

```
int logical(int x, int y)
{
  int t1 = x^y;
  int t2 = t1 >> 17;
  int mask = (1<<13) - 7;
  int rval = t2 & mask;
  return rval;
}
```

$2^{13} = 8192, 2^{13} – 7 = 8185$

```
logical:
    pushl %ebp
    movl %esp,%ebp
```

Set Up

```
    movl 8(%ebp),%eax
    xorl 12(%ebp),%eax
    sarl $17,%eax
    andl $8185,%eax
```

Body

```
    movl %ebp,%esp
    popl %ebp
    ret
```

Finish

```
movl 8(%ebp),%eax       eax = x
xorl 12(%ebp),%eax      eax = x^y      (t1)
sarl $17,%eax           eax = t1>>17   (t2)
andl $8185,%eax         eax = t2 & 8185
```

ISA Properties

CISC - Complex Instruction Set Computer (e.g. x86)

- Instruction can reference different operand types
  - Immediate, register, memory
- Arithmetic operations can read/write memory
- Memory reference can involve complex computation
  - Rb + S*Ri + D
  - Useful for arithmetic expressions, too
- Instructions can have varying lengths
  - x86 IA32 instructions can range from 1 to 15 bytes

"RISC" - Reduced Instruction Set Computer

- e.g. ARM, PowerPC, Sparc
- Memory operations separate from arithmetic (load/store)
- Fixed length instructions (often 4 bytes each)
- Fewer complex computational instructions (e.g. stringe compare)

Summary

Today

- C and x86 memory addressing
- Arithmetic instructions

Next Time

- Control instructions (branch, etc.)

Extra slides

Pentium Pro (P6)

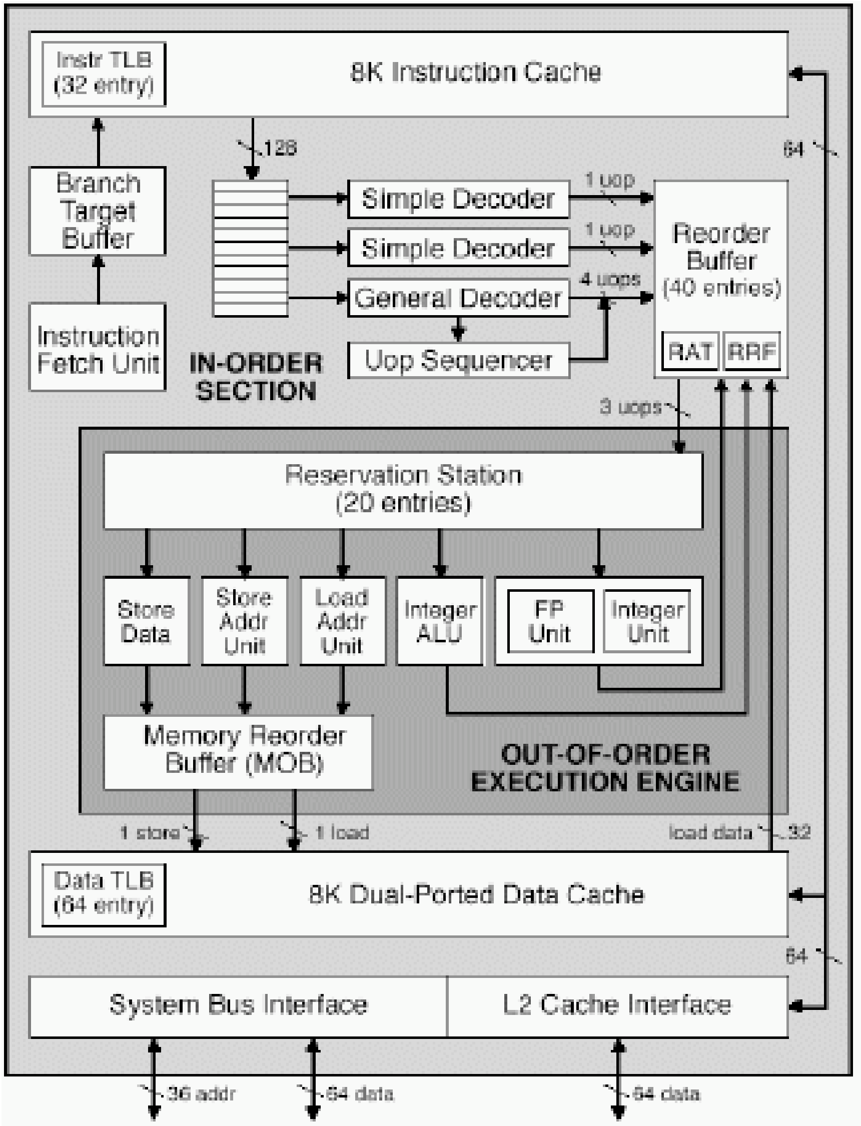
History

- Announced in Feb. ‘95
- Basis for Pentium II, Pentium III, and Celeron processors
- Pentium 4 similar idea, but different details

Features

- Dynamically translates instructions to more regular format
  - Very wide, but simple instructions
- Executes operations in parallel
  - Up to 5 at once
- Very deep pipeline
  - 12–18 cycle latency

PentiumPro Block Diagram



Microprocessor Report
2/16/95

PentiumPro Operation

Translates instructions dynamically into “Uops”

- 118 bits wide
- Holds operation, two sources, and destination

Executes Uops with “Out of Order” engine

- Uop executed when
  - Operands available
  - Functional unit available
- Execution controlled by “Reservation Stations”
  - Keeps track of data dependencies between uops
  - Allocates resources

Consequences

- Indirect relationship between IA32 code & what actually gets executed
- Tricky to predict / optimize performance at assembly level