

Programming at Scale: Dataflow

cs378h

Today

Questions?

Administrivia

- Project Proposal Due Soon!

Agenda:

- MPI Wrapup
- Dataflow

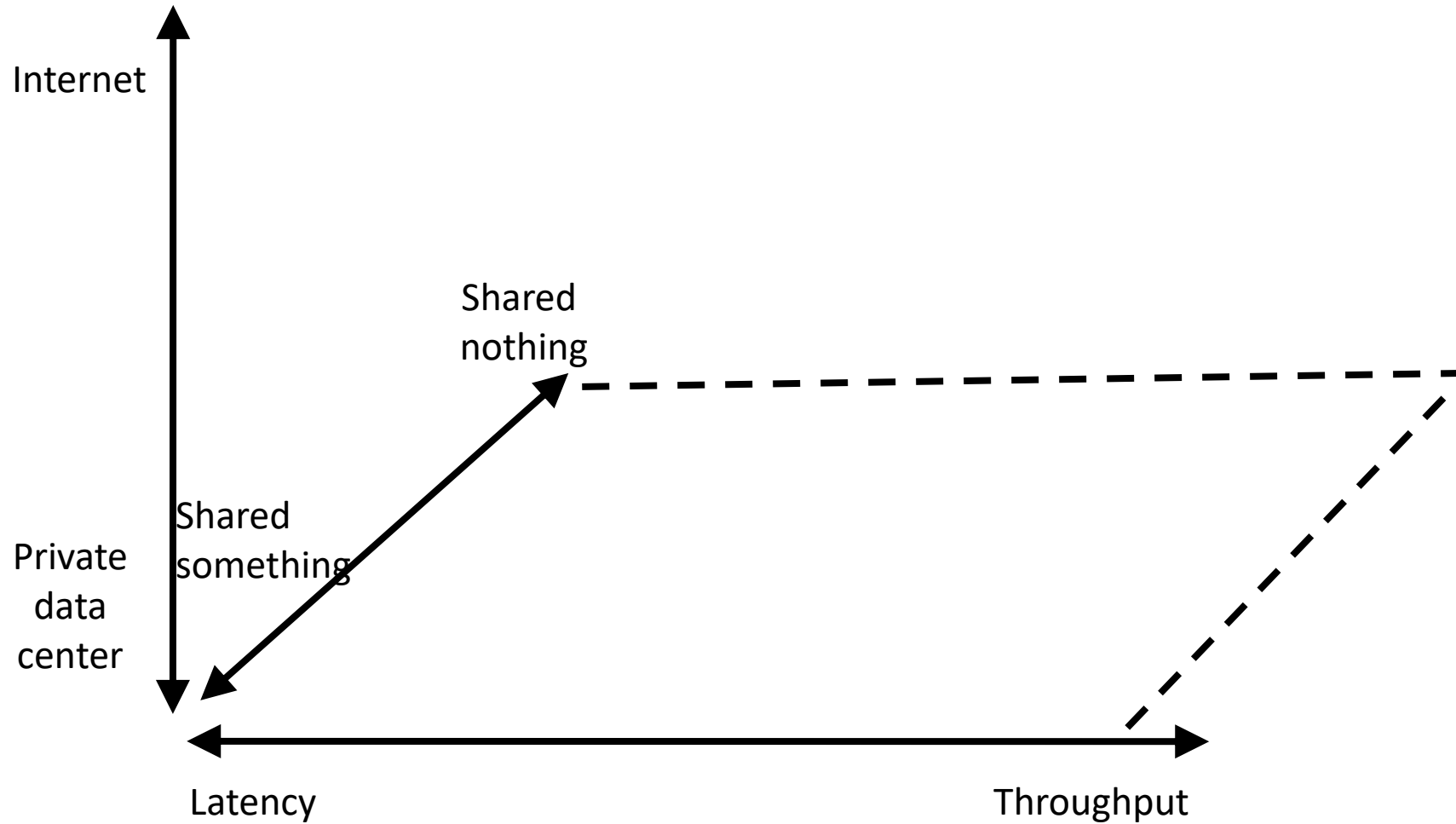
MapReduce faux quiz (5 min, any 2):

- Have you ever written a MapReduce program? If so, what was it?
- What phenomena can slow down a map task?
- Do reducers wait for all their mappers before starting? Why/why not?
- What machine resources does the shuffle phase consume most?
- Is it safe to re-execute a failed map/reduce job sans cleanup? Why [not]?
- How does MR handle master failure? What are the alternatives?

