A Self-Timed Radix-2 FFT Design

Mertcan Temel

mert@utexas.edu

February 6, 2018

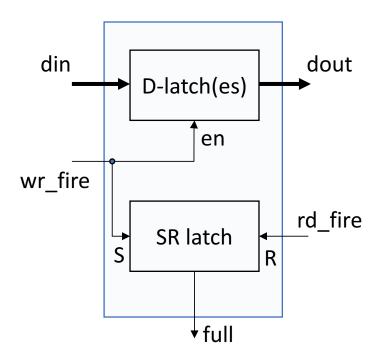
Introduction

- A design methodology/State machine examples using the Link-joint model (introduced in 2015 by Roncken et al.)
 - Unsigned Multiplier
 - Signed Multiplier
 - Complex Multiplier
 - Radix-2 FFT
- Everything implemented as circuit generators in DE (but no proofs just yet)

Outline

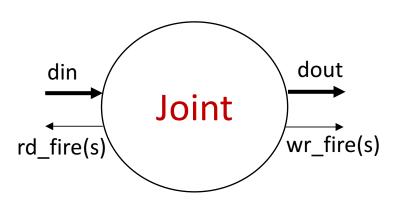
- Link-Joint Model
- Asynchronous Register
- Multipliers
 - Unsigned
 - Signed
 - Complex
- Radix-2 Decimation-in-time FFT Summary
- Self-timed Radix-2 FFT Module
- Summary and Future Work

Link-Joint Model



Links

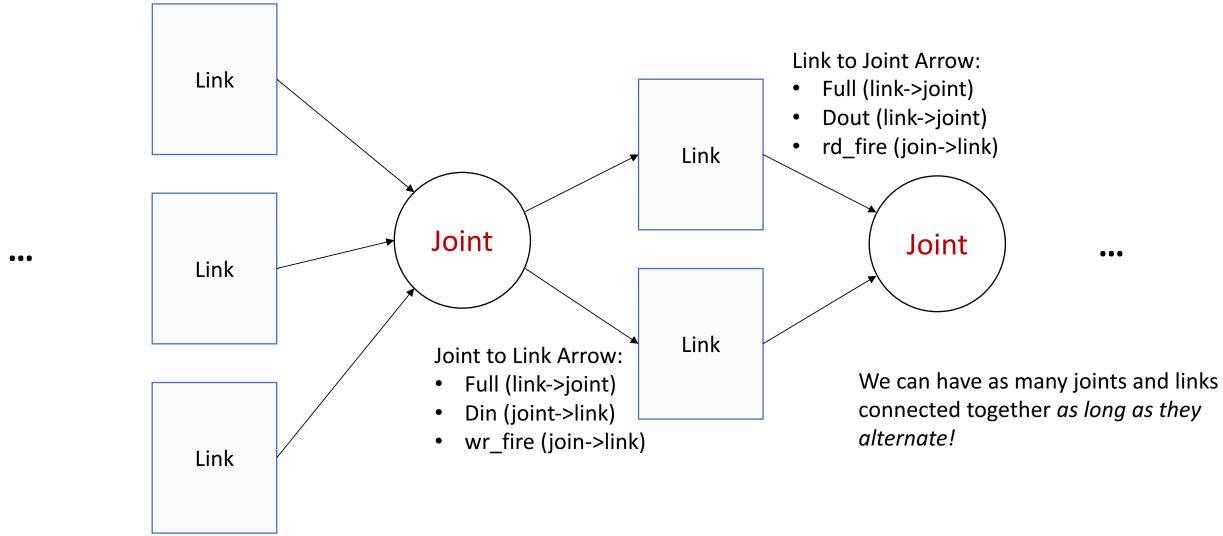
- Storage elements
- D-latches for Data
- SR latch for full status: is data valid?
- wr_fire: write new data & mark as full
- rd_fire: read data & mark as empty



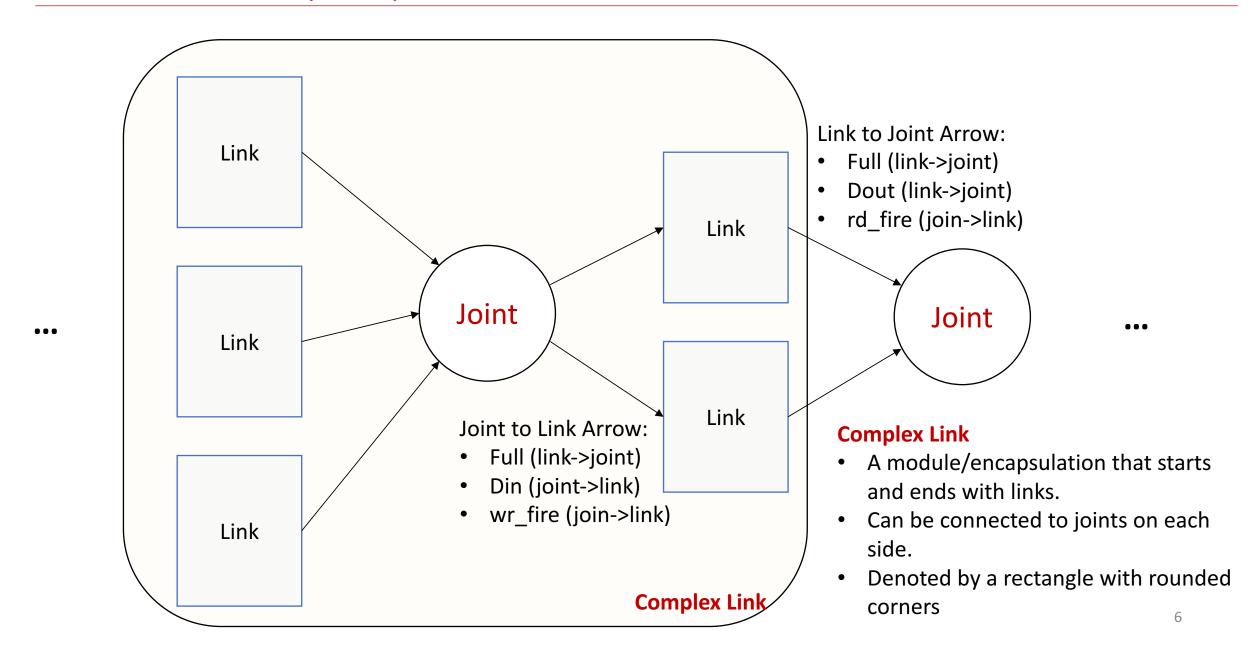
Joints

- Data Processing Elements
- Implements fire rules: decides when data should proceed
- Goes between two links
- wr_fire: write to links
- rd_fire: read from links

Link-Joint Model (cntd.)



