

# EXPANDED DATA TRANSFORMATION

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# Assumptions

$$A \xrightarrow[\alpha']{\alpha} A' \vee A' \xrightarrow[\alpha']{\alpha} A$$

equivalent



$$\alpha A \wedge \alpha' A' \wedge \alpha' \alpha \wedge \alpha \alpha' \wedge \alpha' \alpha' \wedge \alpha \alpha \wedge \alpha' \alpha' \wedge \alpha \alpha'$$

and optionally

$$G[A \xrightarrow[\alpha']{\alpha} A'] \vee G[A' \xrightarrow[\alpha']{\alpha} A]$$

equivalent

$$B \xrightarrow[\beta']{\beta} B' \vee B' \xrightarrow[\beta']{\beta} B$$

equivalent



$$\beta B \wedge \beta' B' \wedge \beta' \beta \wedge \beta \beta' \wedge \beta' \beta' \wedge \beta \beta \wedge \beta' \beta' \wedge \beta \beta'$$

and optionally

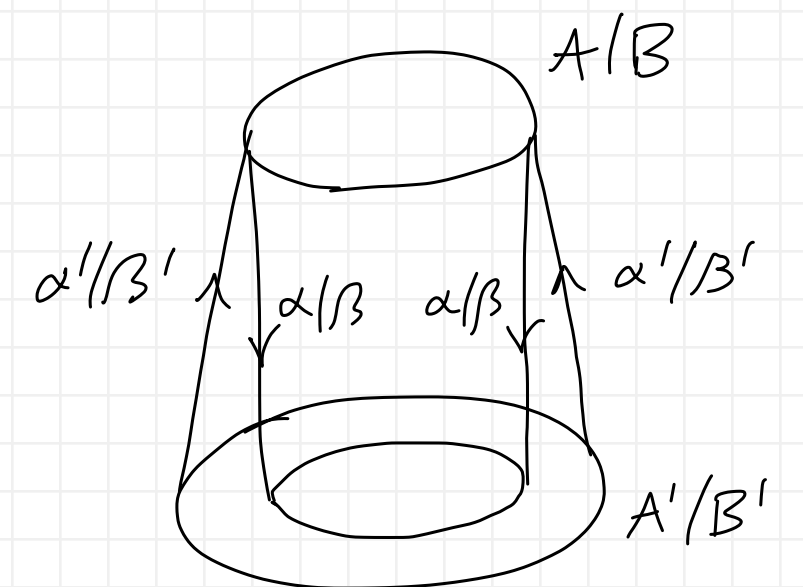
$$G[B \xrightarrow[\beta']{\beta} B'] \vee G[B' \xrightarrow[\beta']{\beta} B]$$

equivalent

(see separate 'Domain Mappings' notes)

key property:  $\alpha', \beta'$  surjective — they map the new expanded data to all the old data

different choices of  $\alpha, \beta$  are possible in general for the same  $\alpha', \beta'$ ,  
i.e. the same old data may be expanded in different ways as new data



# Default Approach

use  $\alpha, \beta$  as inverses of  $\alpha', \beta'$



everything works the same as in the isomorphic data transformation, whose proofs only use  $\alpha'\alpha$  and  $\beta'\beta$  but not  $\alpha\alpha'$  and  $\beta\beta'$  (which are missing here), but some pictures are slightly different since things may not be isomorphic:

