Asynchronous Programming Promises + Futures

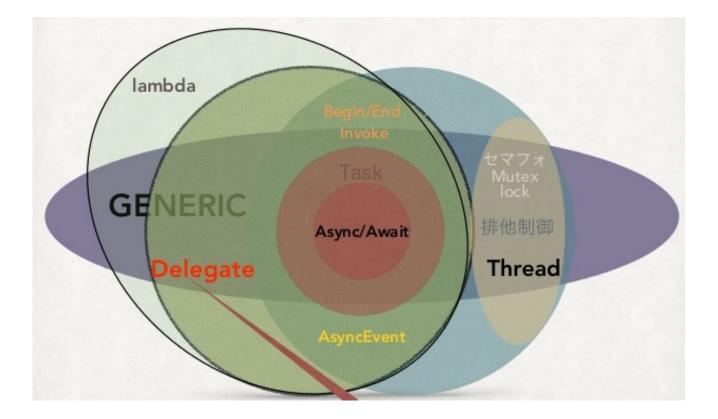
Chris Rossbach

CS378H

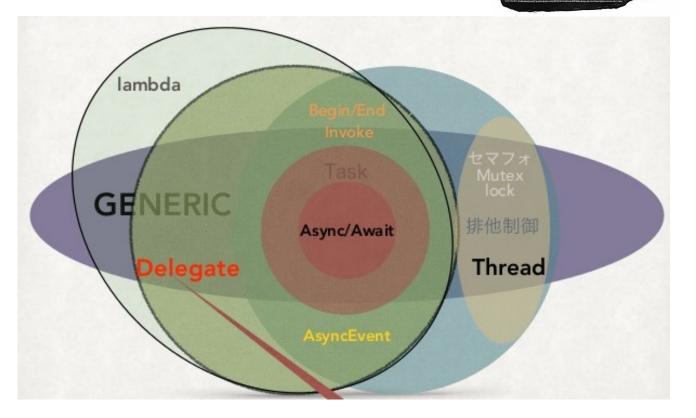
Today

- Questions?
- Administrivia
- Material for the day
 - Events / Asynchronous programming
 - Promises & Futures
 - Bonus: memory consistency models
- Acknowledgements
 - Consistency slides borrow some materials from Kevin Boos. Thanks!

Asynchronous Programming Events, Promises, and Futures



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      CompletableFuture cf = CompletableFuture.runAsync(() -> {
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 - Anonymous function
 - Functor

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- runAsync() immediately returns a waitable object (cf)
- Where (on what thread) does the lambda expression run?

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Futures:

- abstraction for concurrent work supported by
 - Compiler: abstractions are *language-level objects*
 - Runtime: scheduler, task queues, thread-pools are transparent
- Programming remains mostly imperative
 - Threads of control peppered with asynchronous/concurrent tasks

Compromise Programming Model between:

- Event-based programming
- Thread-based programming

• Where (on what thread) does the lambda expression run?

GUI Programming Distilled

```
⊟winmain(...) {
2
      while(true) {
   3
          message = GetMessage();
 4
          switch(message) {
 5
          case WM THIS: DoThis(); break;
 6
          case WM THAT: DoThat(); break;
 7
          case WM OTHERTHING: DoOtherThing(); break;
          case WM DONE: return;
8
9
10
```