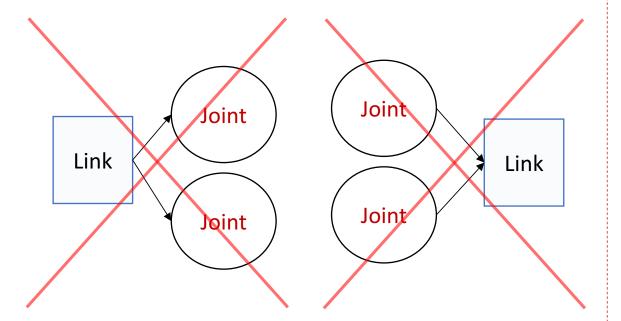
Link-Joint Model (cntd.)



Link-Joint Model Restrictions

Restriction 1:

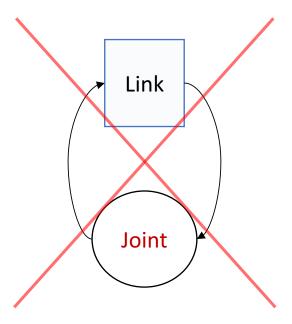
A link cannot have multiple writer/reader joints



How do you intervene and load data to modules?

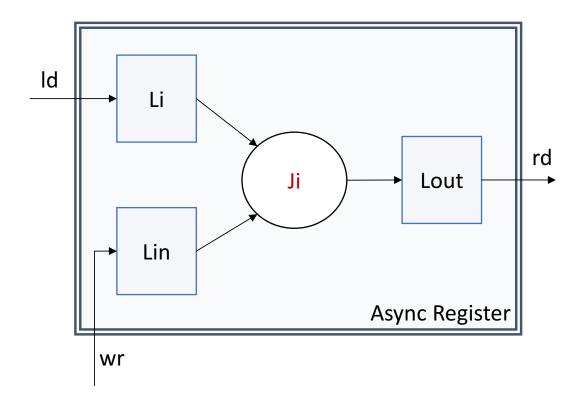
Restriction 2:

A link cannot be read and written by the same joint



How do you update the same data when processing?

Asynchronous Register



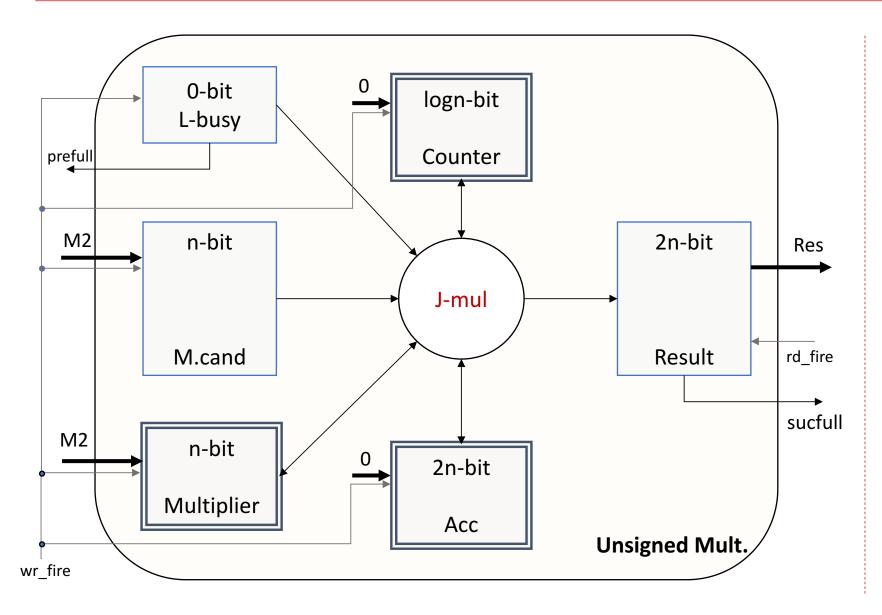
Here, we propose an asynchronous register module.

- Load initial data to Li
- Write data to Lin (during processing)
- Read data from Lout (during processing)
- Denote the module with double lined rectangles

Function of Ji:

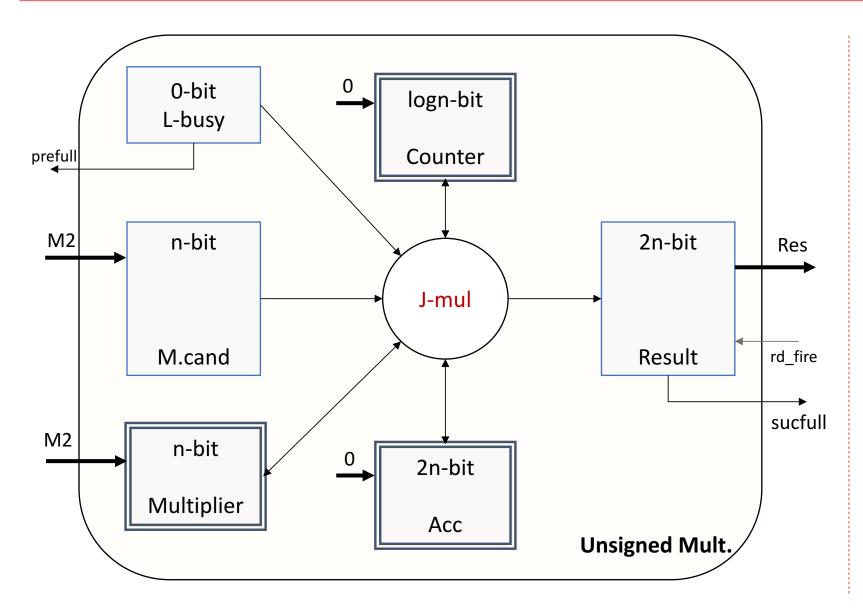
```
if (full(Li))
  Lout = Li;
  fire Li, Lin, Lout;
elsif (full(Lin) && empty(Lout))
  Lout = Lin;
  fire Lin, Lout;
end if;
```

- Now, we can load initial data, and use the same joint to update data
- We can write and still read the previous data

