

Stateman: Using Metafunctions to Manage Large Terms Representing Machine States

J Strother Moore
Department of Computer Science
University of Texas at Austin

Terms Representing Machine States

```
(!I 6 ; set pc to 6
(!S NIL ; set status flag to NIL
(!R 0 8 4280 ; set mem[0..7] to 4280
(!R 16 8 (LOGAND x y)
(!R (+ 40 (* 8 i)) 8 111
(!R 4280 8 999
ST))))))
```

Terms Representing Machine States

	<i>addr</i>	<i>n</i>	<i>val</i>
;			
(!I			6
(!S			NIL
(!R	0	8	4280
(!R	16	8	(LOGAND <i>x y</i>)
(!R	(+ 40 (* 8 <i>i</i>))	8	111
(!R	4280	8	999
ST))))))			

Terms Representing Machine States

```
(I
  (!I                                     6
    (!S                                   NIL
      (!R      0                          8      4280
        (!R      16                        8      (LOGAND x y)
          (!R      (+ 40 (* 8 i))        8      111
            (!R      4280                  8      999
              ST))))))
=
6
```

Terms Representing Machine States

(S

(!I			6
(!S			NIL
(!R	0	8	4280
(!R	16	8	(LOGAND <i>x y</i>)
(!R	(+ 40 (* 8 <i>i</i>))	8	111
(!R	4280	8	999
	ST))))))		

=

NIL

Terms Representing Machine States

```
(R 16 8
 (!I                                     6
  (!S                                     NIL
   (!R 0 8 4280
    (!R 16 8 (LOGAND x y)
     (!R (+ 40 (* 8 i)) 8 111
      (!R 4280 8 999
       ST))))))
=
(MOD (LOGAND x y) 264)
```

Terms Representing Machine States

```
(R 4280 8
  (!I                                     6
    (!S                                   NIL
      (!R 0 8 4280
        (!R 16 8 (LOGAND x y)
          (!R (+ 40 (* 8 i)) 8 111
            (!R 4280 8 999
              ST))))))
```

=

999

provided $(+ 40 (* 8 i) 8) \leq 4280$

$\vee (+ 4280 8) \leq (+ 40 (* 8 i))$

Rewrite Rules

- $(I (!S v st)) = (I st)$
- $((NATP a) \wedge (NATP b) \wedge (+ b k) \leq a)$
 $\rightarrow (R a n (!R b k v st)) = (R a n st)$
- $((NATP a) \wedge (NATP b) \wedge (+ a n) \leq b)$
 $\rightarrow (R a n (!R b k v st)) = (R a n st)$

Such rules suffice to manipulate state expressions.

Rewrite Rules

- $(I (!S v st)) = (I st)$
- $((NATP a) \wedge (NATP b) \wedge (+ b k) \leq a)$
 $\rightarrow (R a n (!R b k v st)) = (R a n st)$
- $((NATP a) \wedge (NATP b) \wedge (+ a n) \leq b)$
 $\rightarrow (R a n (!R b k v st)) = (R a n st)$

Such rules suffice to manipulate state expressions — **except** when there are deep nests of **!R**-expressions and a , b , n , and k are large expressions.

Terms Representing Machine States

```
(!I                                     6
  (!S                                    NIL
    (!R      0                          8      4280
      (!R      16                        8      (LOGAND x y)
        (!R      (+ 40 (* 8 i))          8      111
          (!R      4280                   8      999
            ST))))))
```

Size (in function applications): 9

Biggest Address or Value Expression: 2

Motivation for This Project

We have recently analyzed a piece of code (15,361 instructions of a formal ISA) involving states with:

Size (in function applications): 2,158,895

Biggest Address or Value Expression: 147,233

Backchaining to decide questions like

$$(+\ 40\ (*\ 8\ i)\ 8) \leq 4280$$

$$\vee (+\ 4280\ 8) \leq (+\ 40\ (*\ 8\ i))$$

for every pair of addresses in such state expressions is prohibitive.

Highlights

- Manage read-over-write and write-over-write expressions exclusively with metafunctions
- Implement a syntactic interval inference mechanism
- Implement syntactic means of deciding some inequalities
- Implement syntactic means of simplifying some MOD expressions
- Use syntactic means to decide overlap questions
- Insist that all byte counts be quoted constants

- Do not put nested !R-expressions into address order
- Eliminate perfectly shadowed writes
- Use `hons` rather than `cons` to create state expressions
- HIDE the state expressions produced by the metafunctions
- HIDE some values extracted by reads from hidden states to avoid re-simplifying them
- Prove guards and well-formedness guarantees of the metafunctions

Ainni — Our Interval Analyzer

Given

```
(+ 288 (* 8 (LOGAND 31 (ASH (R 4520 8 st) -3))))
```

our analyzer reports an interval of **[288, 536]**.

But if $(R\ 4520\ 8\ st) < 24$ is known by context, then the interval shrinks to **[288, 304]**.

The analyzer can compute the interval **[0, $2^{32} - 1$]** for the largest value term encountered (147,233 function applications) in 0.01 seconds.

Examples of Ainni

(switch to *shell* buffer)

Finding Assignments

(R 4280 8

(!I

6

(!S

NIL

(!R 0

8

4280

(!R 16

8

(LOGAND *x y*)

(!R (+ 40 (* 8 *i*))

8

111

(!R 4280

8

999

ST))))))

=

Finding Assignments

```
(R 4280 8
  (!S                                     NIL
    (!R      0      8      4280
      (!R      16      8      (LOGAND x y)
        (!R      (+ 40 (* 8 i)) 8      111
          (!R 4280      8      999
            ST))))))
=
```

Finding Assignments

```
(R 4280 8 ; [4280,4287] vs [0,7]
  (!R 0 8 4280
    (!R 16 8 (LOGAND x y)
      (!R (+ 40 (* 8 i)) 8 111
        (!R 4280 8 999
          ST))))))
```

=

Finding Assignments

```
(R 4280 8 ; [4280,4287] vs [16,23]
  (!R 16 8 (LOGAND x y)
    (!R (+ 40 (* 8 i)) 8 111
      (!R 4280 8 999
        ST))))))
```

=

Finding Assignments

```
(R 4280 8 ; [4280,4287] vs [40,167] w/  $i < 16$ 
  (!R (+ 40 (* 8 i)) 8 111
    (!R 4280 8 999
      ST))))))
```

=

Finding Assignments

(R 4280 8 ; same!

(!R 4280 8 999
ST))))))

=

999

Preliminary Performance Results

A: guard verification

B: well-formedness

C: honsing

D: memoization

— 988 secs

A 955 secs

A+B 618 secs

A+B+C 494 secs

A+B+C+D 375 secs

Future Work

- provide a metafunction to prove state equality
- engineer ACL2 to cope better with large definitions

More Generally

This project illustrates a very common industrial application of ACL2: as a programming language suitable for writing verified programs.

By mixing verified metafunctions with the rest of ACL2, one can build a powerful domain-specific prover.