GUI Programming Distilled

```
How can we
   ⊟winmain(...) {
      while(true) {
2
                                              parallelize
3
          message = GetMessage();
4
          switch(message) {
                                                 this?
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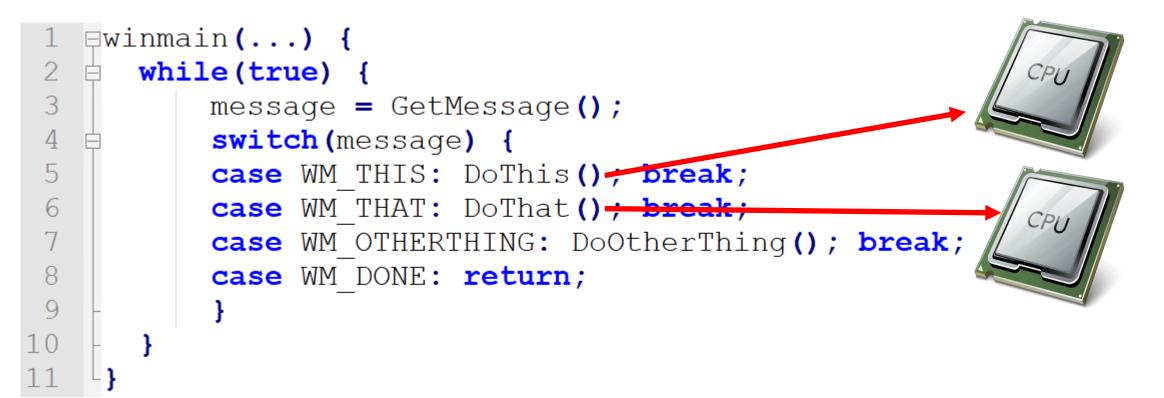


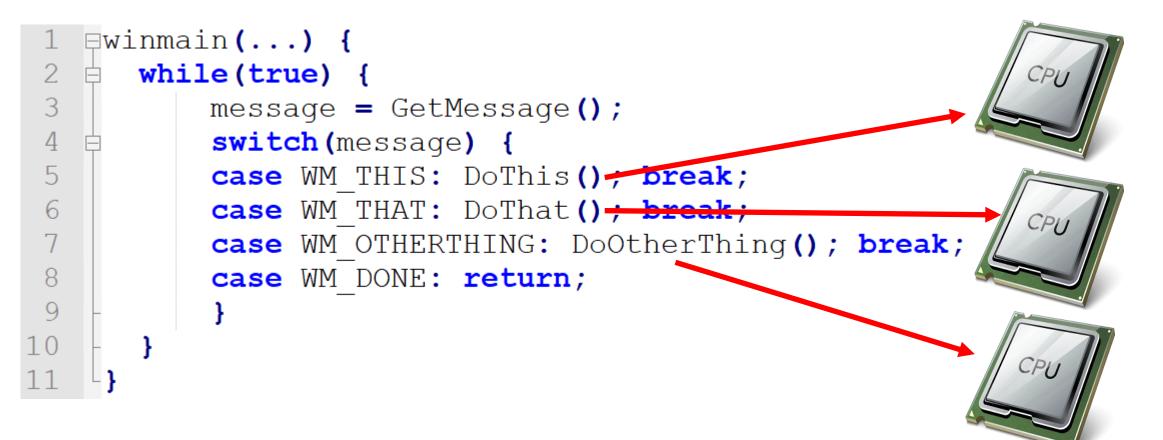




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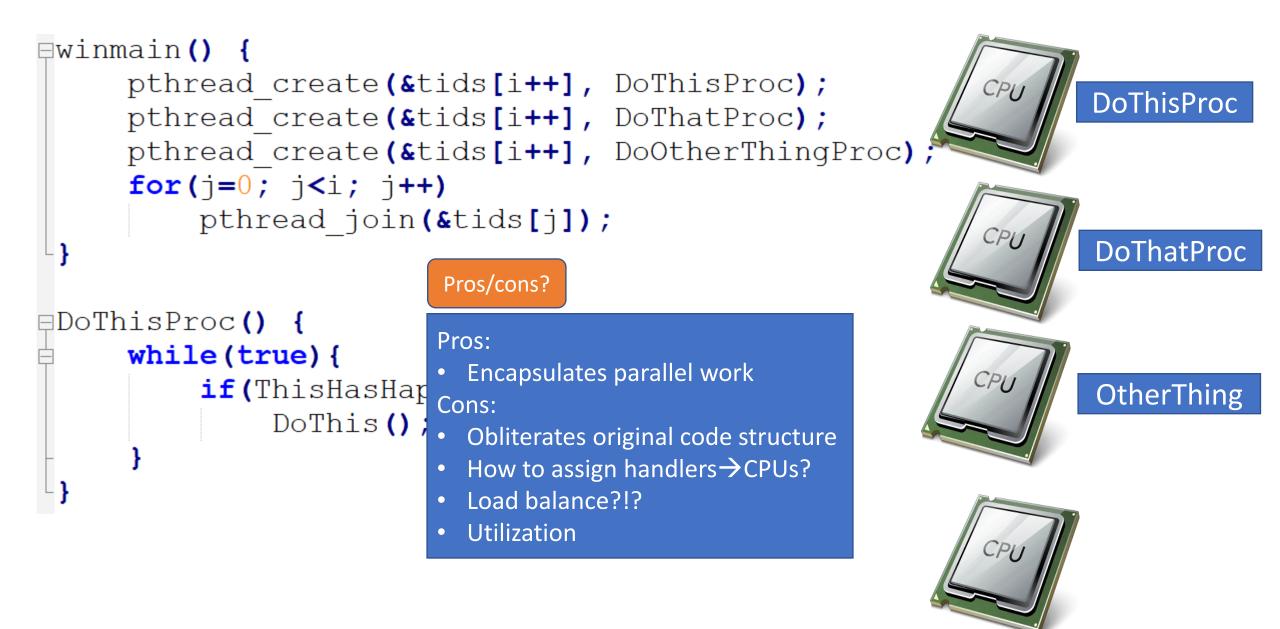
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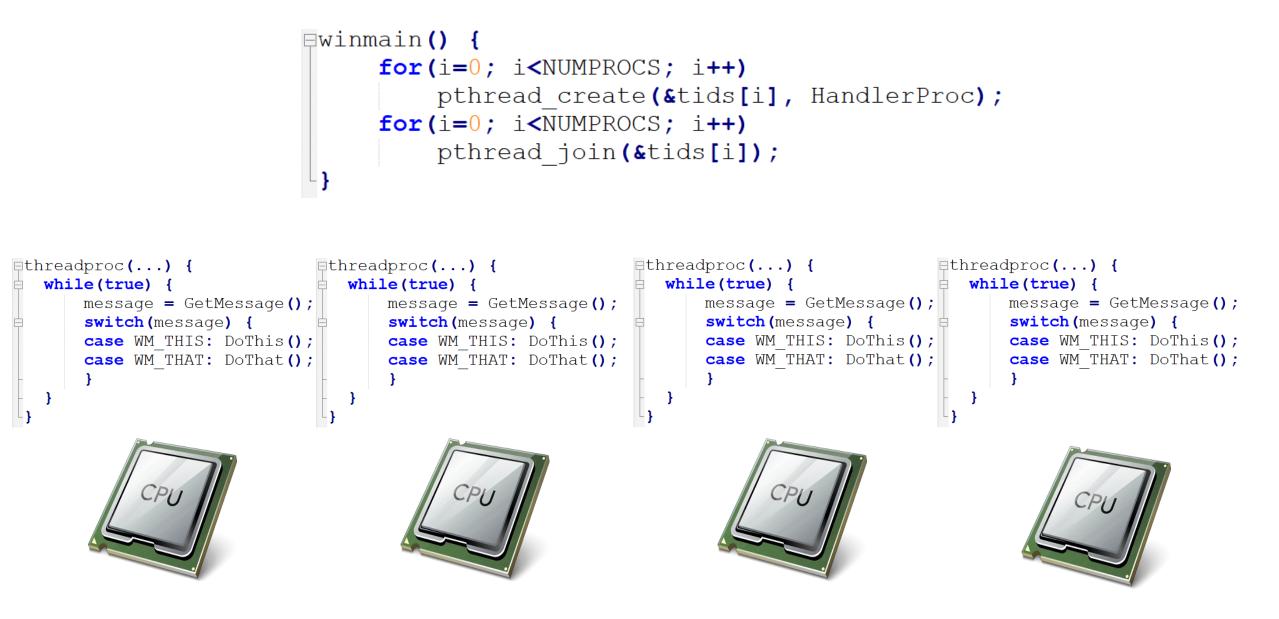




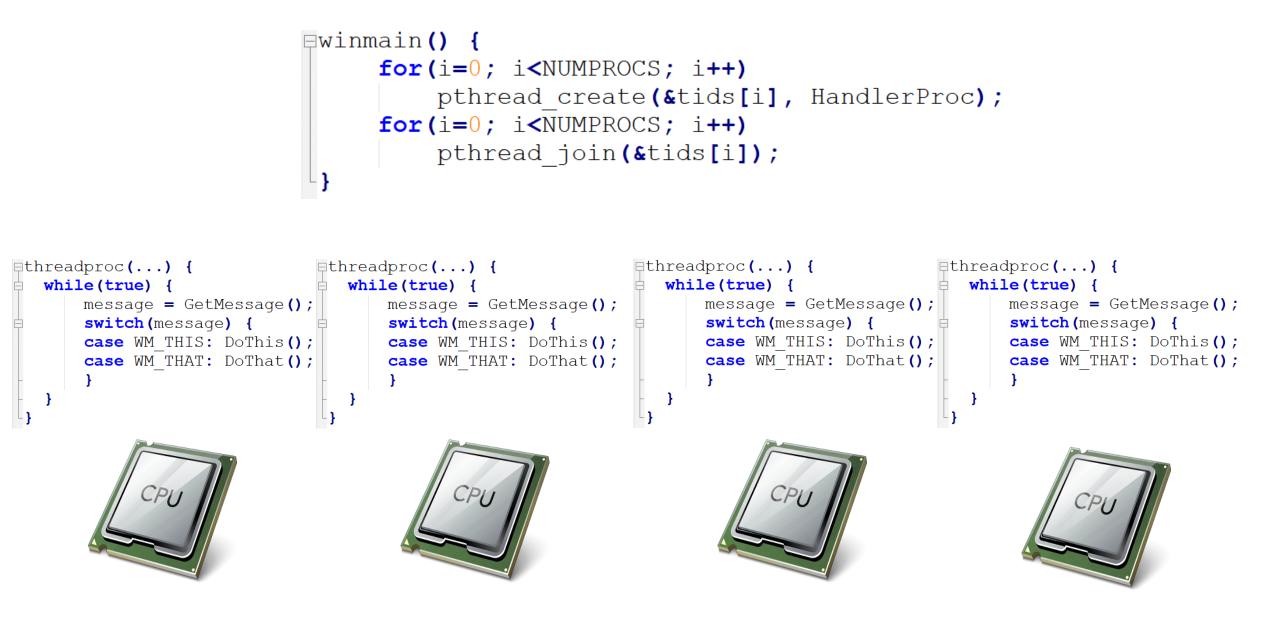
```
⊟winmain() {
