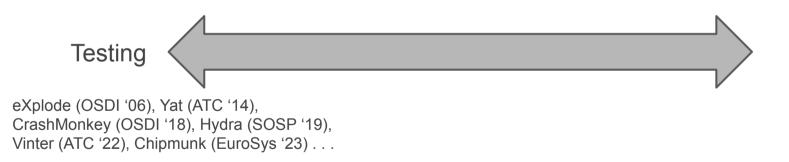
# SquirrelFS: using the Rust compiler to check file-system crash consistency

#### Hayley LeBlanc, Nathan Taylor, James Bornholt, Vijay Chidambaram



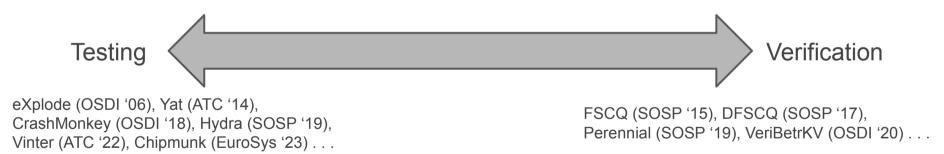




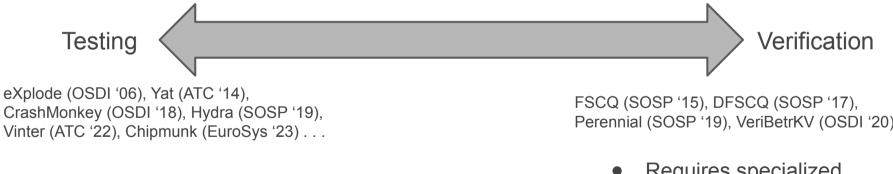


eXplode (OSDI '06), Yat (ATC '14), CrashMonkey (OSDI '18), Hydra (SOSP '19), Vinter (ATC '22), Chipmunk (EuroSys '23) . . .

- Incomplete
- Requires specialized tools



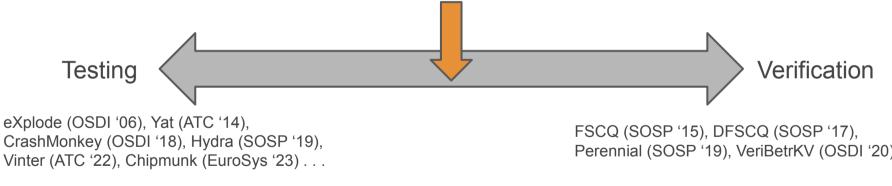
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- Incomplete
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Perennial (SOSP '19), VeriBetrKV (OSDI '20) . . .

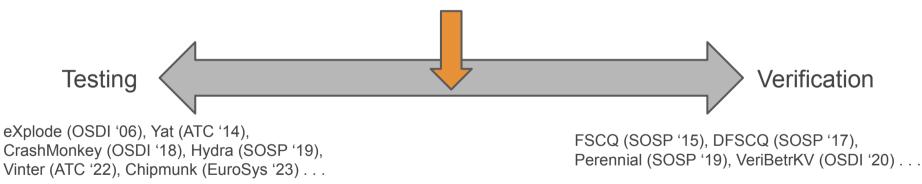
- Requires specialized expertise
- Development takes longer
- Often impacts performance 6



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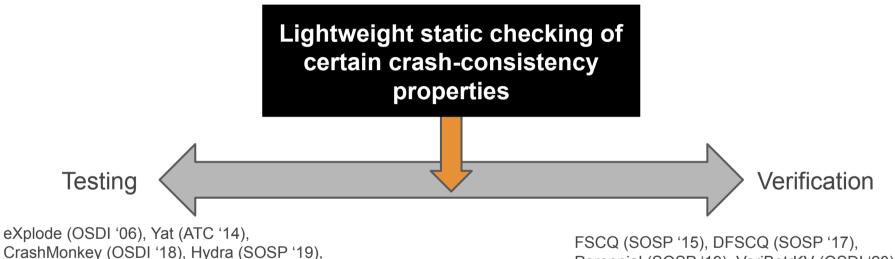
Perennial (SOSP '19), VeriBetrKV (OSDI '20) . . .

- **Requires specialized** expertise
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- Often impacts performance 7



- Incomplete
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- No specialized expertise or tools required
- Statically check critical properties

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Vinter (ATC '22), Chipmunk (EuroSys '23) . . .

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High-performance, low-level systems programming language

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**Strong type system** that can statically prevent:

High-performance, low-level systems programming language

**Strong type system** that can statically prevent:

• Data races

High-performance, low-level systems programming language

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- Memory safety issues

High-performance, low-level systems programming language

**Strong type system** that can statically prevent:

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- **Some crash-consistency bugs!** (Corundum ASPLOS '21)

High-performance, low-level systems programming language

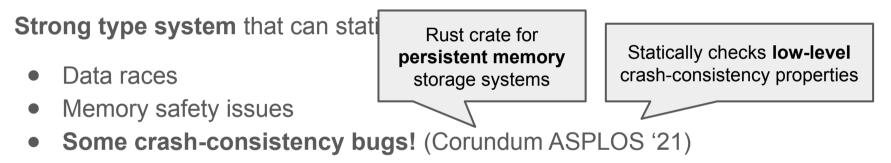
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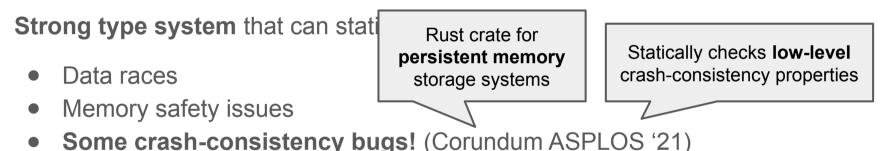
Rust crate for persistent memory storage systems

• Some crash-consistency bugs! (Corundum ASPLOS '21)

High-performance, low-level systems programming language

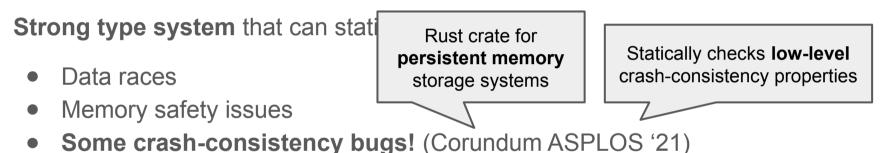


High-performance, low-level systems programming language



This work: use Rust to statically check **higherlevel** crash-consistency properties in a persistent memory file system

High-performance, low-level systems programming language



This work: use Rust to statically check **higherlevel** crash-consistency properties in a persistent memory file system

Atomicity of system calls

Persistent memory file system with statically-checked ordering-related crashconsistency properties

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Static checks rely only on existing Rust features

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Uses the **typestate pattern** to statically check ordering of durable updates

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Achieves similar or better performance to other PM file systems

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https://github.com/utsaslab/squirrelfs