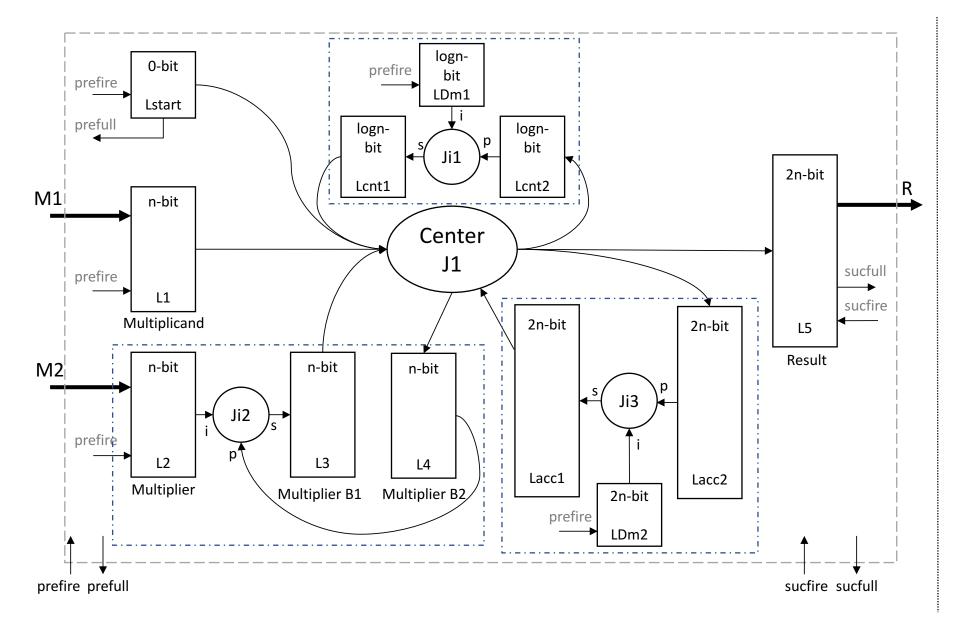


- Registers are initialized by some external joint writing to the module
- L-busy indicates the module is not ready (prefull) for new input
- Registers, M.cand, and L-busy all use the same wr_fire signal that originates externally.
- The module behaves like a link (complex link), only it has two full signals
- Double arrow indicates that the joint both reads and writes.

Asynchronous n-bit Unsigned Multiplier (cntd.)

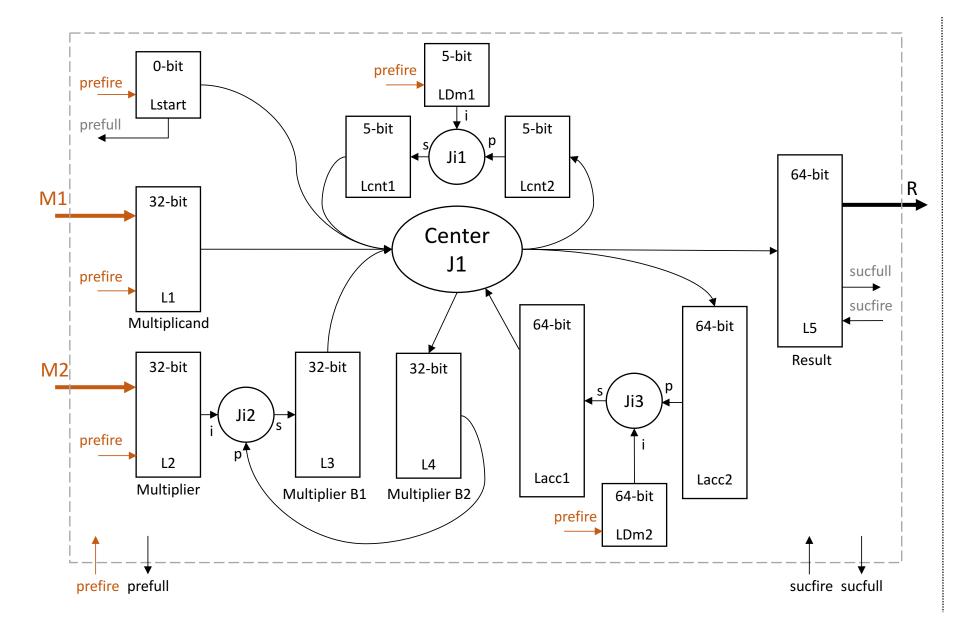


- J-mul performs shift and add to multiply.
 - if MSB(Multiplier) is 1, add
 M.cand to Acc
 - Shift Multiplier and Acc
 - When Counter=n, clear everything and buffer Acc to Result.
- M.cand is just a link as we only read from it.
- The module can pipeline requests. Can start another calculation before Result is cleared.



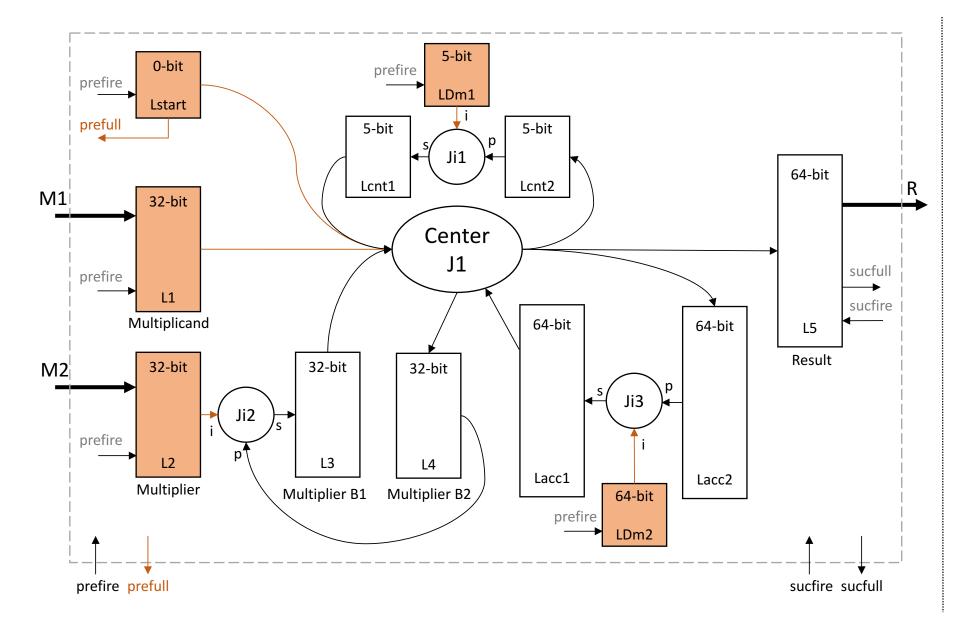
- 1. An external joint prefires to start
- 2. Ji1, Ji2, Ji3 initialize successor links
 - Lcnt1=0
 - Lacc1=0
 - L3=L2
- 3. Center Joint J1 processes and propagates data
 - If MSB(L3) then Ishf and add L1 else