     pthread create(&tids[i++], DoThisProc);
                                                               DoThisProc
     pthread create(&tids[i++], DoThatProc);
     pthread create(&tids[i++], DoOtherThingProc)
     for(j=0; j<i; j++)</pre>
         pthread join(&tids[j]);
                                                               DoThatProc
⊟DoThisProc() {
     while(true) {
         if (ThisHasHappened)
                                                               OtherThing
              DoThis();
```







Pros/cons?



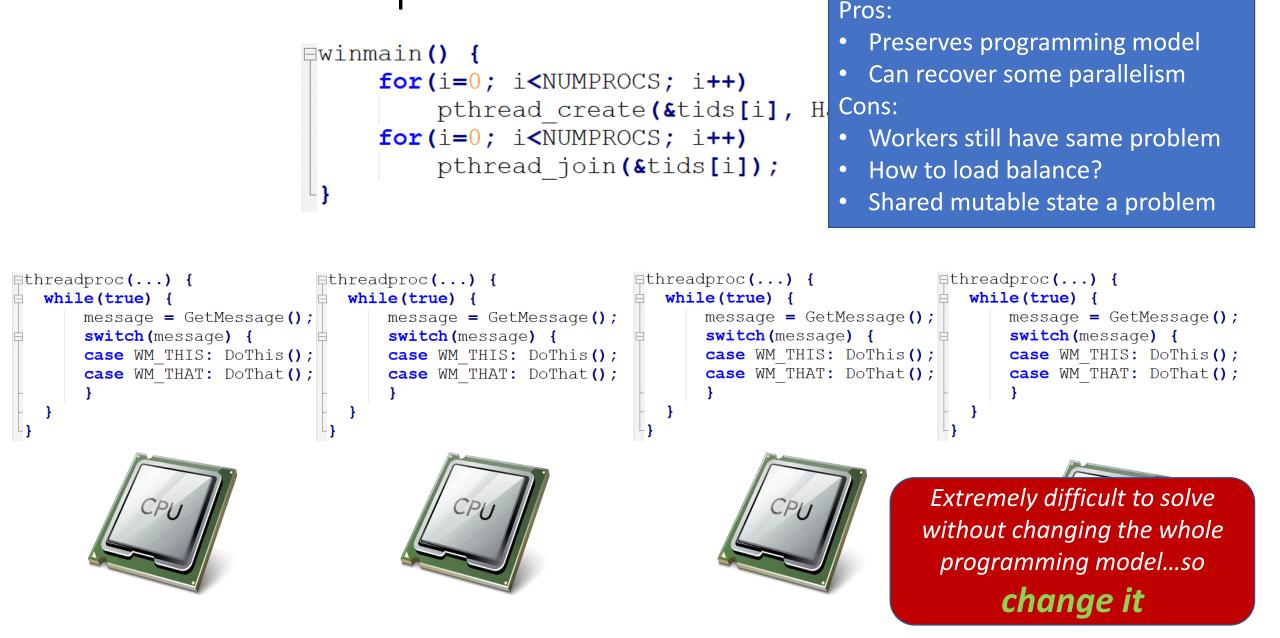
```
Preserves programming model
                          ⊟winmain() {
                                                                              Can recover some parallelism
                                 for(i=0; i<NUMPROCS; i++)</pre>
                                       pthread create (&tids[i], H. Cons:
                                 for(i=0; i<NUMPROCS; i++)</pre>
                                                                             • Workers still have same problem
                                       pthread join(&tids[i]);
                                                                               How to load balance?
                                                                             • Shared mutable state a problem