#### Roadmap

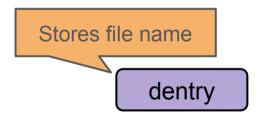
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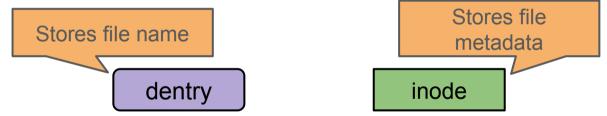
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Crash consistency depends on the **order of durable updates** (Ganger & Patt '94, Frost et al. '07, Chidambaram et al. '12)

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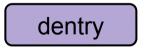


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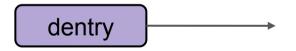


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## Ordering for crash consistency

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Simple example: creating a new file

Setting dentry pointer depends on inode initialization

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Crash consistency depends on the **order of durable updates** (Ganger & Patt '94, Frost et al. '07, Chidambaram et al. '12)

Simple example: creating a new file

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#### Statically enforcing durable update ordering can prevent many crash-consistency bugs

Track dependencies between durable in-place updates to enforce crash-consistent ordering (Ganger & Patt OSDI '94)

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Bob Beck, OpenBSD commit message, 2023

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 Fast synchrony with persistent memory

 Getting the update ordering right
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- Intel Optane DC PM
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Statically check ordering with typestate

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Encode runtime state in an object's type with no runtime overhead

Encode runtime state in an object's type with no runtime overhead

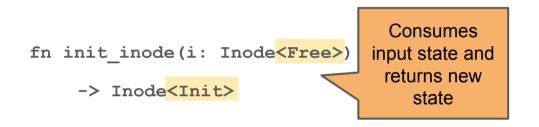
fn init\_inode(i: &mut Inode)

Encode runtime state in an object's type with no runtime overhead

fn init\_inode(i: Inode)

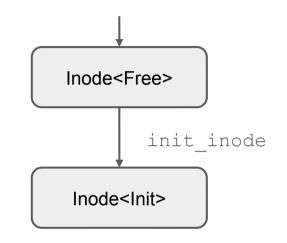
-> Inode<Init>

Encode runtime state in an object's type with no runtime overhead

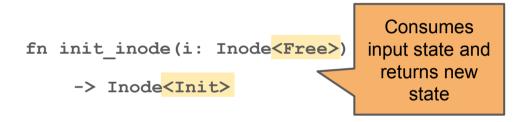


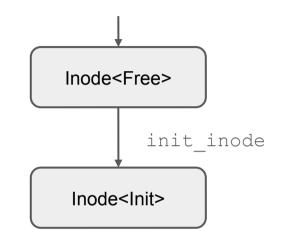
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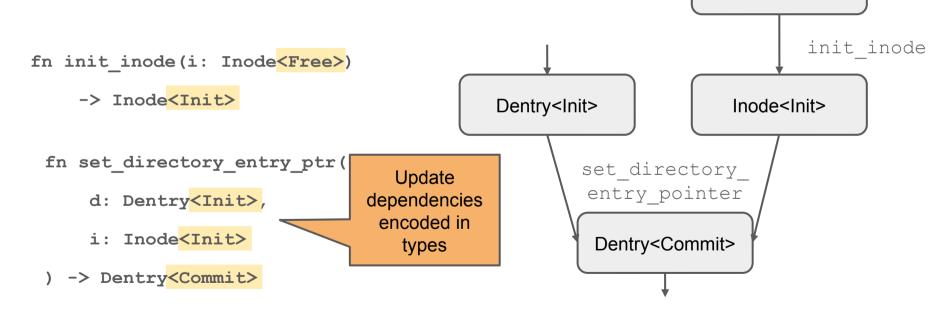
fn set\_directory\_entry\_ptr(

- d: Dentry<Init>,
- i: Inode<Init>
- ) -> Dentry<Commit>

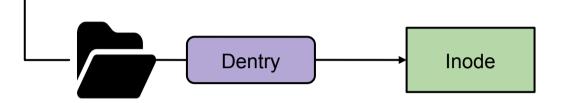
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Inode<Free>

Encode runtime state in an object's type with no runtime overhead



Inode<Free>



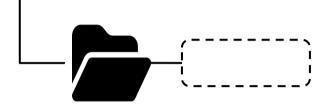
fn create file(name: String) {

}

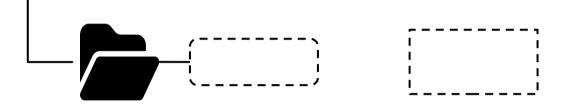
fn create\_file(name: String) {

}

```
let d = Dentry::get_free_dentry(); // obtain Dentry<Free>
let i = Inode::get_free_ino(); // obtain Inode<Free>
let d = d.set_name(name); // Dentry<Free> -> Dentry<Init>
let d = d.set_directory_entry_ptr(i); // BUG!!
```



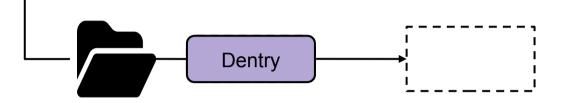
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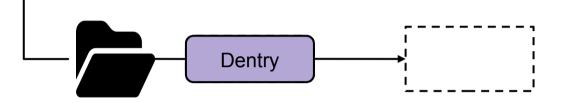
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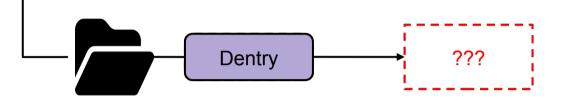
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## SquirrelFS implementation

Typestate-checked Synchronous Soft Updates for crash consistency 7500 LOC of Rust

Simple durable layout with volatile indexes and allocators

Atomic metadata-related system calls (including rename)

Modeled as a transition system and model checked in Alloy

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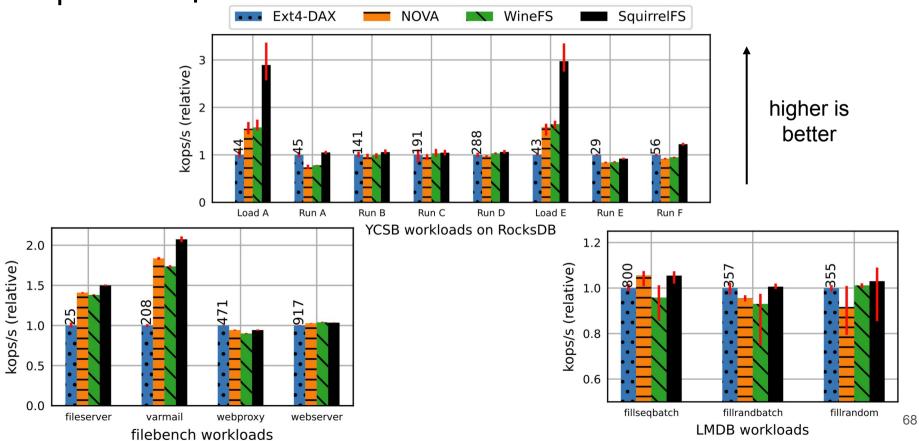
## Evaluation

Evaluated on 128GB Intel Optane DC Persistent Memory Module

Compared against Ext4-DAX, NOVA, and WineFS

- 1. How does SquirrelFS compare to other PM file systems?
- 2. How long does it take to statically check SquirrelFS's crash-consistency properties?

