CS378 Autonomous Multiagent Systems Spring 2005

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Week 13b: Thursday, April 21st

Good Afternoon, Colleagues

Are there any questions?

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- How did they the FCC make the decisions in the end?
 What were they?
- Is TAC ongoing? What changes have been made?
- How do people at different places co-author papers?

Logistics

• Include Dr. Rich on your response

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- Some games from last year available

Open vs. Sealed Bid

- Open increases information, reducing winner's curse
 - Leads to higher bids
- But...
 - Risk aversion leads to higher bids in sealed bid auctions
 - Sealed bid auctions deter colusion
- Decided former outweighed latter
- Went with announcing bids, but not the bidders
 - Circumvented!

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 - End auction quickly
 - Close licenses almost simultaneously
 - be simple and understandable

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Went with activity rules

Combinatorial Bids

Nationwide bidding could decrease efficiency and revenue

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- Nationwide bidding could decrease efficiency and revenue
- Full combinatorial bidding too complex
 - Winner determination problem
 - Active research area

Aiding Designated Bidders

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Royalties vs. Up-front Payments

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- Decided against

Reserve Prices

- Not necessary in such a competitive market
- Did include withdrawal penalties

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 - Lots of bidders
 - Lots of revenue

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- Incremental design changes
 - New problems always arise
 - Bidders indeed find ways to circumvent mechanisms
- Lessons to be learned via agent-based experiments

Class Discussion

Ankit Srivastava on auctions in RoboCup Rescue

Trading Agent Competition

- Put forth as a benchmark problem for e-marketplaces (Wellman, Wurman, et al., 2000)
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 - Client: TACtown ↔ Tampa within 5-day period

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- Put forth as a benchmark problem for e-marketplaces (Wellman, Wurman, et al., 2000)
- Autonomous agents act as travel agents
 - Game: 8 agents, 12 min.
 - Agent: simulated travel agent with 8 clients
 - Client: TACtown → Tampa within 5-day period
- Auctions for flights, hotels, entertainment tickets
 - Server maintains markets, sends prices to agents
 - Agent sends bids to server over network

28 Simultaneous Auctions

Flights: Inflight days 1-4, Outflight days 2-5 (8)

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- 16 rooms per auction; 16th-price ascending auction; quote is ask price; no resale
- Random auction closes minutes 4 11

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Entertainment: Wrestling/Museum/Park days 1-4 (12)

 Continuous double auction; initial endowments; quote is bid-ask spread; resale allowed

Client Preferences and Utility

Preferences: randomly generated per client

- Ideal arrival, departure days
- Good Hotel Value
- Entertainment Values

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Score: Sum of client utilities – expenditures

Allocation

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G \equiv \text{complete} \text{ allocation of goods to clients} v(G) \equiv \text{utility of } G - \text{cost of needed goods} G^* \equiv \text{argmax } v(G)
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Prices known $\Rightarrow G^*$ known \Rightarrow optimal bids known

Learn model of expected hotel price

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Goal: analytically calculate optimal bids

Features:

- Current hotel and flight prices
- Current time in game
- Hotel closing times
- Agents in the game (when known)
- Variations of the above

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New algorithm for conditional density estimation

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- Value of *i*th copy is avg($V_i V_{i-1}$)

Other Uses of Sampling

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Entertainment: Bid more (ask less) than expected value of having one more (fewer) ticket

Finals

Team	Avg.	Adj.	Institution
ATTac	3622	4154	AT&T
livingagents	3670	4094	Living Systems (Germ.)
whitebear	3513	3931	Cornell
Urlaub01	3421	3909	Penn State
Retsina	3352	3812	CMU
CaiserSose	3074	3766	Essex (UK)
Southampton	3253*	3679	Southampton (UK)
TacsMan	2859	3338	Stanford

- ATTac improves over time
- livingagents is an open-loop strategy

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- EarlyBidder: motivated by TAC-01 entry livingagents
 - Immediately bids high for G^* (with $SimpleMean_{ns}$)
 - Goes to sleep



Stability

• 7 EarlyBidder's with 1 ATTac

Agent	Score	Utility
ATTac	2431 ± 464	8909 ± 264
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EarlyBidder gets more utility; ATTac pays less

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Agent	Relative Score		
	Phase I	Phase III	
$ATTac_{ns}$	$105.2 \pm 49.5 \ (2)$	$166.2 \pm 20.8 \ (1)$	
ATTac _s	$27.8 \pm 42.1 (3)$	$122.3 \pm 19.4 \ (2)$	
EarlyBidder	$140.3 \pm 38.6 \ (1)$	$117.0 \pm 18.0 \ (3)$	
$SimpleMean_{ns}$	$-28.8 \pm 45.1 (5)$	$-11.5 \pm 21.7 \ \ (4)$	
SimpleMean _s	$-72.0 \pm 47.5 (7)$	$-44.1 \pm 18.2 (5)$	
$Conditional Mean_{ns}$	$8.6 \pm 41.2 \ (4)$	$-60.1 \pm 19.7 (6)$	
Conditional Mean _s	$-147.5 \pm 35.6 \ (8)$	$-91.1 \pm 17.6 \ (7)$	
CurrentPrice	$-33.7 \pm 52.4 \ \ (6)$	$-198.8 \pm 26.0 \ \ (8)$	



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- Other complex rules



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It's a poker game!