threadproc(...) {

                                                          □threadproc(...) {
                                                                                      □threadproc(...) {
threadproc(...) {
  while(true) {
                              while(true) {
                                                            while(true) {
                                                                                         while(true) {
     message = GetMessage();
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                                                                                             case WM THIS: DoThis();
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Event-based Programming: Motivation

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- Threads have a *lot* of down-sides:
 - Tuning parallelism for different environments
 - Load balancing/assignment brittle
 - Shared state requires locks \rightarrow
 - Priority inversion
 - Deadlock
 - Incorrect synchronization
 - ...

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- Events: restructure programming model so threads are not exposed!

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 - create_event_queue(handler) → event_q
 - enqueue_event(event_q, event-object)
 - Invokes handler (eventually)

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 - Invokes handler (eventually)
- Scheduler decides which event to execute next
 - E.g. based on priority, CPU usage, etc.

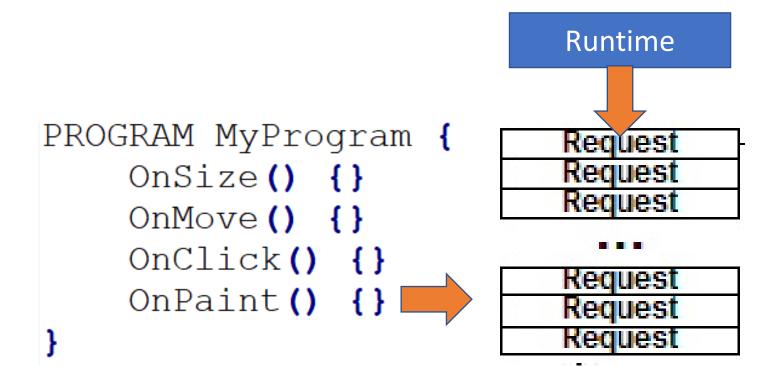
Event-based programming

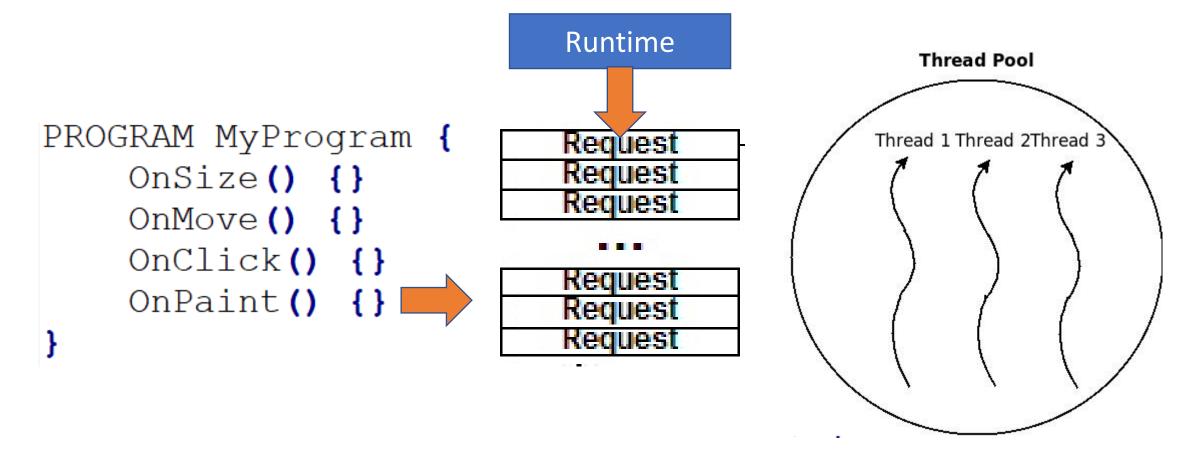
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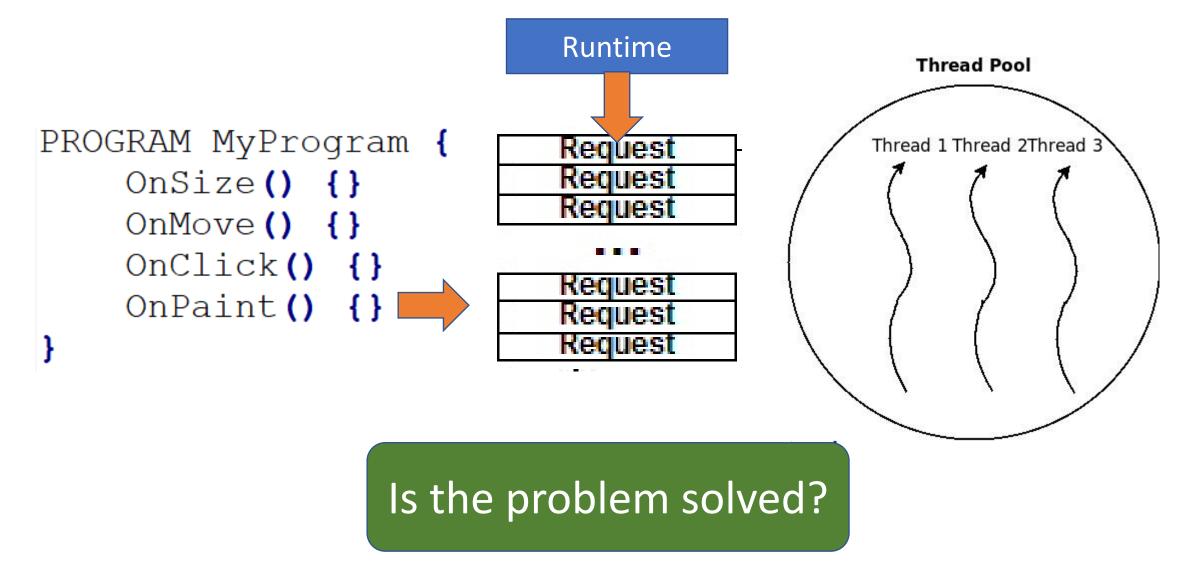
```
switch (message)
{
        //case WM_COMMAND:
          // handle menu selections etc.
       //break;
        //case WM PAINT:
         // draw our window - note: you must paint something here or not trap it!
        //break;
        case WM DESTROY:
             PostQuitMessage(0);
        break;
        default:
           // We do not want to handle this message so pass back to Windows
           // to handle it in a default way
            return DefWindowProc(hWnd, message, wParam, lParam);
}
```

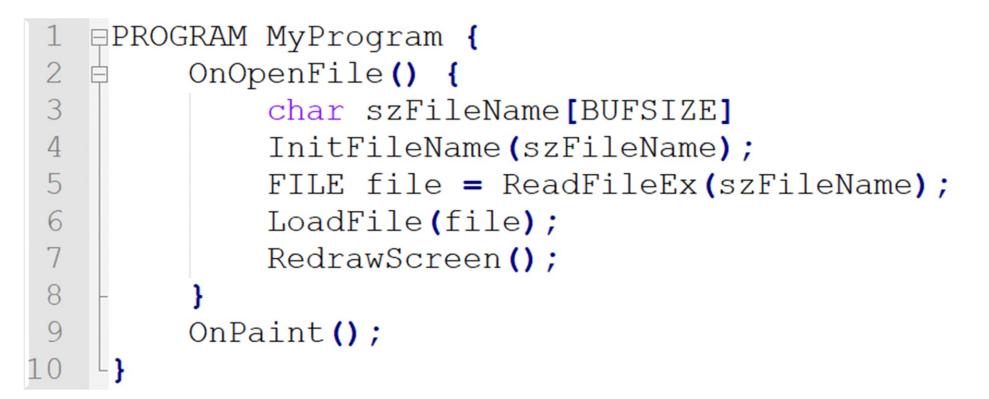
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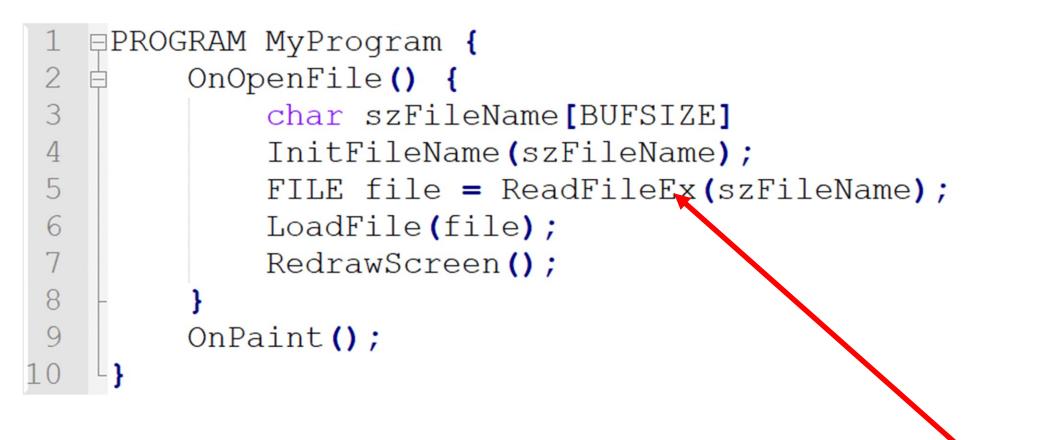
```
PROGRAM MyProgram {
    OnSize() {}
    OnMove() {}
    OnClick() {}
    OnPaint() {}
```



