

# Proving Unbounded Theorems with the Help of GL

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# Agenda

- A technique for proving unbounded theorems with the help of GL.
- Benefit of using that technique in certifying the 32-bit physical memory model.

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- Proving **unbounded** theorems might not be trivial if they are complicated. However, if they can be transformed into **bounded** theorems, then we can use **GL** to solve the problem.

# Approach

- **GL** is a symbolic simulation framework for proving **bounded** ACL2 theorems. It cannot prove theorems that contain **unbounded** variables.
- Proving **unbounded** theorems might not be trivial if they are complicated. However, if they can be transformed into **bounded** theorems, then we can use **GL** to solve the problem.
- Present a trick that proves **unbounded** theorems with the help of **GL**.

# Simple Example

- (implies (integerp  $x$ )  
    (equal (mod  $x$  8)  
          (logand  $x$  7))))

Although  $x$  is **unbounded**, only its **3 low bits** affect the computation in the above theorem. So, we can transform it to the **bounded** lemma:

- (implies (integerp  $x$ )  
    (equal (mod  $x[3:0]$  8)  
          (logand  $x[3:0]$  7))))

where  $x[j:i]$  represents the bit string from index  $i$  to  $j$  of  $x$  ( $0 \leq i \leq j$ ).

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    (equal (mod  $x[3:0]$  8)  
          (logand  $x[3:0]$  7))))

Then, the **unbounded** theorem will follow by applying two following rewrite rules:

- (equal (mod  $x[3:0]$  8) (mod  $x$  8))
- (equal (logand  $x[3:0]$  7) (logand  $x$  7))



# Main Theorem

```
(defthm main
  (implies (and (natp i02)
                 (<= i02 2))
            (equal (logior (mod (ash mem-val (* -8 i02))
                                *2^8*))
                   (* *2^8*
                      (mod (ash mem-val (+ -8 (* -8 i02)))
                          *2^8*))))
            (mod (ash mem-val (* -8 i02))
                 *2^16*))))
```

# Main Theorem

```
(defthm main
  (implies (and (natp i02)
                (<= i02 2))
    (equal (logior E0
      (* *2^8*
        (mod (ash mem-val (+ -8 (* -8 i02)))
          *2^8*)))
      (mod (ash mem-val (* -8 i02))
        *2^16*))))))
```

# Main Theorem

```
(defthm main
  (implies (and (natp i02)
                (<= i02 2))
            (equal (logior E0
                      (* *2^8*
                        E1))
                   (mod (ash mem-val (* -8 i02))
                         *2^16*))))
```

# Main Theorem

```
(defthm main
  (implies (and (natp i02)
                (<= i02 2))
            (equal (logior E0
                    (* *2^8*
                      E1))
                  E2))))
```





# Analyze E1

- $E1 = (\text{mod} (\text{ash } \text{mem-val} (+ -8 (* -8 i02)))$   
     $*2^8*)$   
     $= \text{mem-val}[(+ 15 (* 8 i02)) : (+ 8 (* 8 i02))]$

# Analyze E1

- $E1 = (\text{mod} (\text{ash } \text{mem-val} (+ -8 (* -8 i02)))$   
 $*2^8*)$   
 $= \text{mem-val}[(+ 15 (* 8 i02)) : (+ 8 (* 8 i02))]$
- $(\text{and} (<= 0 i02) \quad \Rightarrow \quad (\text{and} (<= 8 (+ 8 (* 8 i02)))$   
 $(<= i02 2))$   
 $(<= (+ 8 (* 8 i02)) 24)$   
 $(<= 15 (+ 15 (* 8 i02)))$   
 $(<= (+ 15 (* 8 i02)) 31))$

$\Rightarrow$  Only  $\text{mem-val}[31 : 8]$  of  $\text{mem-val}$  affects E1.



# Analyze E2

- $E2 = (\text{mod} (\text{ash } \text{mem-val} (* -8 i02))$   
     $*2^{16}*)$   
     $= \text{mem-val}[(+ 15 (* 8 i02)) : (* 8 i02)]$

# Analyze E2

- $E2 = (\text{mod} (\text{ash } \text{mem-val} (* -8 i02))$   
 $*2^{16})$   
 $= \text{mem-val}[(+ 15 (* 8 i02)) : (* 8 i02)]$
- $(\text{and} (<= 0 i02) \Rightarrow (\text{and} (<= 0 (* 8 i02)$   
 $(<= i02 2)))$   
 $(<= (* 8 i02) 16)$   
 $(<= 15 (+ 15 (* 8 i02)))$   
 $(<= (+ 15 (* 8 i02)) 31))$

$\Rightarrow$  Only  $\text{mem-val}[31 : 0]$  of  $\text{mem-val}$  affects E2.

# Claim

- Only `mem-val[23 : 0]` of `mem-val` affects E0.
- Only `mem-val[31 : 8]` of `mem-val` affects E1.
- Only `mem-val[31 : 0]` of `mem-val` affects E2.

⇒ Only `mem-val[31 : 0]` of `mem-val` affects E0, E1, and E2.