## SquirrelFS performance



System (verified)	Lines of code	Verification time (s)
FSCQ	31K	39600
VeriBetrKV	45K	6480

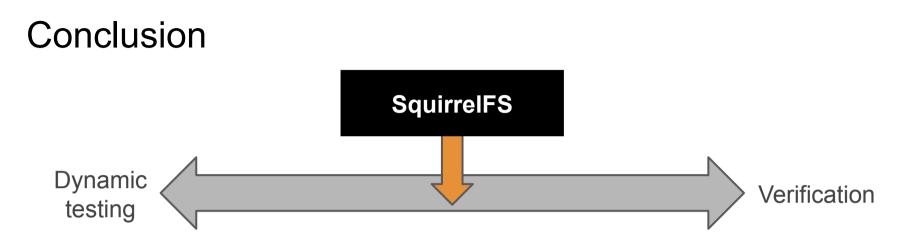
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System (unverified)	Lines of code	Compilation time (s)
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NOVA	16K	20
WineFS	9К	13

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Ext4	45K	38
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WineFS	9К	13

System (typestate-checked)	Lines of code	Compile+check time (s)
SquirrelFS	7.5K	10



Typestate pattern statically checks ordering for crash consistency

Synchronous Soft Updates crash-consistency mechanism

Comparable performance to existing PM file systems

https://github.com/utsaslab/squirrelfs



# Extra slides

# Background: persistent memory

Low latency on the order of DRAM

Byte-addressable via memory loads and stores

Cache-line flushes and memory fences for durability and ordering

Examples:

- Intel Optane DC Persistent Memory Module
- Battery-backed DRAM
- Future devices: Micron, startups, CXL.mem, ...

# Background: soft updates

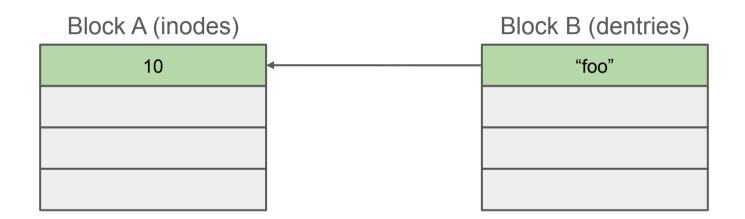
Crash-consistency mechanism based on ordering in-place updates

Rules:

- 1. Never point to a structure before it is initialized
- 2. Never reuse a resource before nullifying existing references to it
- 3. Never reset the old pointer to a resource before setting the new one

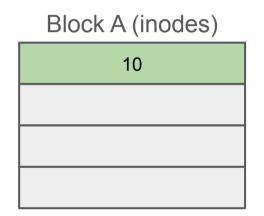
Enforced by tracking update dependencies and ordering durable updates

Reduces write amplification, but increases complexity

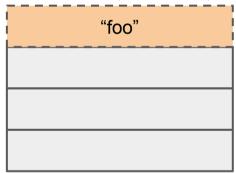


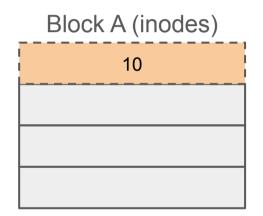




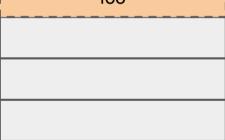


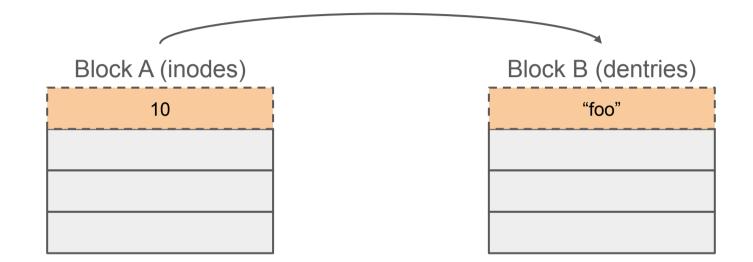
#### Block B (dentries)

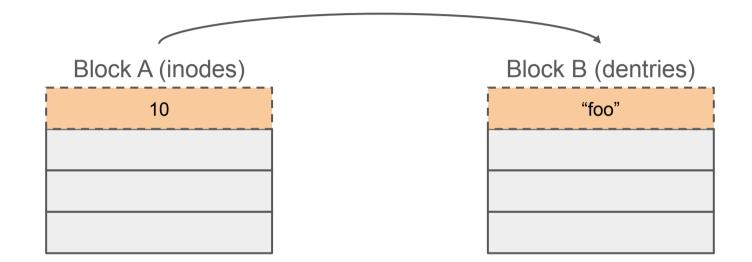






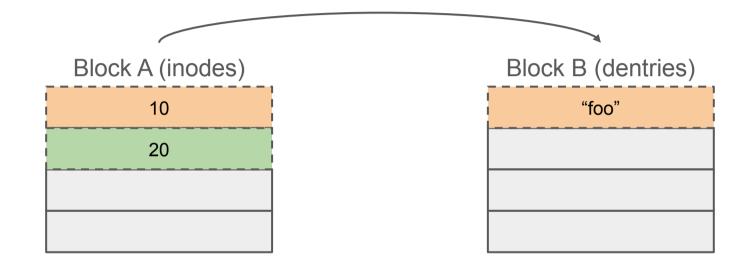






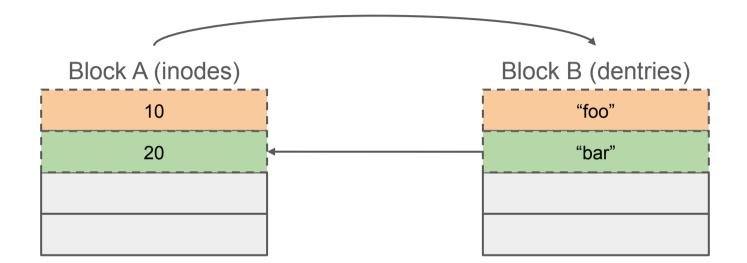
1. unlink foo

#### 2. create bar

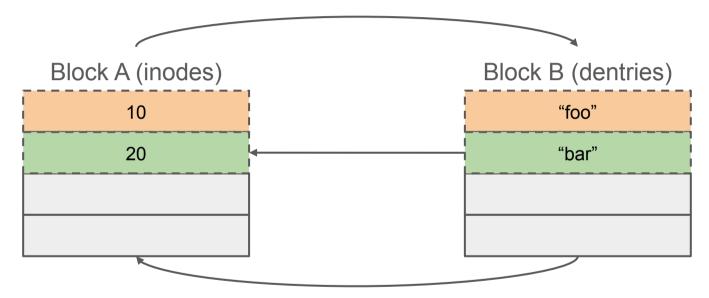


1. unlink foo

#### 2. create bar



- 1. unlink foo
- 2. create bar



1. unlink foo

2. create bar

# Background: verification

Mathematical proof that a program is correct

Prove that the complex implementation matches a simpler specification of correctness