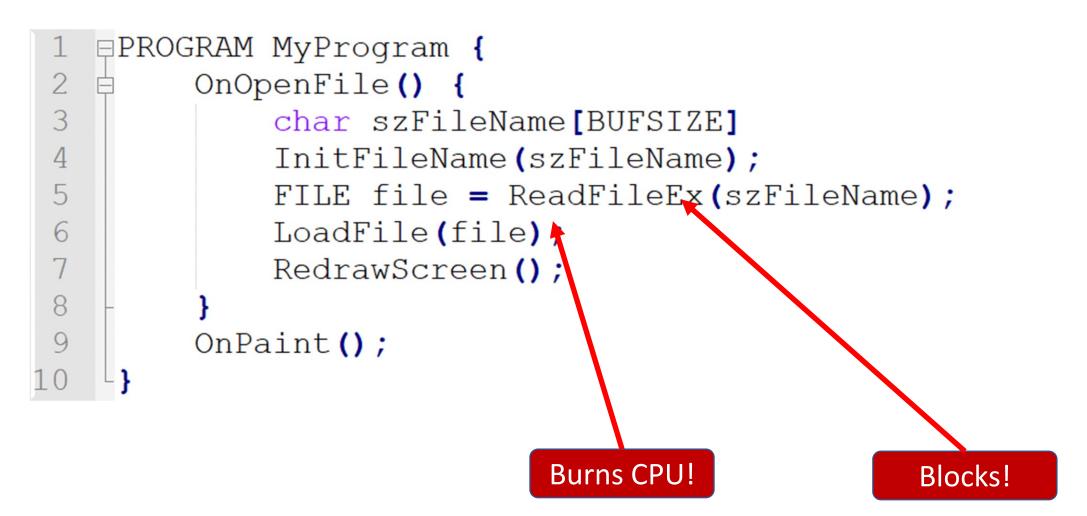


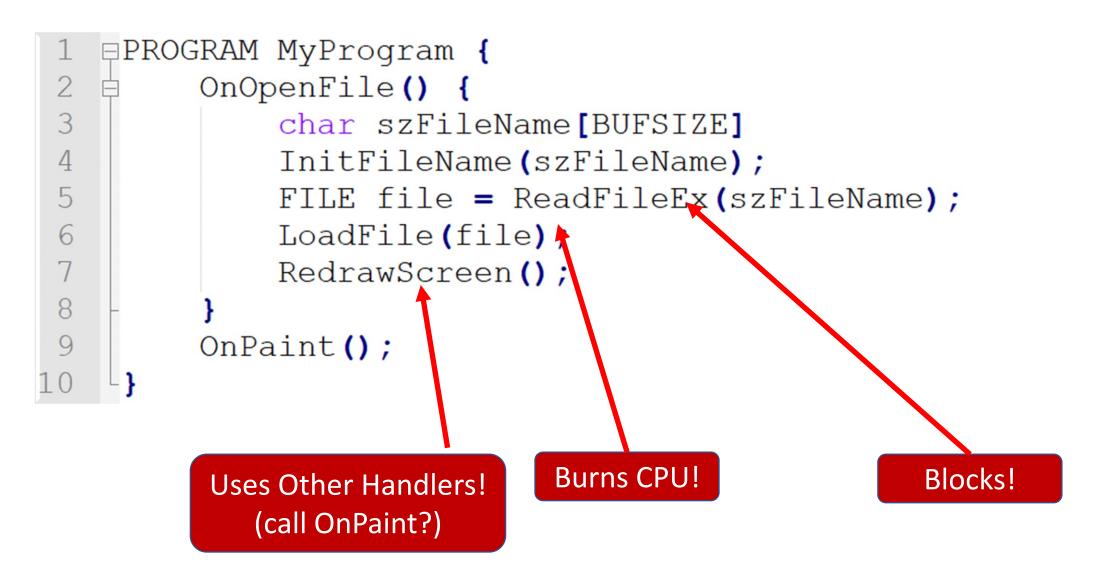






Blocks!