# Claim

- Only `mem-val[23 : 0]` of `mem-val` affects E0.
- Only `mem-val[31 : 8]` of `mem-val` affects E1.
- Only `mem-val[31 : 0]` of `mem-val` affects E2.

⇒ Only `mem-val[31 : 0]` of `mem-val` affects E0, E1, and E2.

- `mem-val[31 : 0]`  
= (mod `mem-val *232*`)  
= (logand `mem-val *232-1*`)  
= ...

# Claim

- Only `mem-val[23 : 0]` of `mem-val` affects E0.
- Only `mem-val[31 : 8]` of `mem-val` affects E1.
- Only `mem-val[31 : 0]` of `mem-val` affects E2.

⇒ Only `mem-val[31 : 0]` of `mem-val` affects E0, E1, and E2.

⇒ The `main` theorem can be transformed into the `bounded` lemma by replacing `mem-val` by `mem-val[31 : 0]` in the `main` theorem.

# Bounded Main-2 Lemma

```
(defthm main-2
  (let ((mem-val (mod mem-val *2^32*)))
    (implies (and (natp i02)
                  (< i02 3))
              (equal (logior (mod (ash mem-val (* -8 i02))
                                   *2^8*)
                              (* *2^8*
                                  (mod (ash mem-val (+ -8 (* -8 i02)))
                                       *2^8*)))
                    (mod (ash mem-val (* -8 i02))
                          *2^16*))))))
```

# Bounded Main-1 Lemma

```
(def-gl-thm main-1
  :hyp (and (natp i02)
            (< i02 3)
            (n32p mem-val))
  :concl (equal (logior (mod (ash mem-val (* -8 i02))
                             *2^8*)
                        (* *2^8*
                          (mod (ash mem-val (+ -8 (* -8 i02)))
                               *2^8*))))
              (mod (ash mem-val (* -8 i02))
                   *2^16*))
  :g-bindings
  `((mem-val (:g-number ,(gl-int 0 2 33)))
     (i02 (:g-number ,(gl-int 1 2 3)))))
```

# Bounded Main-2 Lemma

```
(defthm main-2
  (let ((mem-val (mod mem-val *2^32*)))
    (implies (and (natp i02)
                  (< i02 3))
              (equal (logior (mod (ash mem-val (* -8 i02))
                                   *2^8*)
                              (* *2^8*
                                  (mod (ash mem-val (+ -8 (* -8 i02)))
                                       *2^8*)))
                    (mod (ash mem-val (* -8 i02))
                          *2^16*))))))
```



# Rewrite Rules

- $(\text{mod } (\text{ash } (\text{mod } \text{mem-val } *2^{32}*) (* -8 \text{i02}))$   
     $*2^8*)$   
=  $(\text{mod } (\text{ash } \text{mem-val } (* -8 \text{i02}))$   
     $*2^8*)$   
= E0

# Rewrite Rules

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 $*2^8*)$   
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 $*2^8*)$   
= E0
- $(\text{mod } (\text{ash } (\text{mod } \text{mem-val } *2^{32}*) (+ -8 (* -8 \text{i02}))))$   
 $*2^8*)$   
= E1

# Rewrite Rules

- $(\text{mod } (\text{ash } (\text{mod } \text{mem-val } *2^{32}*) (* -8 \text{ i02}))$   
 $*2^8*)$   
=  $(\text{mod } (\text{ash } \text{mem-val } (* -8 \text{ i02}))$   
 $*2^8*)$   
= E0
- $(\text{mod } (\text{ash } (\text{mod } \text{mem-val } *2^{32}*) (+ -8 (* -8 \text{ i02}))))$   
 $*2^8*)$   
= E1
- $(\text{mod } (\text{ash } (\text{mod } \text{mem-val } *2^{32}*) (* -8 \text{ i02}))$   
 $*2^{16}*)$   
= E2

# Main Theorem

```
(defthm main
  (let ((mem-val (mod mem-val *232*)))
    (implies (and (natp i02)
                  (< i02 3))
              (equal (logior (mod (ash mem-val (* -8 i02))
                                   *28*)
                          (* *28*
                             (mod (ash mem-val (+ -8 (* -8 i02)))
                                   *28*)))
                     (mod (ash mem-val (* -8 i02))
                           *216*)))
              :hints ((“Goal” :use (main-2))))))
```

# Agenda

- A technique for proving unbounded theorems with the help of GL.
- Benefit of using that technique in certifying the 32-bit physical memory model.

# Benefit

- The **main** theorem will help to prove **16-bit read-over-write** theorems in the 32-bit physical memory model without requiring the **(x86p x86)** hypothesis.

```
(defthm |rm-low-16 over wm-low-16 at diff-addr & non-overlap|
  (implies (and (or (< (1+ addr1) addr2)
                    (< (1+ addr2) addr1))
            (n16p val)
            (x86p x86)
            ...))
  (equal (rm-low-16 addr2 (wm-low-16 addr1 val x86))
         (rm-low-16 addr2 x86))))
```

# 16-Bit Read-Over-Write

```
(defthm |rm-low-16 over wm-low-16 at diff-addr & non-overlap|
  (implies (and (or (< (1+ addr1) addr2)
                   (< (1+ addr2) addr1))
            (n16p val)
            (x86p x86)
            ...))
  (equal (rm-low-16 addr2 (wm-low-16 addr1 val x86))
         (rm-low-16 addr2 x86))))
```

- (rm-low-<i> addr2 x86) performs reading an <i>-bit value from addr2 in x86 memory field.
- (wm-low-<j> addr1 val x86) performs writing a <j>-bit value val into x86 memory field at addr1.