- 4. Repeat when cnt<31
- 5. J1 finishes if cnt==31.
 - Wr. res to L5
 - Release all prelinks



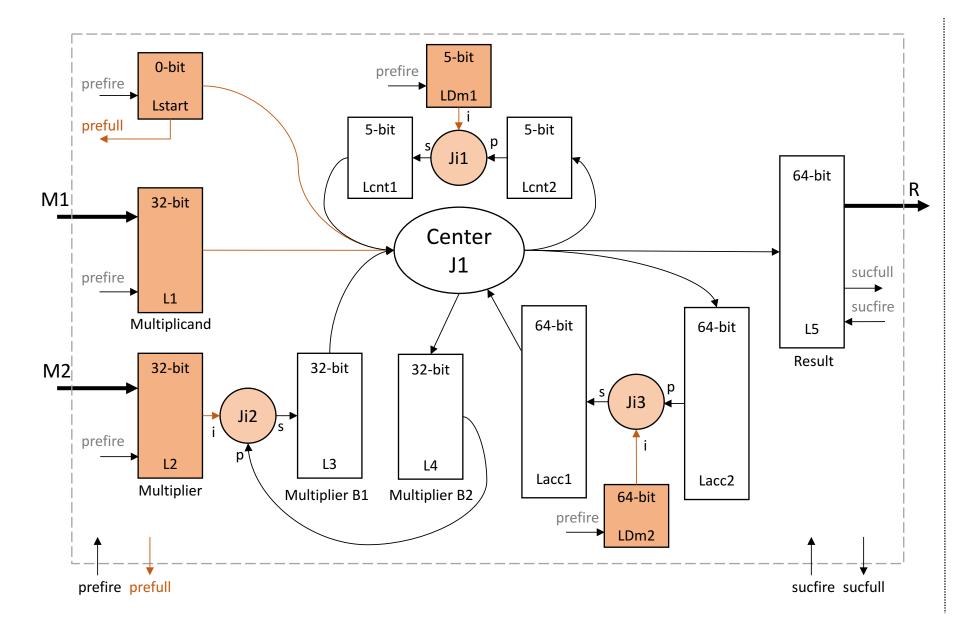
- An external joint prefires to start
- 2. Ji1, Ji2, Ji3 initialize successor links
 - Lcnt1=0
 - Lacc1=0
 - L3=L2
- 3. Center Joint J1 processes and propagates data
 - If MSB(L3) then Ishf and add L1 else

- 4. Repeat when cnt<31
- 5. J1 finishes if cnt==31.
 - Wr. res to L5
 - Release all prelinks



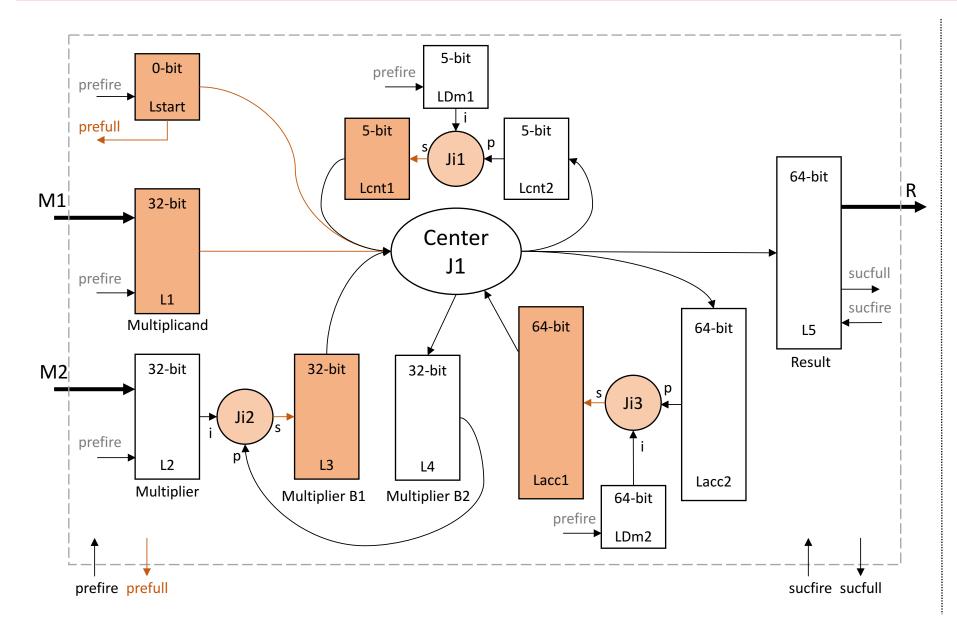
- An external joint prefires to start
- 2. Ji1, Ji2, Ji3 initialize successor links
 - Lcnt1=0
 - Lacc1=0
 - L3=L2
- 3. Center Joint J1 processes and propagates data
 - If MSB(L3) then Ishf and add L1 else

- 4. Repeat when cnt<31
- J1 finishes if cnt==31.
 - Wr. res to L5
 - Release all prelinks