Developer writes a proof, computer checks it

Uses verification-aware programming languages or interactive theorem provers

E.g.: Verus verification framework for Rust

#### Typestate in Rust: update operations

```
impl Inode<Clean,Free> {
    fn init(self,...) -> Inode<Dirty,Init> {...}
}
impl Dentry<Clean,Free> {
    fn set name(self, name: String) -> Dentry<Dirty,Init> {...}
}
impl Dentry<Clean,Init> {
    fn set ino(self, ino: Inode<Clean,Init>) -> Dentry<Dirty,Committed> {...}
}
```

#### Typestate: ensuring persistence

```
impl<S> Inode<Dirty,S> {
```

}

}

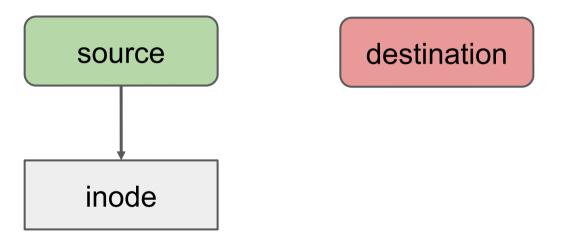
```
fn flush(self) -> Inode<InFlight,S> {...}
```

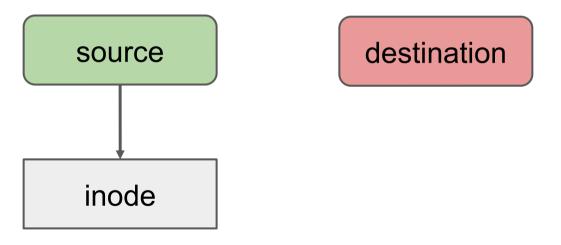
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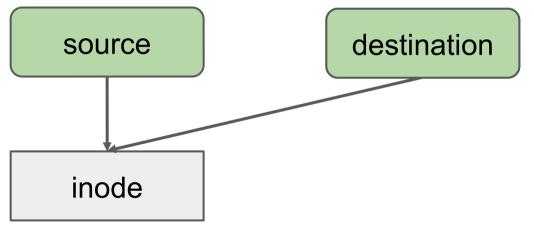
```
fn fence(self) -> Inode<Clean,S> {...}
```

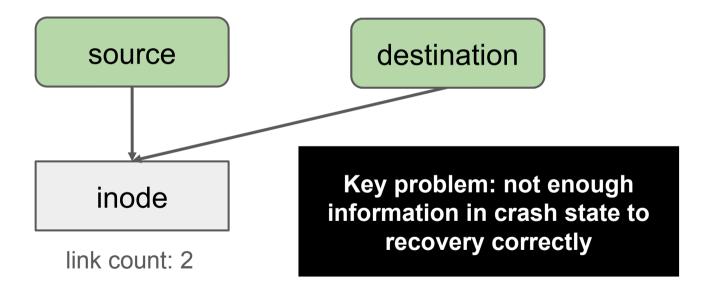
# Directory entry validity rules

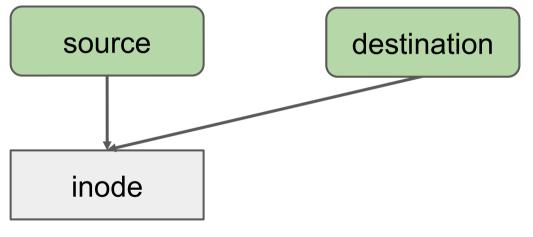
- 1. If a dentry's inode number is 0, the dentry is invalid.
- 2. If dst's rename pointer points to src, then:
  - a. If dst.inode != src.inode, both dentries are valid
  - b. If dst.inode == src.inode, src is invalid