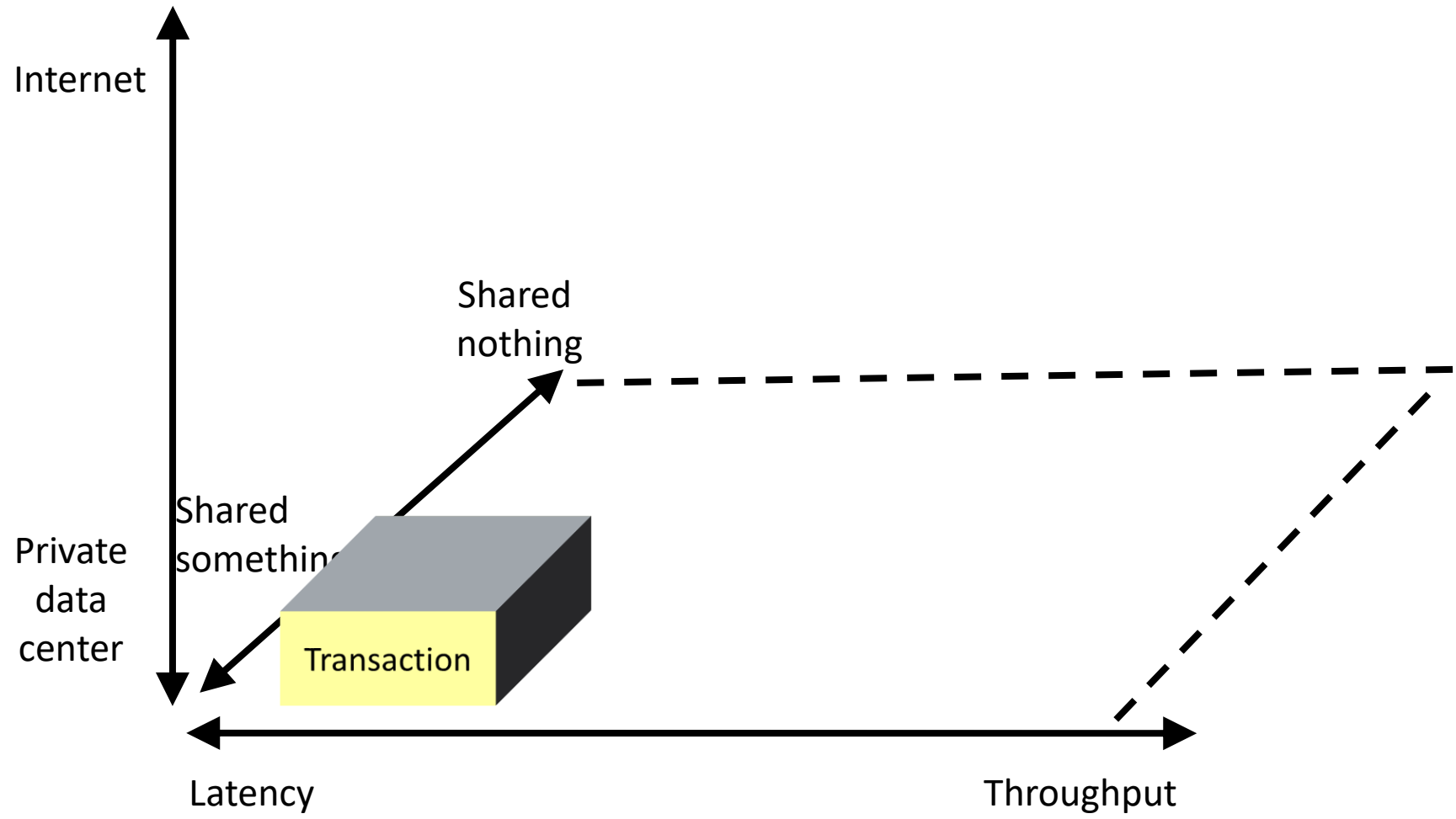
Review: Scale: Goal