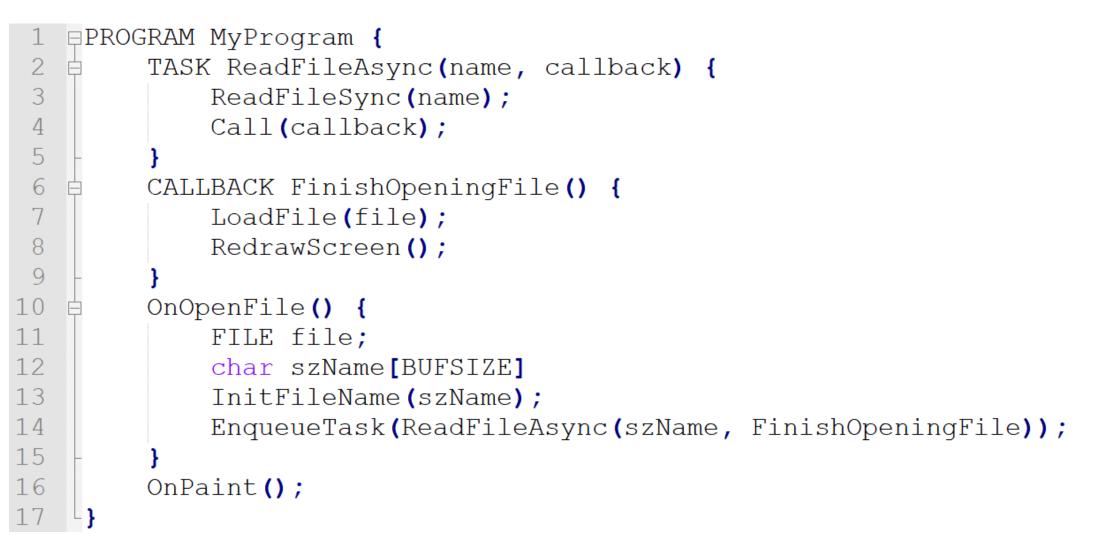


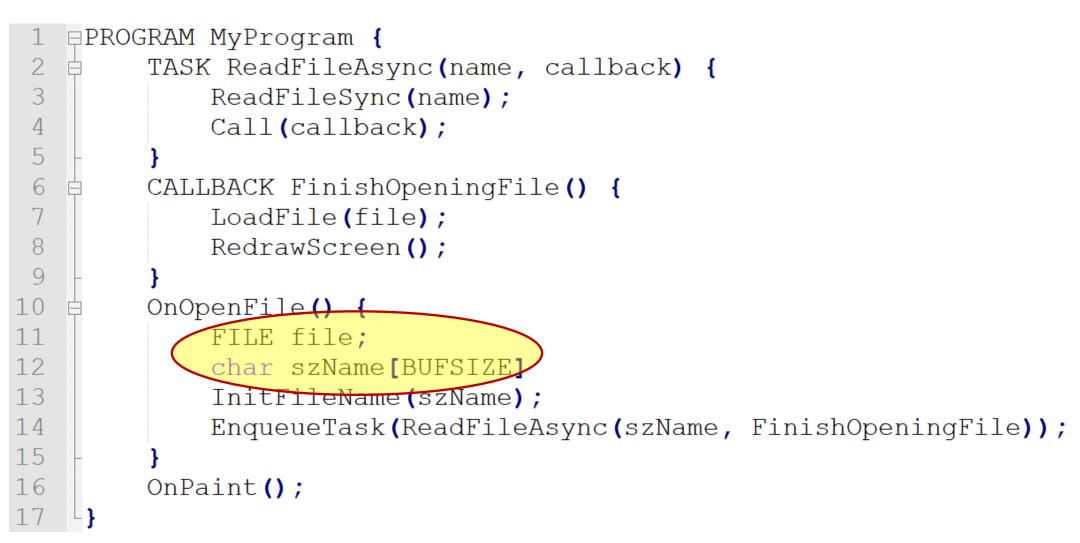
No problem! Just use more events/handlers, right?

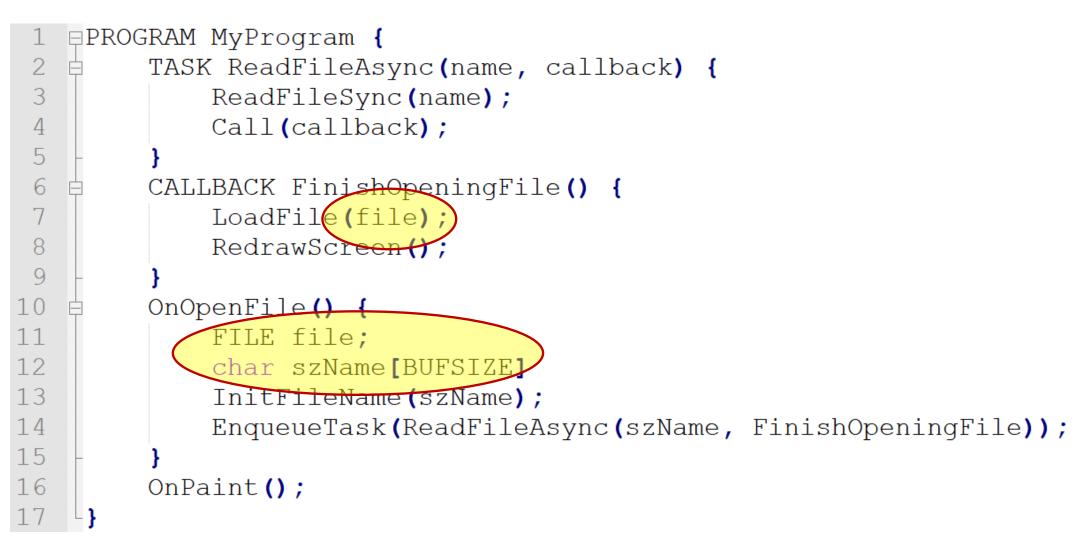
```
⊟PROGRAM MyProgram {
 2
         TASK ReadFileAsync(name, callback) {
 3
             ReadFileSync(name);
             Call(callback);
 4
 5
         CALLBACK FinishOpeningFile() {
 6
             LoadFile(file);
             RedrawScreen();
 8
 9
         OnOpenFile() {
10
11
             FILE file;
12
             char szName[BUFSIZE]
13
             InitFileName(szName);
14
             EnqueueTask(ReadFileAsync(szName, FinishOpeningFile));
15
16
         OnPaint();
```