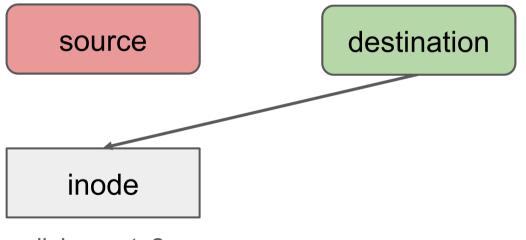


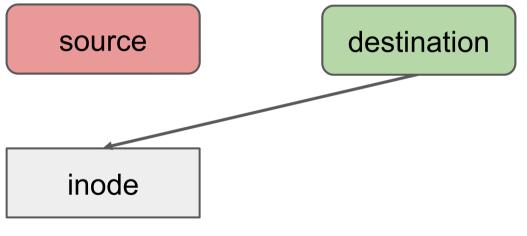


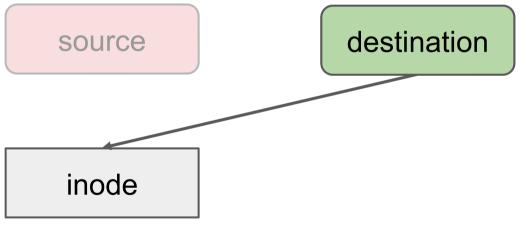


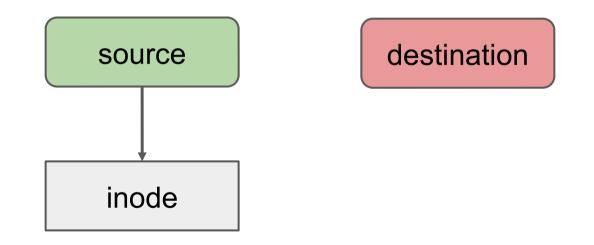


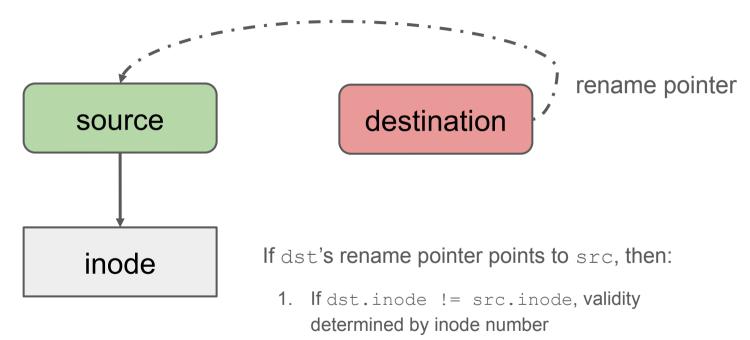




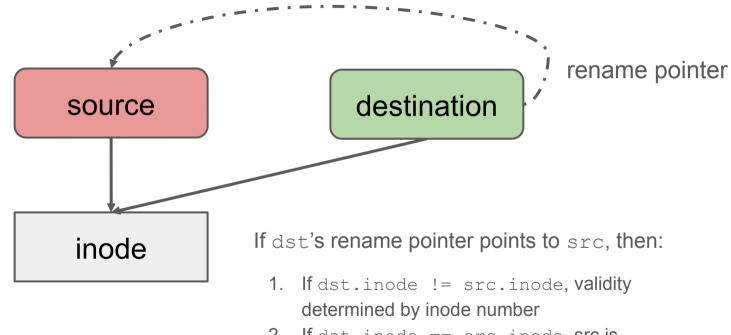




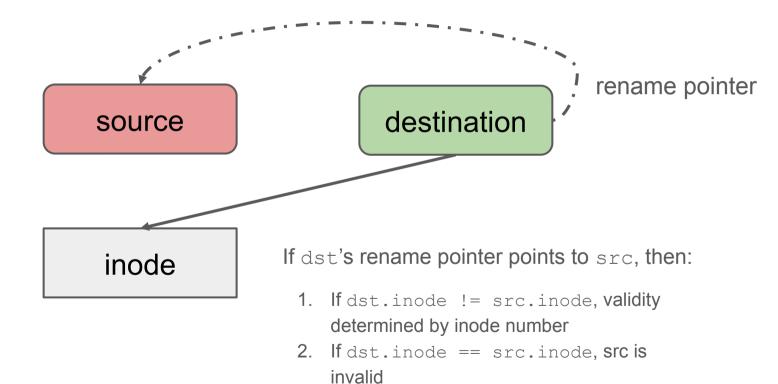


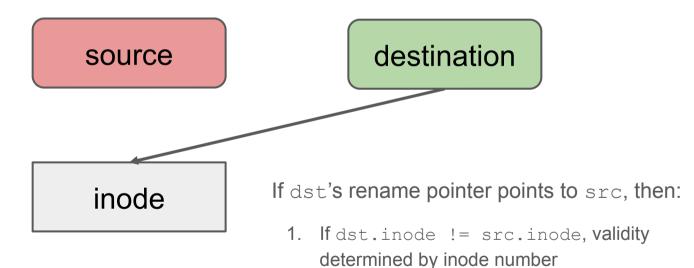


2. If dst.inode == src.inode, src is
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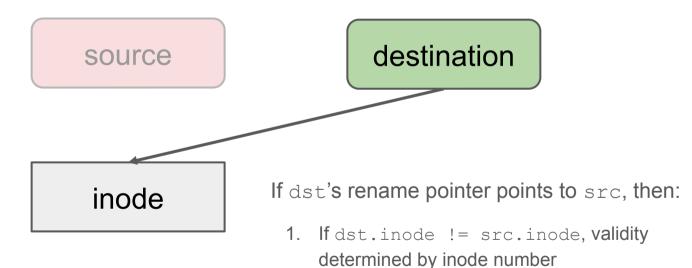


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# Typestate in SquirrelFS

Operational typestate

- What operations have been performed on this object?
- Is it free? Initialized? Allocated but not initialized?

Persistence typestate

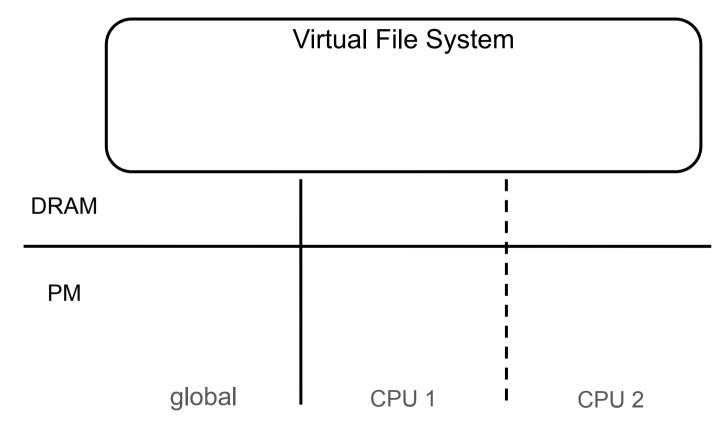
• Have the most recent updates been made durable?

Typestate transition functions make persistent updates and return the new typestate

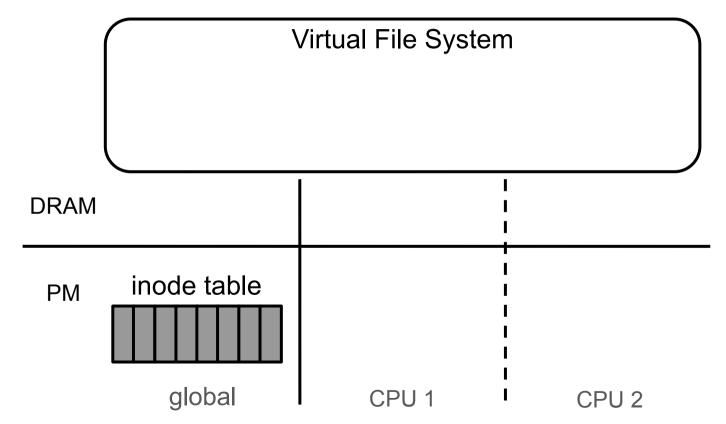
# SquirrelFS crash consistency bugs

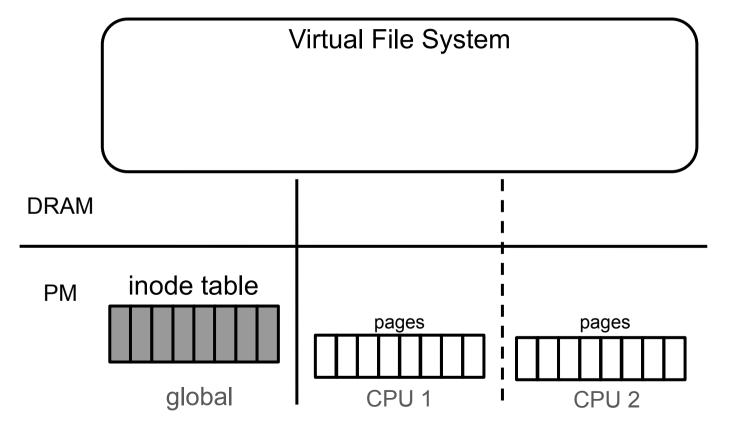
- 1. A cache line flush persistence function was passed a reference to a page pointer, rather than the page pointer itself (typestate transition body)
- 2. Missing case to free orphaned dir pages (recovery code)
- 3. Allocated but orphaned directory entries towards parent link count (recovery code)
- 4. Used persistent inode number, rather than inode table index, in inode table scan (recovery code)