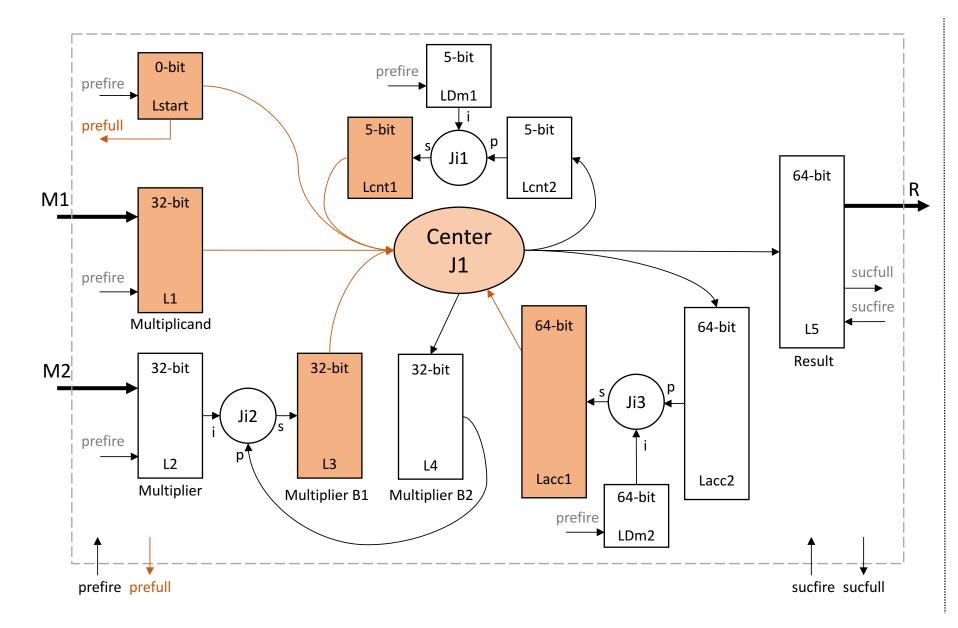
- 1. An external joint prefires to start
- 2. Ji1, Ji2, Ji3 initialize successor links
 - Lcnt1=0
 - Lacc1=0
 - L3=L2
- Center Joint J1 processes and propagates data
 - If MSB(L3) then Ishf and add L1 else

- 4. Repeat when cnt<31
- 5. J1 finishes if cnt==31.
 - Wr. res to L5
 - Release all prelinks



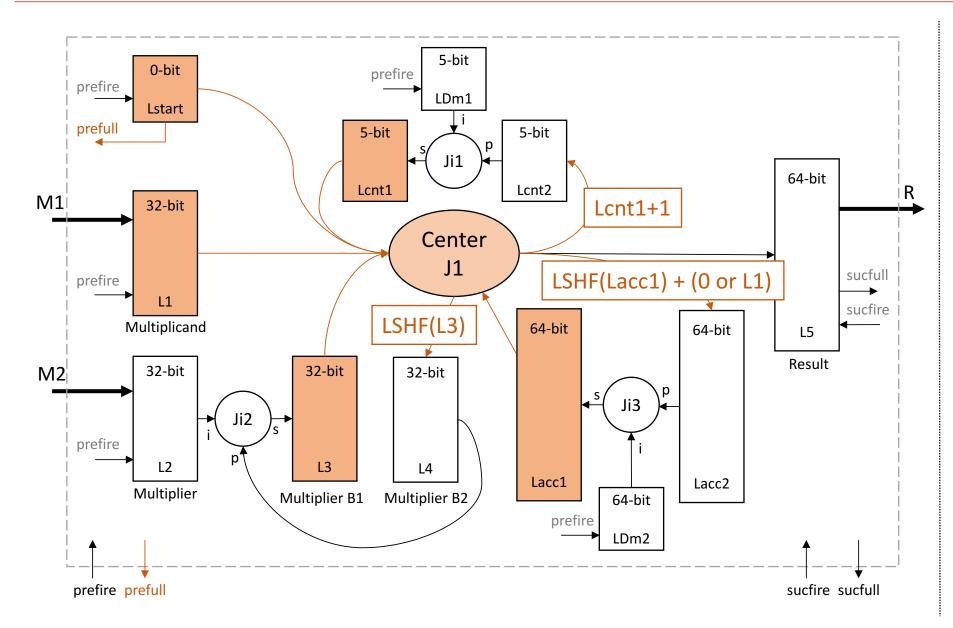
- 1. An external joint prefires to start
- 2. Ji1, Ji2, Ji3 initialize successor links
 - Lcnt1=0
 - Lacc1=0
 - L3=L2
- Center Joint J1 processes and propagates data
 - If MSB(L3) then Ishf and add L1 else

- 4. Repeat when cnt<31
- J1 finishes if cnt==31.
 - Wr. res to L5
 - Release all prelinks



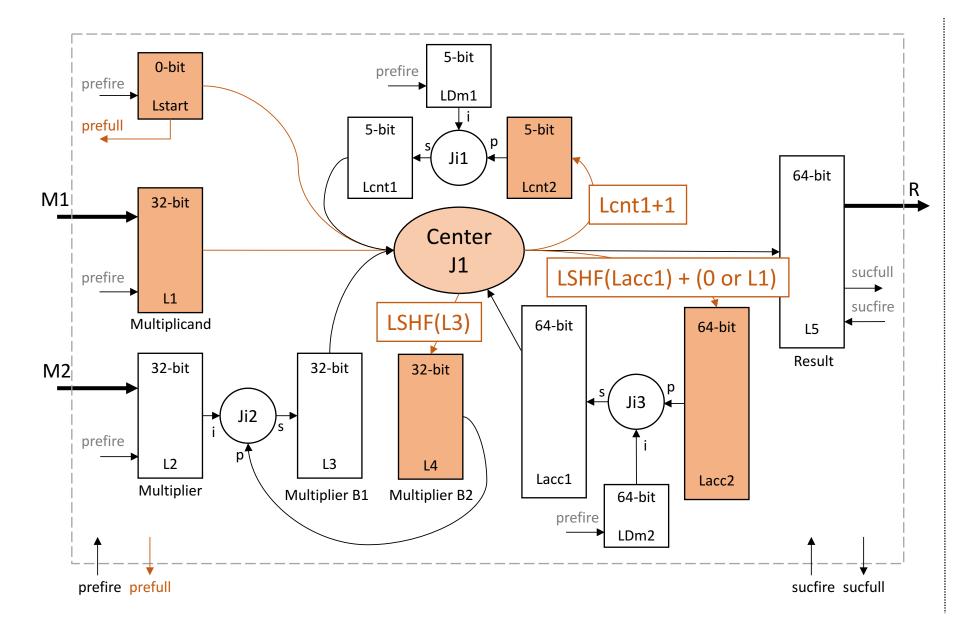
- 1. An external joint prefires to start
- 2. Ji1, Ji2, Ji3 initialize successor links
 - Lcnt1=0
 - Lacc1=0
 - L3=L2
- Center Joint J1 processes and propagates data
 - If MSB(L3) then Ishf and add L1 else

- 4. Repeat when cnt<31
- 5. J1 finishes if cnt==31.
 - Wr. res to L5
 - Release all prelinks