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- Utility is profit $\Rightarrow \Sigma_l(value-cost)$
- modeled 5 most important bidders
 - others served mainly to raise prices (small bidders)
 - lower valuations (75% → pessimistic)

Bidding Strategies

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 - explicit communication not allowed...

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 - convention: unlucky bidders may take licences until fair
- Small bidders take licences from owners
 - big bidders remember licence's owner
 - allocate while small bidders active

RSDR vs. Knapsack

Method	Agent	Profit (\$M)		Ratio	Cost
Knapsack	1	980	(± 170)	1.00	.82
	2	650	(±85)	1.00	.82
	3	830	(±91)	1.00	.84
	4	170	(±20)	1.00	.84
	5	550	(±96)	1.00	.86
PRSDR	1	1240	(±210)	1.26	.76
	2	820	(±83)	1.25	.77
	3	1300	(±290)	1.58	.74
	4	300	(±44)	1.78	.79
	5	930	(±240)	1.68	.76

44% more profit; avg. ratio 1.51

Robustness

- What if someone cheats?
 - cheat: defect back to knapsack
 - others stay out of its way \Rightarrow big win
- Solution: Punishing RSDR (PRSDR)
 - cheaters may not own licences
 - recall: non-cheaters take licence from owner = fairing
 - convention: cheater takes your licence ⇒ take it back
 - * take it back first while still have money
 - * aggressively punitive: skips optimizers

Simplification: pointing out cheaters by hand

PRSDR Results

Method	Ratio	Cost
Knapsack	1.00	.84
(P)RSDR	1.51	.76
RSDR Cheater	1.63	.76
RSDR Victims	1.22	.79
PRSDR Cheater	1.02	.83
PRSDR Enforcers	1.17	.81

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Threats work!

Extensions

- Change small bidder valuations
 - test robustness
 - RSDR is optimal for preserving profit
- Multiple cheaters
 - current punishment too aggressive
 - collapse back to knapsack instead

Extentions

Method	Ratio	Local Ratio	Cost
Multiple Cheater	1.03		.84
Multiple Enforcer	1.01		.83

Extentions

Method	Ratio	Local Ratio	Cost
Multiple Cheater	1.03		.84
Multiple Enforcer	1.01		.83
50%, Knapsack	1.70	1.00	.74
50%, PRSDR	3.42	2.02	.51
75%, Knapsack	1.00	1.00	.84
75%, PRSDR	1.51	1.51	.76
85%, Knapsack	0.68	1.00	.89
85%, PRSDR	0.81	1.19	.87

Future Work

- Robustness enhancements
 - better punishment method
- More complex value functions
 - inter-market dependencies
- Automatic cheater detection
 - partial cheating vs. detection arms race
 - smack back into compliance
- Generalization to other auctions
 - more robust to tie-breaking procedure variations

Summary

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- Enables much higher profits
- Works even uncertain knowledge
- Real-world functionality relies on simple assumptions:

Summary

- Communication-free coordination
- Enables much higher profits
- Works even uncertain knowledge
- Real-world functionality relies on simple assumptions:
 - bidders want more profit
 - bidders familiar with PRSDR and its benefits
 - bidders willing to try it risk-free

Last-minute bidding (R,O, 2001)

- eBay: first-price, ascending auction
- Amazon: auction extended if bid in last 10 minutes
- eBay: bots exist to incrementally raise your bid to a maximum
- Still people snipe. Why?
 - There's a risk that the bid might not make it
 - However, common-value ⇒ bid conveys info
 - Late-bidding can be seen as implicit collusion
 - Or . . . , lazy, unaware, etc. (Amazon and eBay)
- Finding: more late-bidding on eBay,
 - even more on antiques rather than computers

Small design-difference matters



Late Bidding as Best Response

- Good vs. incremental bidders
 - They start bidding low, plan to respond
 - Doesn't give them time to respond
- Good vs. other snipers
 - Implicit collusion
 - Both bid low, chance that one bid doesn't get in
- Good in common-value case
 - protects information

Overall, the analysis of multiple bids supports the hypothesis that last-minute bidding arises at least in part as a response by sophisticated bidders to unsophisticated incremental bidding.