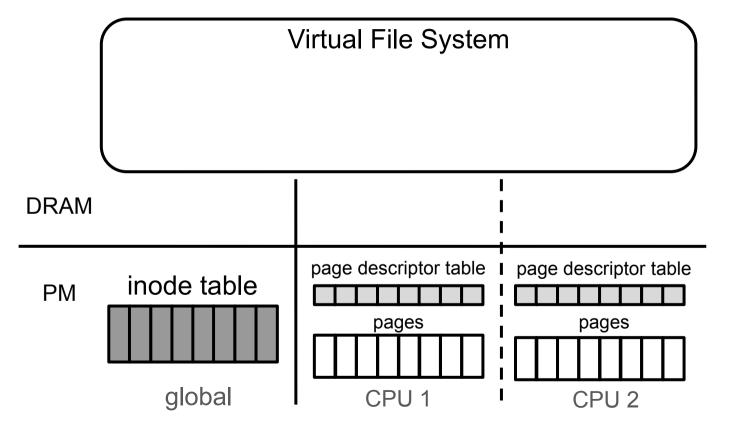
# SquirrelFS architecture

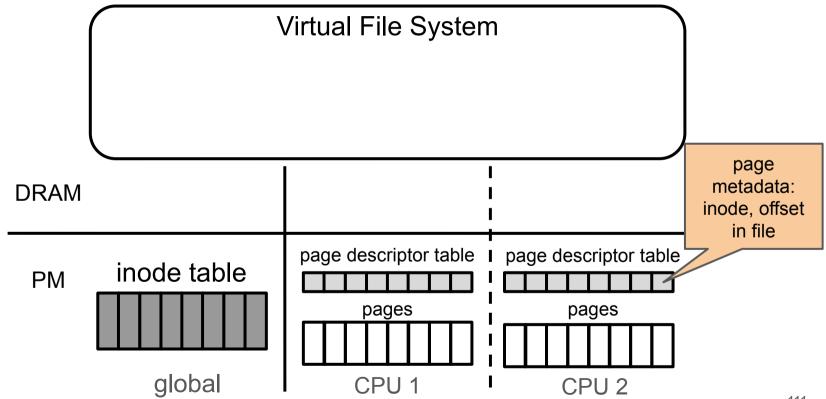


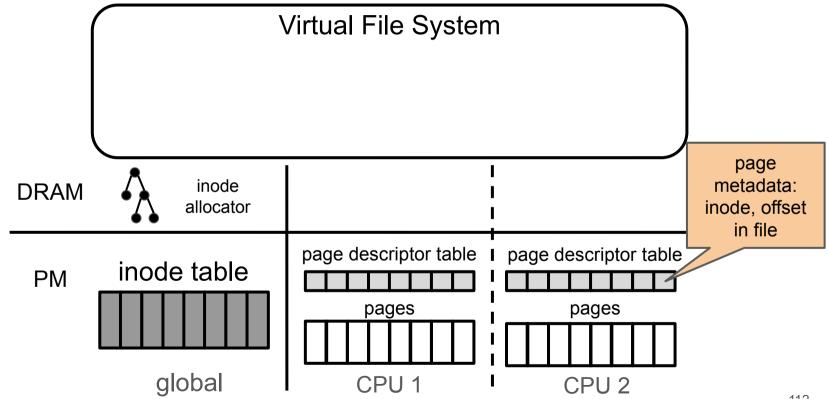
### SquirrelFS architecture

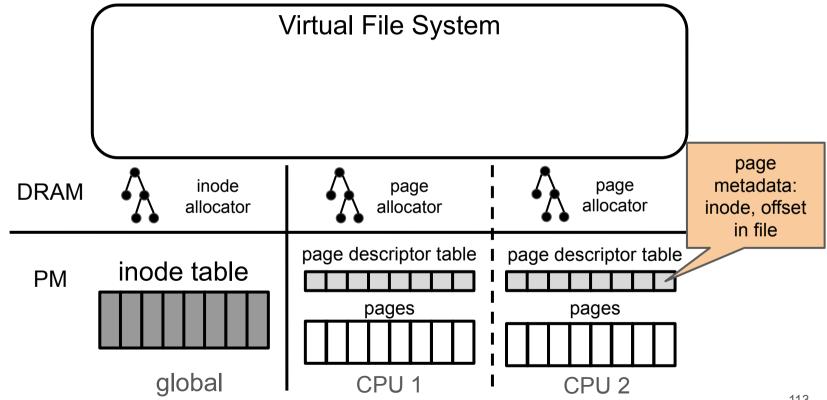


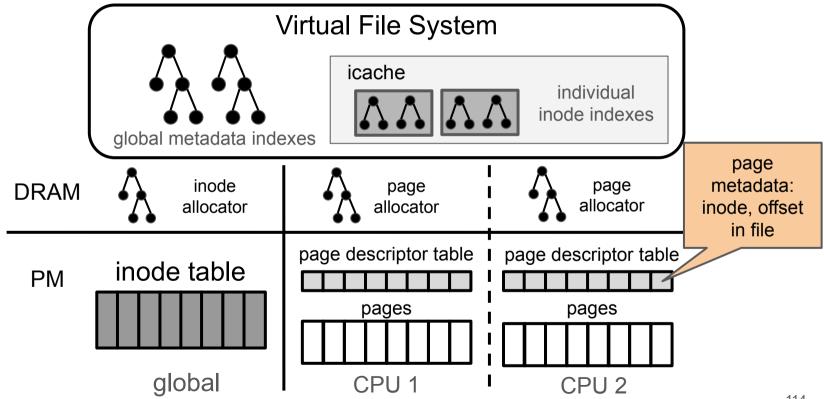












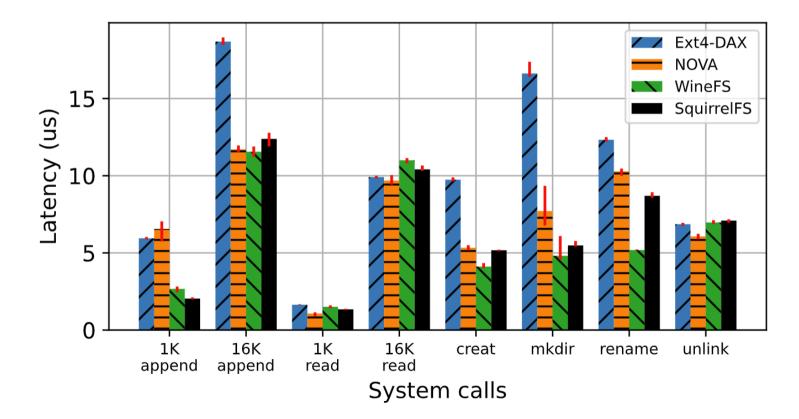
## Bugs found with typestate

- Missing persistence primitives
  - E.g.: initial implementation of write was missing flush/fence calls after setting new page backpointer
- Incorrect ordering
  - E.g.: initial rename implementation incorrectly updated link count before clearing a directory entry, which could result in a dangling link later on

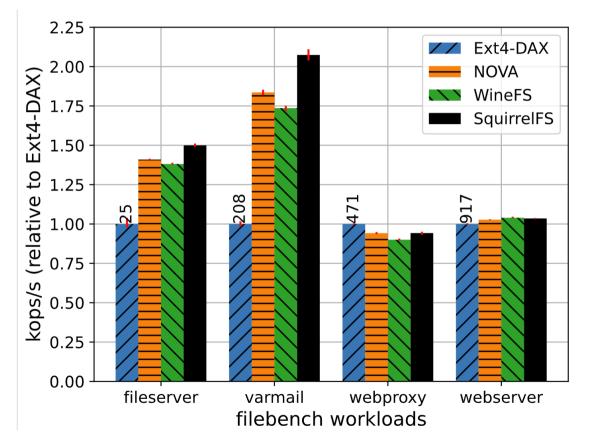
# Bugs found with Alloy model

- Recovering from renames
  - Initial model did not include any crash recovery logic; we believed it was not necessary
  - Model found a counterexample where invalid directory entries could reappear after a crash during rename
  - Fixed by adding mandatory post-crash cleanup of rename pointers
- Handling . and .. dentries
  - Originally stored durably and included in update ordering rules
  - Alloy model repeatedly found issues with these rules, particularly during rename
  - Now stored only in volatile memory