# Supporting Lemma

```
(defthm rm-low-16-as-rm-low-08
  (implies (and (natp addr)
                 (< (+ 1 addr) *mem-size-in-bytes*))
            (equal (rm-low-16 addr x86)
                   (let* ((byte0 (rm-low-08 addr x86))
                          (byte1 (rm-low-08 (+ 1 addr) x86)))
                     (logior (ash byte1 8)
                              byte0))))))
```



# Key Checkpoint

```
(implies (and (natp addr)
              (< addr 4503599627370495)
              (<= (mod addr 4) 2)
              (integerp (memi (ash addr -2) x86))))
  (equal (mod (ash (memi (ash addr -2) x86)
                  (* -8 (mod addr 4)))
            65536)
         (logior (mod (ash (memi (ash addr -2) x86)
                          (* -8 (mod addr 4)))
                    256)
                 (* 256
                  (mod (ash (memi (ash addr -2) x86)
                          (+ -8 (* -8 (mod addr 4))))
                    256))))))
```

# Problem

- The key checkpoint is the **main** theorem we discussed earlier, where **i02** is replaced with **(mod addr 4)**, and **mem-val** is replaced with **(memi (ash addr -2) x86)**.
- Although **(memi (ash addr -2) x86)** returns a **32-bit** value, proving **(n32p (memi (ash addr -2) x86))** requires **(x86p x86)** hypothesis by the following lemma:

```
(defthm memi-is-n32p
  (implies (and (x86p x86)
                (natp i)
                (< i *mem-size-in-dwords*))
           (n32p (memi i x86))))
```

# Problem with (x86p x86)

- The present of (x86p x86) hypothesis in read-over-write and write-over-write theorems causes significant slowdown when proving lemmas containing read-over-long-nested-writes as well as write-over-long-nested-writes into memory.

=> The main theorem is a solution for avoiding (x86p x86) hypothesis in read-over-write theorems.

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- The present of (x86p x86) hypothesis in read-over-write and write-over-write theorems causes significant slowdown when proving lemmas containing read-over-long-nested-writes as well as write-over-long-nested-writes into memory.
- => The main theorem is a solution for avoiding (x86p x86) hypothesis in read-over-write theorems.
- How about write-over-write theorems?

# Supporting Lemma

```
(defthm wm-low-16-as-wm-low-08-lemma-1
  (implies (and (n02p i02) (< i02 3) (n16p val) (n32p mem-val))
    (equal (logior (* (mod (ash val -8) 256)
      (expt 2 (+ 8 (* 8 i02)))))
      (logand (lognot (* 255 (expt 2 (+ 8 (* 8 i02)))))
        (* (mod val 256) (expt 256 i02)))
      (logand (lognot (* 255 (expt 256 i02)))
        (lognot (* 255 (expt 2 (+ 8 (* 8 i02)))))
        mem-val))
      (logior (* val (expt 256 i02))
        (logand (lognot (* 65535 (expt 256 i02)))
          mem-val))))))
```

# Problem

- We cannot transform `wm-low-16-as-wm-low-08-lemma-1` into a `bounded` lemma because the following condition is `not` satisfied:
  - Only `fixed finite bits` of `unbounded` variables affect the computation.

# Supporting Lemma

```
(defthm wm-low-16-as-wm-low-08-lemma-1
  (implies (and (n02p i02) (< i02 3) (n16p val) (n32p mem-val))
    (equal (logior (* (mod (ash val -8) 256)
      (expt 2 (+ 8 (* 8 i02)))))
      (logand (lognot (* 255 (expt 2 (+ 8 (* 8 i02)))))
        (* (mod val 256) (expt 256 i02)))
      (logand (lognot (* 255 (expt 256 i02)))
        (lognot (* 255 (expt 2 (+ 8 (* 8 i02)))))
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```

# Strategy

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    (equal (logior (* (mod (ash val -8) 256)
      (expt 2 (+ 8 (* 8 i02))))
      (logand (lognot (* 255 (expt 2 (+ 8 (* 8 i02)))))
        (* (mod val 256) (expt 256 i02)))
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```



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```

# Timing Results

The experiments below were performed on “eld” using /projects/acl2/svn-recent/ccl-saved\_acl2hp

Certification time	8-bit	32-bit	32-bit (x86p x86)
Lemma <b>loop-effects</b>	29.90s	<b>32.83s</b>	498.17s
Lemma <b>prime-effects</b>	20.80s	<b>22.84s</b>	475.02s
Whole model	32:32.29s	<b>34:35.85s</b>	43:42.12s

Lemma **loop-effects** and **prime-effects** contain 8-bit read-over-**80**-nested-writes.

Questions!