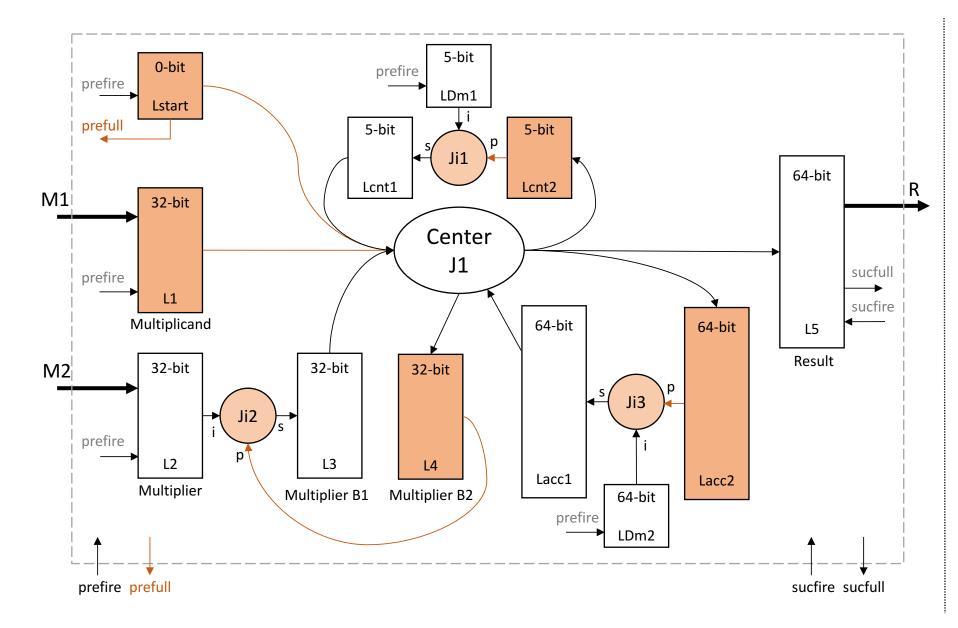
- 1. An external joint prefires to start
- 2. Ji1, Ji2, Ji3 initialize successor links
 - Lcnt1=0
 - Lacc1=0
 - L3=L2
- 3. Center Joint J1 processes and propagates data
 - If MSB(L3) then Ishf and add L1 else

- 4. Repeat when cnt<31
- 5. J1 finishes if cnt==31.
 - Wr. res to L5
 - Release all prelinks

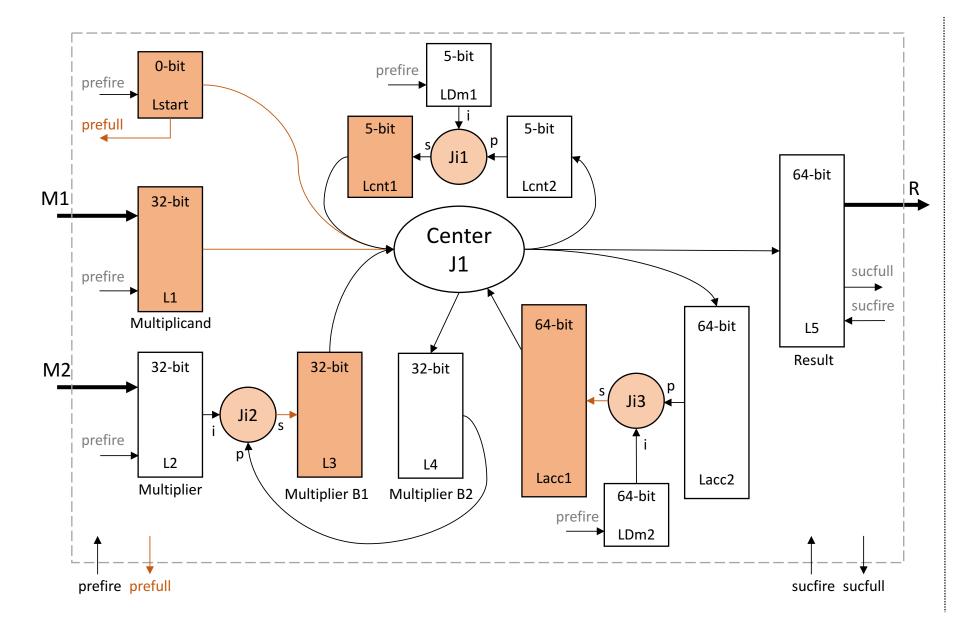


- 1. An external joint prefires to start
- 2. Ji1, Ji2, Ji3 initialize successor links
 - Lcnt1=0
 - Lacc1=0
 - L3=L2
- 3. Center Joint J1 processes and propagates data
 - If MSB(L3) then Ishf and add L1 else

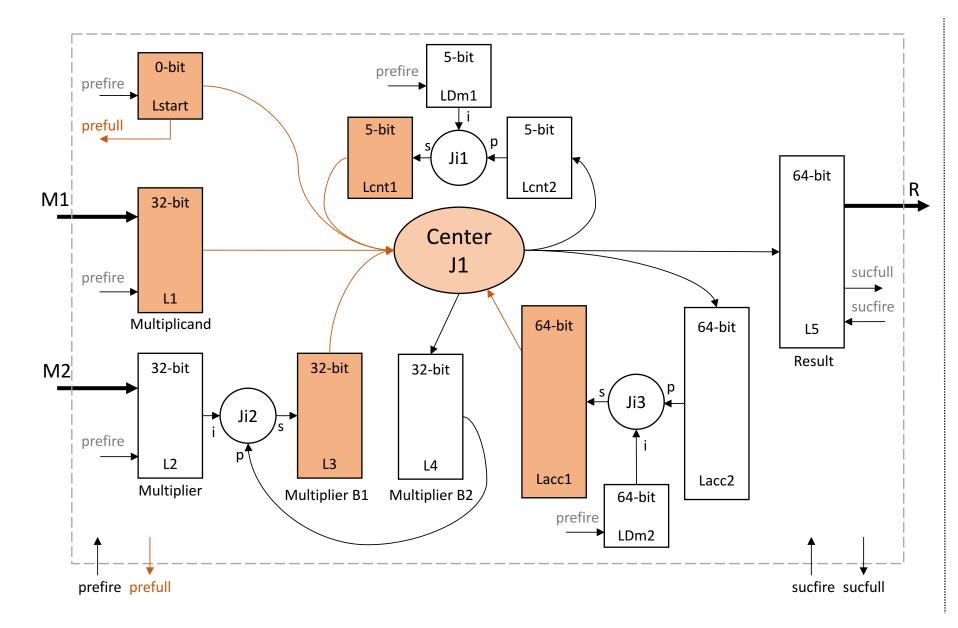
- 4. Repeat when cnt<31
- 5. J1 finishes if cnt==31.
 - Wr. res to L5
 - Release all prelinks