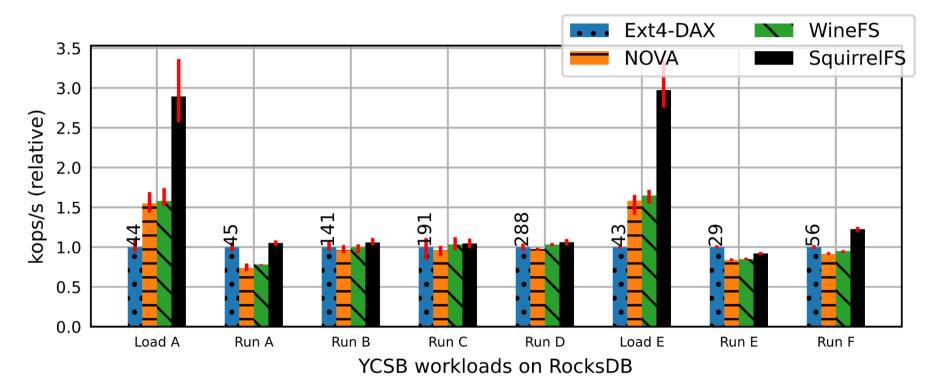
#### Microbenchmark: system call latency



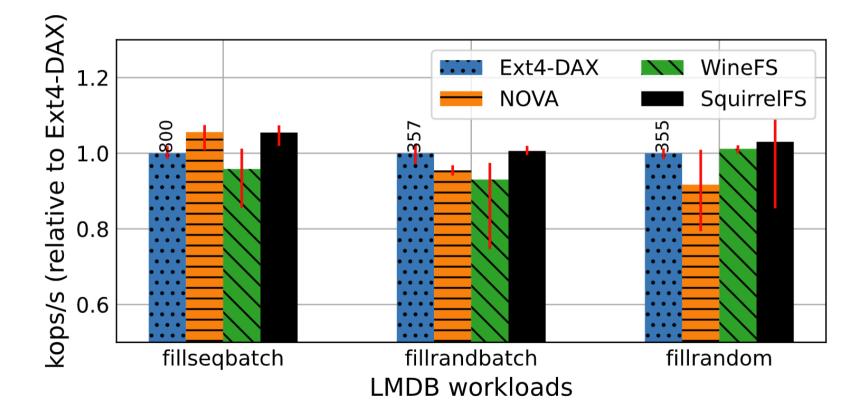
### Macrobenchmark: filebench



### Application benchmark: RocksDB

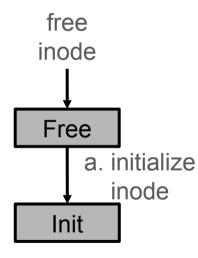


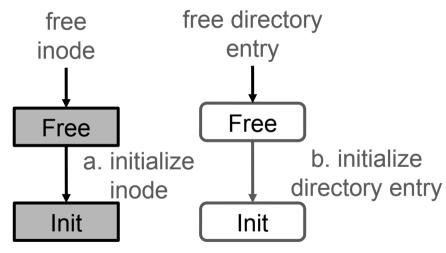
### Application benchmark: LMDB

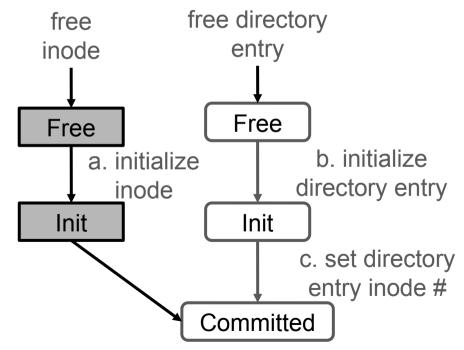


### SquirrelFS mount times

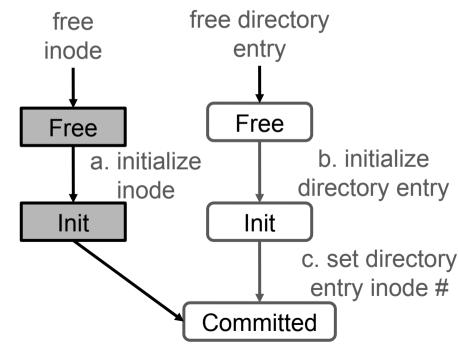
	System state	Mount time (s)
Normal mount	mkfs	5.80
	Empty	5.51
	Full	30.50
Recovery mount	Empty	5.76
	Full	55.50





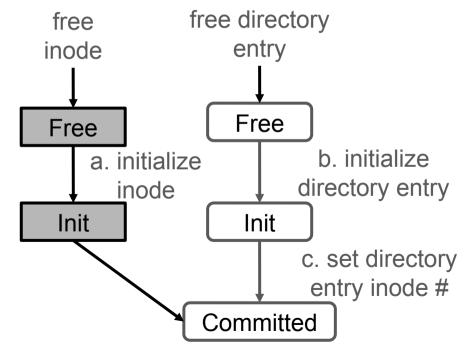


Soft updates: crash consistency from ordered in-place durable updates



Managing update dependencies in **asynchronous** soft updates is notoriously difficult

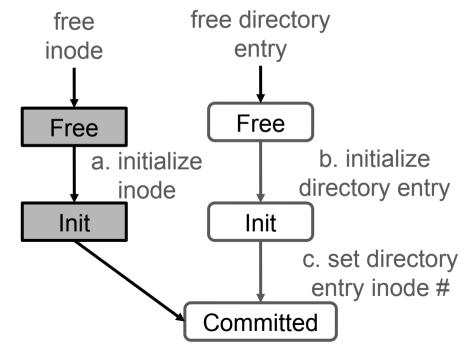
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**Synchronous** soft updates eliminates most complexity!