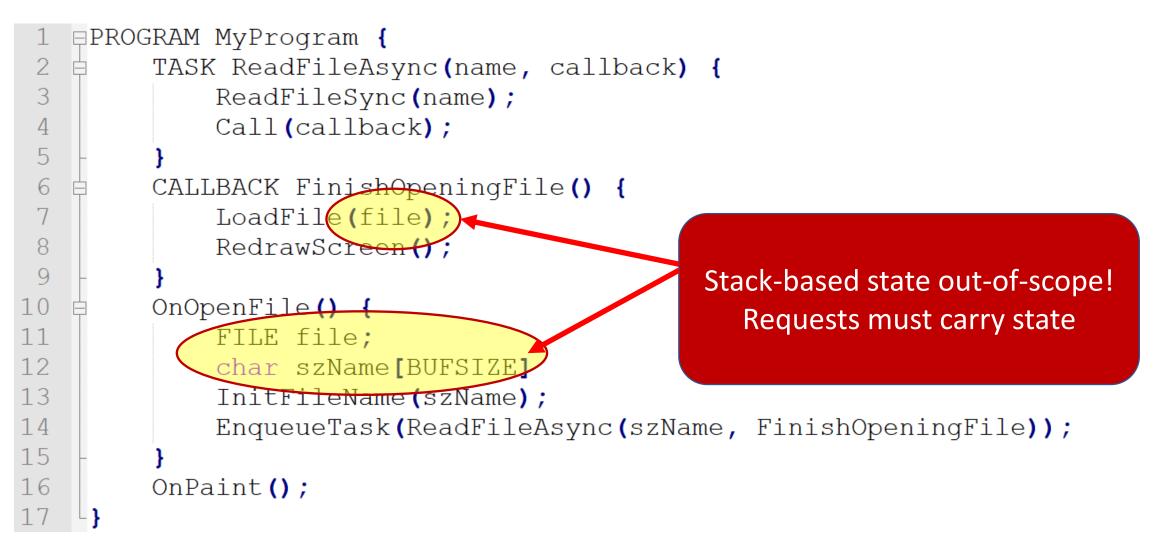
Continuations, BTW

```
□PROGRAM MyProgram {
2
        OnOpenFile() {
3
            ReadFile(file, FinishOpeningFile);
4
5
        OnFinishOpeningFile() {
            LoadFile(file, OnFinishLoadingFile);
6
 7
        OnFinishLoadingFile() {
8
            RedrawScreen();
9
10
        OnPaint();
12
```









• Thread Pros

• Event Pros

• Thread Cons

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- Thread Pros
 - Overlap I/O and computation
 - While looking sequential
 - Intermediate state on stack
 - Control flow naturally expressed
- Thread Cons
 - Synchronization required
 - Overflowable stack
 - Stack memory pressure

- Event Pros
 - Easier to create well-conditioned system
 - Easier to express dynamic change in level of parallelism

- Event Cons
 - Difficult to program
 - Control flow between callbacks obscure
 - When to deallocate memory
 - Incomplete language/tool/debugger support
 - Difficult to exploit concurrent hardware

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Language-level Futures: the sweet spot?

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Thread Pool Implementation

```
void
```

```
ThreadPool::StartThreads (
    in UINT uiThreads,
      in BOOL bWaitAllThreadsAlive
    Lock();
    if (uiThreads != 0 && m vhThreadDescs.size() < m uiTargetSize)
        ResetEvent (m hAllThreadsAlive);
    while(m vhThreadDescs.size() < m uiTargetSize) {</pre>
        for(UINT i=0; i<uiThreads; i++) {</pre>
            THREADDESC* pDesc = new THREADDESC(this);
            HANDLE * phThread = &pDesc->hThread;
            *phThread = CreateThread(NULL, 0, ThreadPoolProc, pDesc, 0, NULL);
            m vhAvailable.push back(*phThread);
            m vhThreadDescs[*phThread] = pDesc;
    m uiThreads = (UINT)m vhThreadDescs.size();
    Unlock();
    if(bWaitAllThreadsAlive)
        WaitThreadsAlive();
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```

Cool project idea: build a thread pool!

Thread Pool Implementation

```
DWORD
ThreadPool::ThreadPoolProc(
     in THREADDESC * pDesc
   HANDLE hThread = pDesc->hThread;
   HANDLE hStartEvent = pDesc->hStartEvent;
   HANDLE hRuntimeTerminate = PTask::Runtime::GetRuntimeTerminateEvent();
   HANDLE vEvents[] = { hStartEvent, hRuntimeTerminate };
   NotifyThreadAlive(hThread);
    while(!pDesc->bTerminate) {
        DWORD dwWait = WaitForMultipleObjects(dwEvents, vEvents, FALSE, INFINITE);
        pDesc->Lock();
        pDesc->bTerminate |= bTerminate;
        if(pDesc->bRoutineValid && !pDesc->bTerminate) {
            LPTHREAD START ROUTINE lpRoutine = pDesc->lpRoutine;
            LPVOID lpParameter = pDesc->lpParameter;
            pDesc->bActive = TRUE;
            pDesc->Unlock();
            dwResult = (*lpRoutine) (lpParameter);
            pDesc->Lock();
            pDesc->bActive = FALSE;
            pDesc->bRoutineValid = FALSE;
        pDesc->Unlock();
        Lock();
        m vhInFlight.erase(pDesc->hThread);
        if(!pDesc->bTerminate)
            m vhAvailable.push back(pDesc->hThread);
        Unlock();
   NotifyThreadExit(hThread);
   return dwResult;
```

ThreadPool Implementation

```
BOOL
ThreadPool::SignalThread(
    _____in HANDLE hThread
    )
{
    Lock();
    BOOL bResult = FALSE;
    std::set<HANDLE>::iterator si = m_vhWaitingStartSignal.find(hThread);
    if(si!=m_vhWaitingStartSignal.end()) {
        m_vhWaitingStartSignal.erase(hThread);
        THREADDESC * pDesc = m_vhThreadDescs[hThread];
        HANDLE hEvent = pDesc->hStartEvent;
        SetEvent(hEvent);
        bResult = TRUE;
    }
    Unlock();
    return bResult;
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