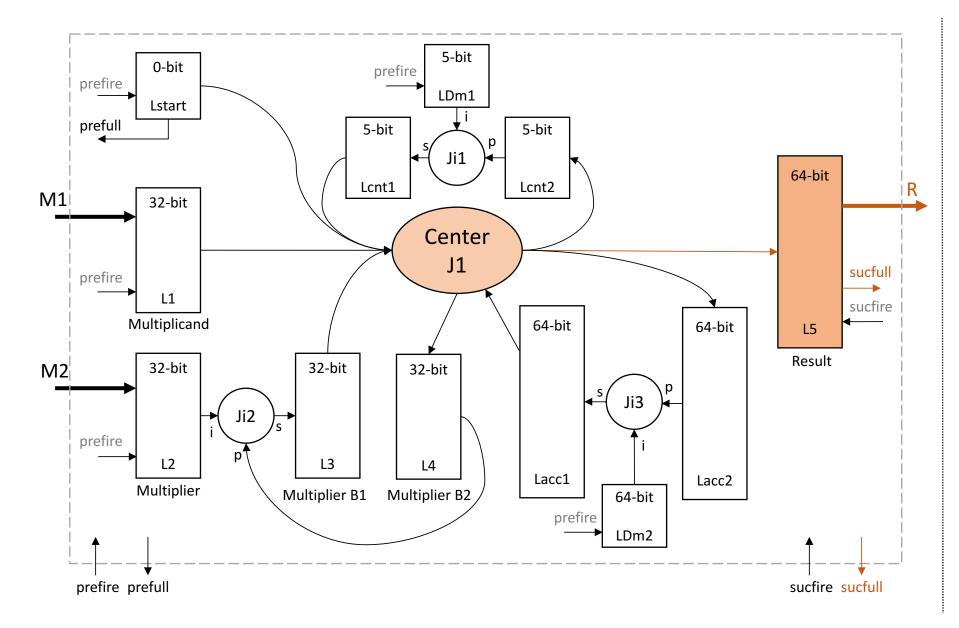
- 1. An external joint prefires to start
- 2. Ji1, Ji2, Ji3 initialize successor links
 - Lcnt1=0
 - Lacc1=0
 - L3=L2
- 3. Center Joint J1 processes and propagates data
 - If MSB(L3) then Ishf and add L1 else Ishf and buffer
- 4. Repeat when cnt<31
- 5. J1 finishes if cnt==31.
 - Wr. res to L5
 - Release all prelinks



- 1. An external joint prefires to start
- 2. Ji1, Ji2, Ji3 initialize successor links
 - Lcnt1=0
 - Lacc1=0
 - L3=L2
- 3. Center Joint J1 processes and propagates data
 - If MSB(L3) then Ishf and add L1 else Ishf and buffer
- 4. Repeat when cnt<31
- 5. J1 finishes if cnt==31.
 - Wr. res to L5
 - Release all prelinks

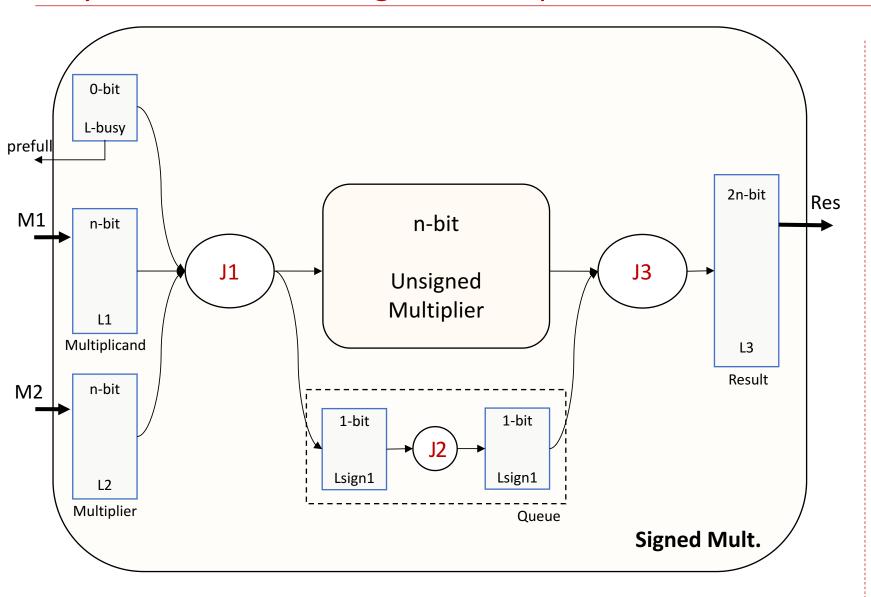


- 1. An external joint prefires to start
- 2. Ji1, Ji2, Ji3 initialize successor links
 - Lcnt1=0
 - Lacc1=0
 - L3=L2
- 3. Center Joint J1 processes and propagates data
 - If MSB(L3) then Ishf and add L1 else Ishf and buffer
- 4. Repeat when cnt<31
- 5. J1 finishes if cnt==31.
 - Wr. res to L5
 - Release all prelinks



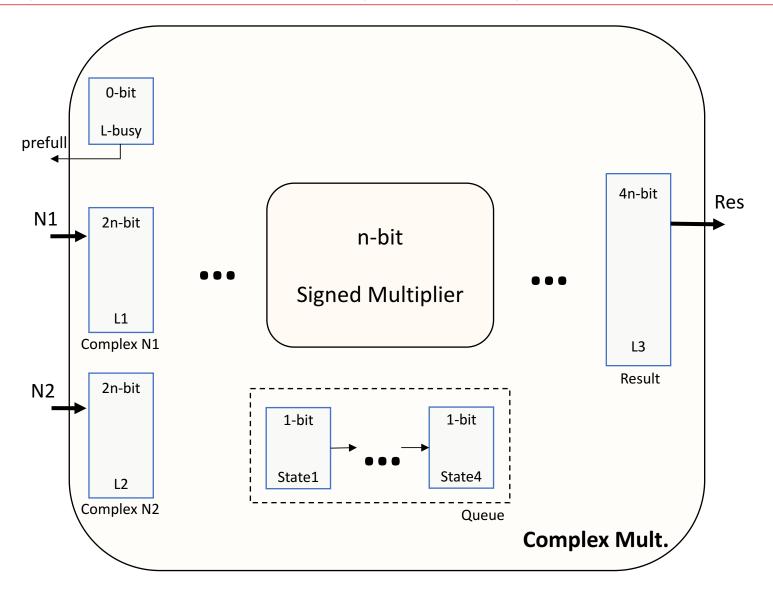
- 1. An external joint prefires to start
- 2. Ji1, Ji2, Ji3 initialize successor links
 - Lcnt1=0
 - Lacc1=0
 - L3=L2
- 3. Center Joint J1 processes and propagates data
 - If MSB(L3) then Ishf and add L1 else