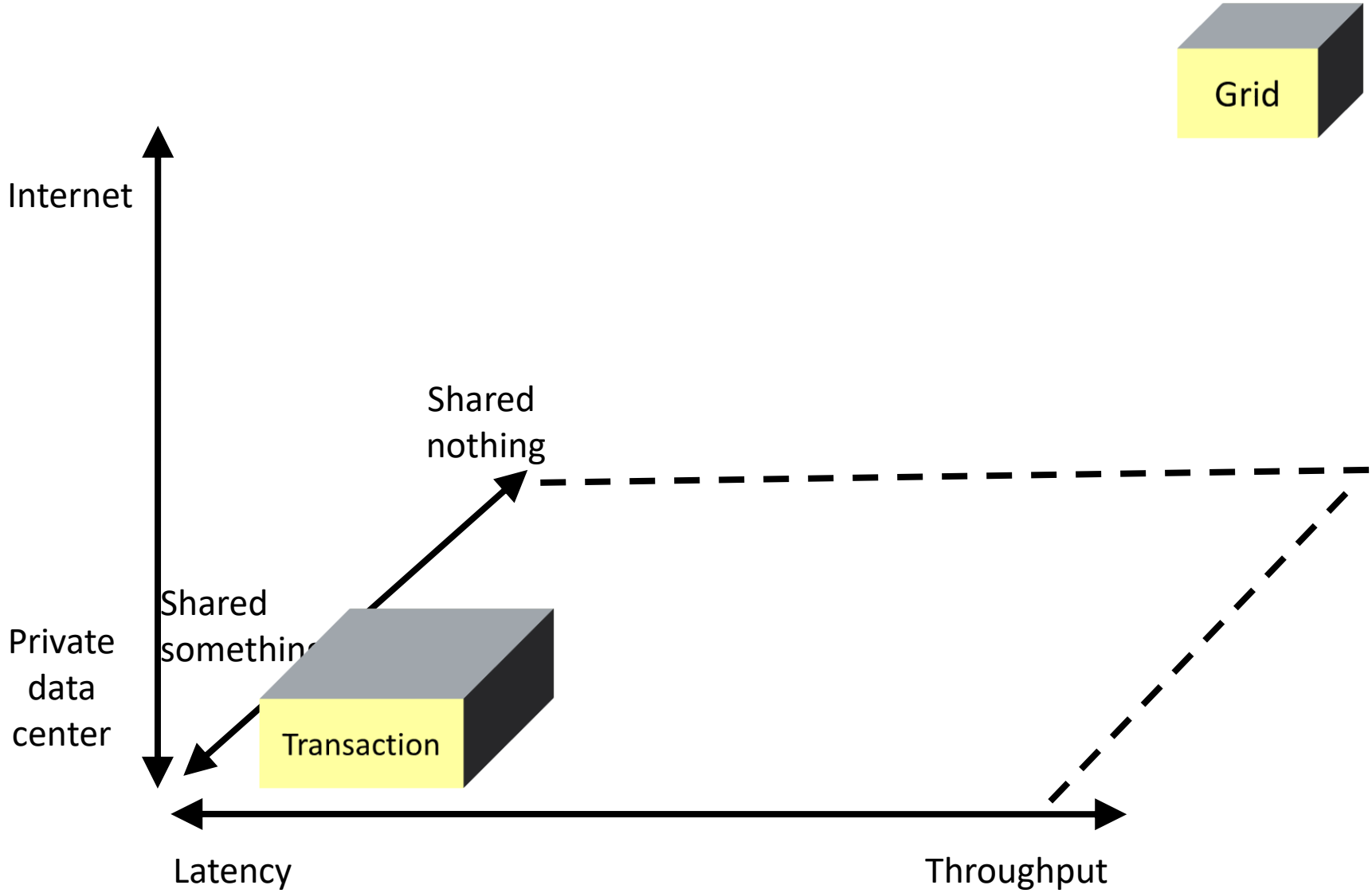
Review: Design Space



