

# Introduction to Rippling

John Erickson  
Dept. of Computer Science  
University of Texas at Austin

October 10, 2007

# Outline

- ▶ What is Rippling?
- ▶ Rippling-Out
- ▶ Rippling-In
- ▶ Rippling-Sideways
- ▶ Rippling-Across
- ▶ Conclusions

# What is Rippling?

- ▶ A way to control rewriting during induction
- ▶ Based on “rippling-out” the differences between the IH and IC
- ▶ Extended to achieve other rewriting goals
- ▶ Uses annotated rewrite rules to guide rewriting and ensure termination

# Rippling-Out

- ▶ IH  $x + (y + z) = (x + y) + z$
- ▶ IC  $\boxed{s(x)} + (y + z) = (\boxed{s(x)} + y) + z$
- ▶  $s(\dots)$  is a *wave-front*
- ▶  $x$  is a *wave-hole*
- ▶  $x + (y + z) = (x + y) + z$  is the *skeleton*

# Wave Rules

- ▶ General Form:  $\eta(\underline{\xi(\underline{\mu})}) \rightarrow \underline{\zeta(\underline{\eta(\underline{\mu})})}$
- ▶  $\underline{s(U)} + V \rightarrow \underline{s(U + V)}$
- ▶  $\underline{s(U)} \times V \rightarrow \underline{U \times V + V}$
- ▶  $\text{even}(\underline{s(s(U))}) \rightarrow \text{even}(U)$
- ▶  $U + (\underline{V + W}) \rightarrow \underline{(U + V) + W}$
- ▶  $(\underline{U + V}) + W \rightarrow \underline{U + (V + W)}$

# Rippling-Out Example

- ▶ IH  $x + (y + z) = (x + y) + z$
- ▶ IC  $\boxed{s(x)} + (y + z) = (\boxed{s(x)} + y) + z$
- ▶  $\boxed{s(x + (y + z))} = \boxed{s(x + y)} + z$
- ▶  $\boxed{s(x + (y + z))} = \boxed{s((x + y) + z)}$

# Rippling-In

- ▶ Useful for when one side of an equality is missing a wave rule
- ▶ IH  $half(x + x) = x$
- ▶ IC  $half(\boxed{s(\underline{x})} + \boxed{s(\underline{x})}) = \boxed{s(\underline{x})}$
- ▶  $half(\boxed{s(x + \boxed{s(\underline{x})})}) = \boxed{s(\underline{x})}$
- ▶ Missing:  $U + \boxed{s(\underline{V})} \rightarrow \boxed{s(U + V)}$
- ▶ XF  $half(s(x + s(x))) = \boxed{s(half(x + x))} \downarrow$
- ▶  $half(s(x + s(x))) = half(\boxed{s(s(\underline{x} + x))} \downarrow)$
- ▶  $x + s(x) = s(x + x)$

# Rippling-Sideways

- ▶ Unmeasured (free) induction variables can be used as “sinks”
- ▶ Rippling-sideways Attempts to ripple wave-front into a sink
- ▶ General form:  $\eta(\boxed{\xi(\underline{\mu})}^{\uparrow}, \nu) \rightarrow \eta(\mu, \boxed{\zeta(\underline{nu})}^{\downarrow})$
- ▶ IH  $rev(l) \langle \rangle M = qrev(l, M)$
- ▶ IC  $rev(\boxed{h :: \underline{l}}^{\uparrow}) \langle \rangle [m] = qrev(\boxed{h :: \underline{l}}^{\uparrow}, [m])$
- ▶  $rev(\boxed{h :: \underline{l}}^{\uparrow}) \langle \rangle [m] = qrev(l, \boxed{h :: \underline{m}}^{\downarrow})$
- ▶  $\boxed{rev(l) \langle \rangle (h :: nil)^{\uparrow}} \langle \rangle [m] = qrev(l, \boxed{h :: \underline{m}}^{\downarrow})$
- ▶  $rev(l) \langle \rangle (\boxed{(h :: nil) \langle \rangle \underline{m}}^{\downarrow}) = qrev(l, \boxed{h :: \underline{m}}^{\downarrow})$



# Rippling-Across

- ▶ Adapts rippling to destructor inductions

- ▶  $U + V =$  if  $U = 0$  then  $V$  else  $s(p(U) + V)$

- ▶ Creational Rule:  $U \neq 0 \implies U + V \rightarrow \boxed{s(\underline{p(U)}^- + V)}^\uparrow$

- ▶  $\boxed{p(x)}^\uparrow + (y+z) = (\boxed{p(x)}^\uparrow + y) + z \vdash x + (y+z) = (x+y) + z$

- ▶  $\boxed{p(x)}^\uparrow + (y+z) = (\boxed{p(x)}^\uparrow + y) + z \vdash$

$$\boxed{s(\underline{p(x)}^- + (y+z))}^\uparrow = \boxed{s(\underline{p(x)}^- + y)}^\uparrow + z$$

- ▶  $p(x) + (y+z) = (p(x) + y) + z \vdash \boxed{s(p(x) + (y+z))}^\uparrow = \boxed{s(p(x) + y)}^\uparrow + z$

# Conclusions

- ▶ Pros: Terminating rewrites, ability to use rules right to left, goal directed
- ▶ Cons: A little bit complicated
- ▶ Doesn't address generalization, lemma generation, or choosing induction scheme
- ▶ Some techniques address these problems by asking “how can I make this choice so that rippling will be facilitated?”
- ▶ Formalizes informal strategies for rewriting