- 4. Repeat when cnt<31
- 5. J1 finishes if cnt==31.
 - Wr. res to L5
 - Release all prelinks



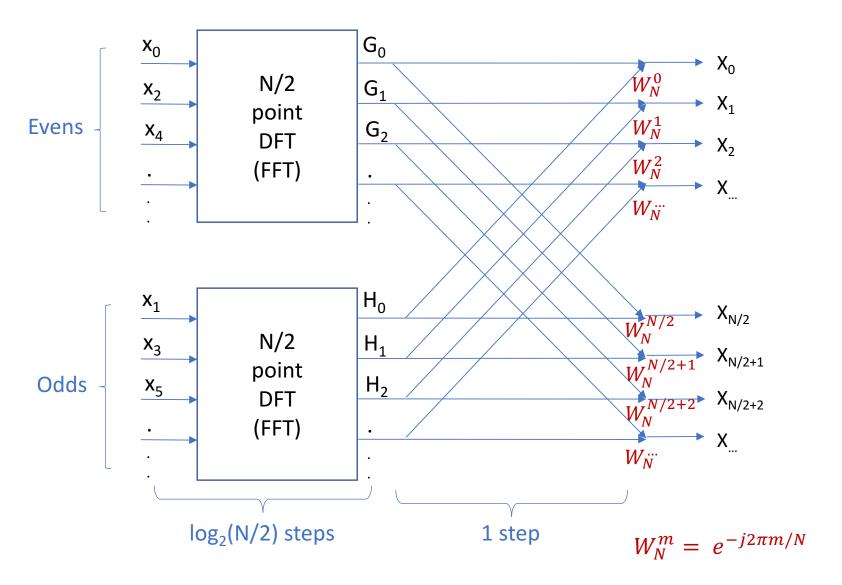
- Unsigned Multiplier is used as a link
- A queue is used to make use of pipelining capability of Unsigned Multiplier.
- J1 two's complement inputs if they're negative
- J2 buffers
- J3 two's complement the result if the result is to be negative.
- Signed Multiplier can pipeline 4
 requests before Result is
 cleared

Asynchronous n-bit Complex Multiplier



- **Signed Multiplier** is used as a link.
- A queue is used to make use of pipelining capability of Signed Multiplier.
- Performs 4 real multiplication per a complex multiplication.
- It can **pipeline 7 requests** before *Result* is cleared

Radix-2 Decimation-in-Time FFT Summary



- Goal is to calculate Discrete Fourier Transform (DFT) efficiently
- A **recursive** algorithm
 - Input divided into two sets and N/2-point DFT is calculated on each
 - Results are paired up and multiplied by a constant W_N^m
- There are log(N) steps in each of which N complex multiplications are performed.

Radix-2 Decimation-in-Time FFT Summary (cntd.)

When initializing:

```
next_i = inputs_{(reverse-bits i)}
```

When calculating:

```
if i>y
next_i = prev_y + prev_i * W_N^z
else
next_i = prev_i + prev_y * W_N^z
```

where

```
i = index of number being processed [0 N)

s = current step [0 logN)

y = i's sth bit flipped

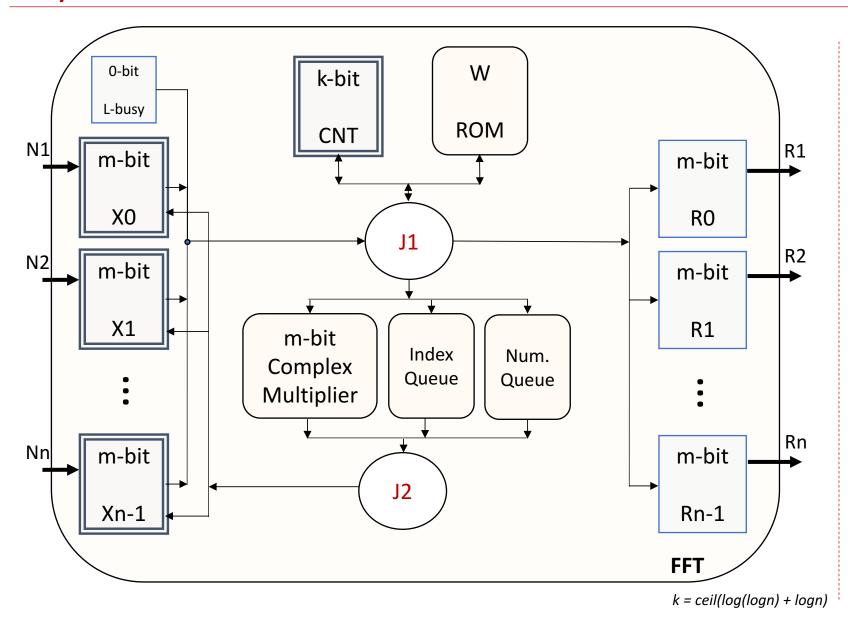
z = i << logN-s
```

Old values (prev) need to be retained as they're used twice!

The iterative solution is needed before implementing

- An formula to redistribute inputs for initialization
- Iterate over each step and then each number
- Find the **formula** to determine the **constant** W_N^m
- Find the formula to determine the other number n2
- Update numbers with n1*
 W_N^m +n2

Asynchronous Radix-2 FFT



- For N-point FFT, numbers are stored in N registers.
- J1
 - Keeps track of state
 - Selects number pair
 - Selects what W to read
- J2
 - Performs addition
 - Writes the resulting number to a register
- Index Q.: index of the number processed
- Number Q. : pair of the number processed
- Old values are kept until next step while J2 writes on the registers!

Summary and Future Work

- An asynchronous "register" module is proposed.
 - Hopefully, it'll be used developing even more complex machines
- A radix-2 FFT module with a single multiplier implemented as circuit generators in DE

- Plan to work on the functional correctness of the given modules
- Plan to introduce new designs using the link-joint model