Soft updates: crash consistency from ordered in-place durable updates

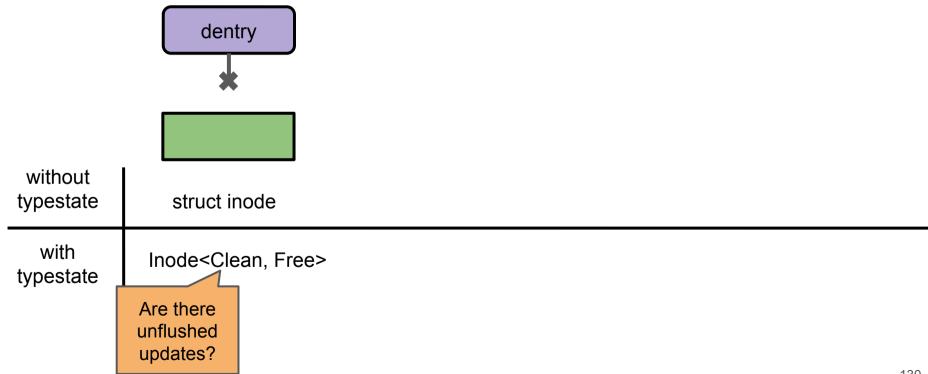


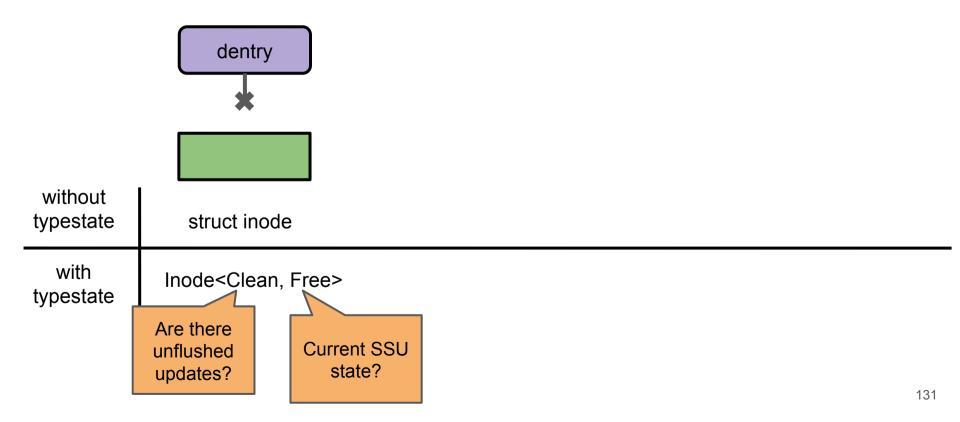
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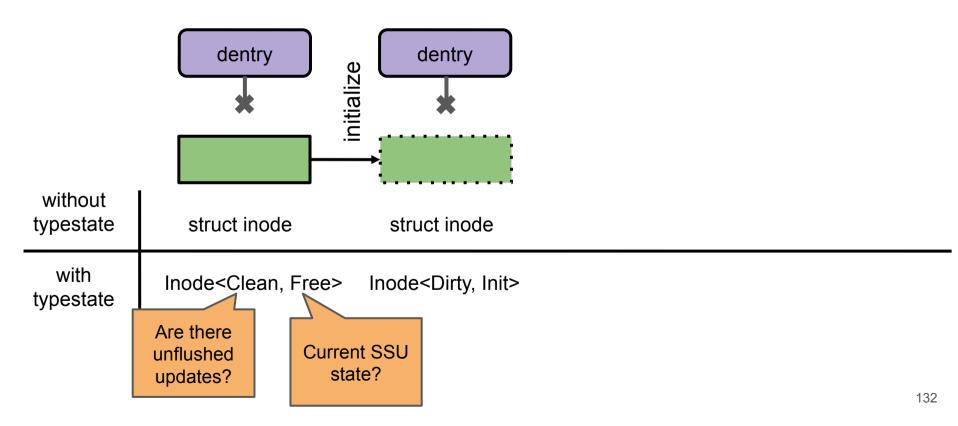
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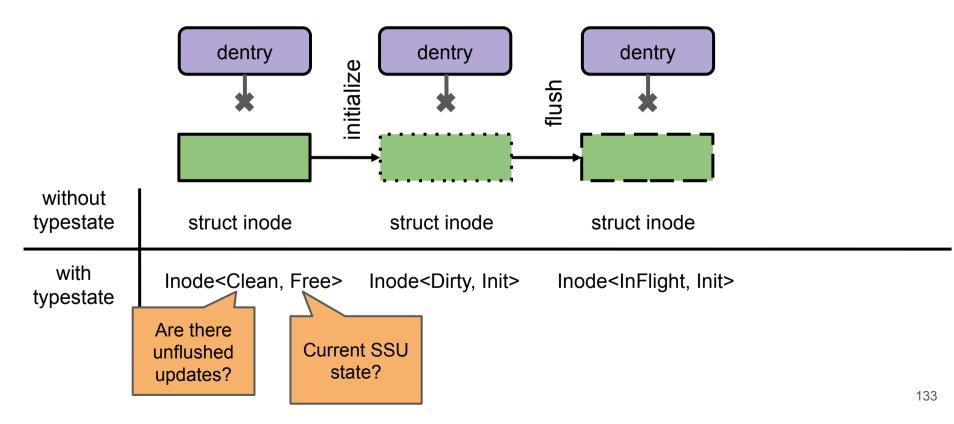
Fast **persistent memory** storage enables performant synchrony

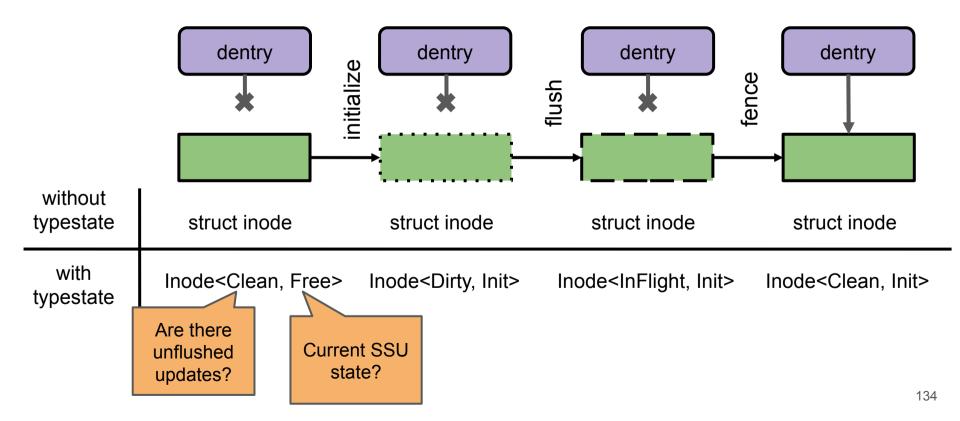


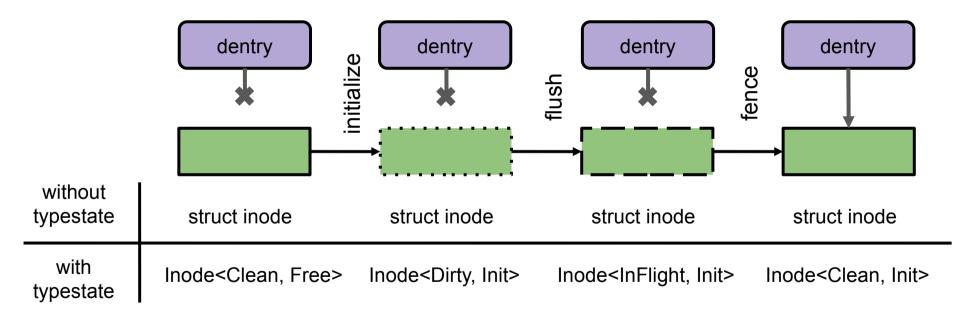












Ordering encoded in function signatures:

impl Inode<Clean,Free> {fn init(self) -> Inode<Dirty, Init> {...}}