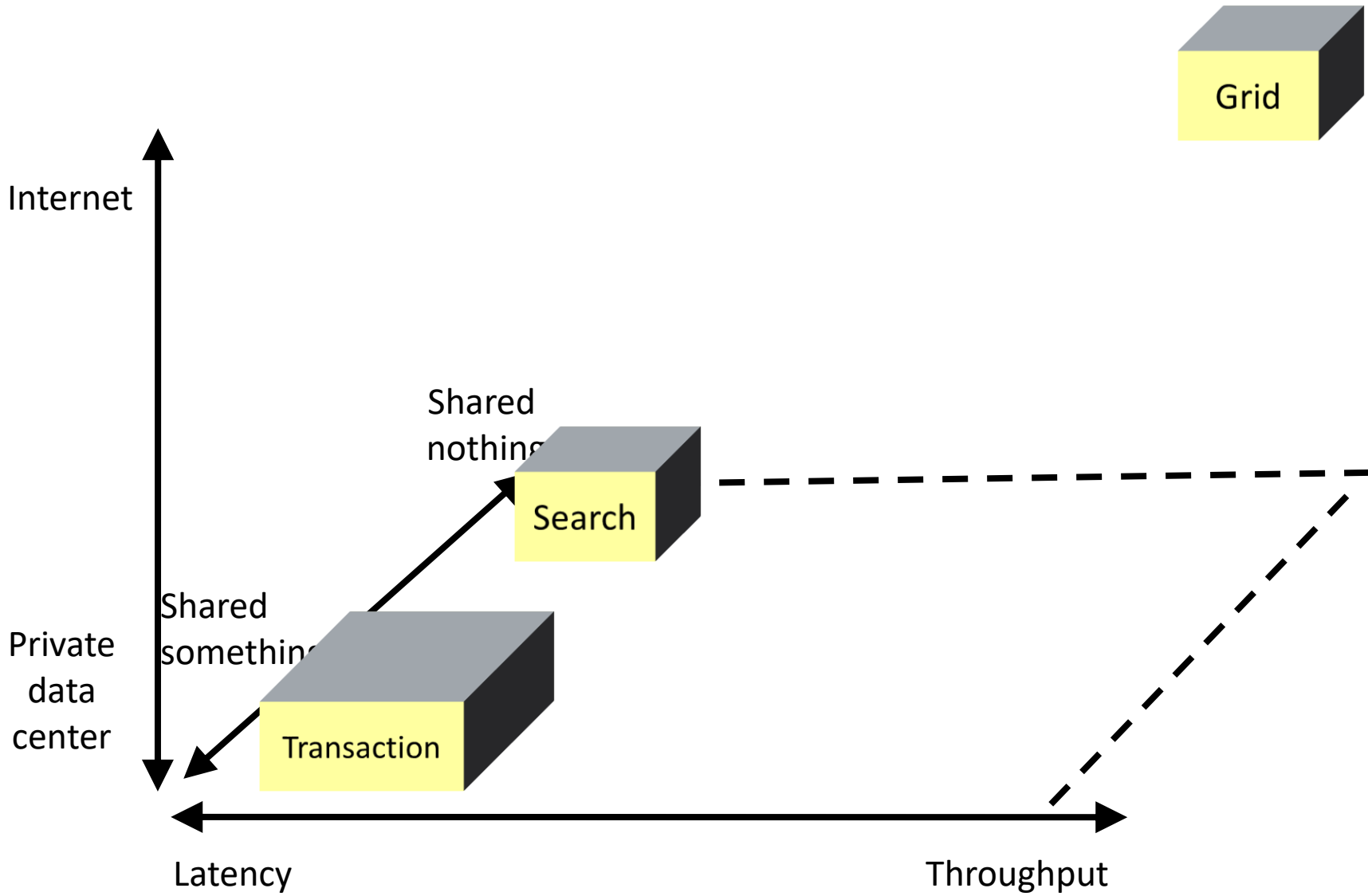
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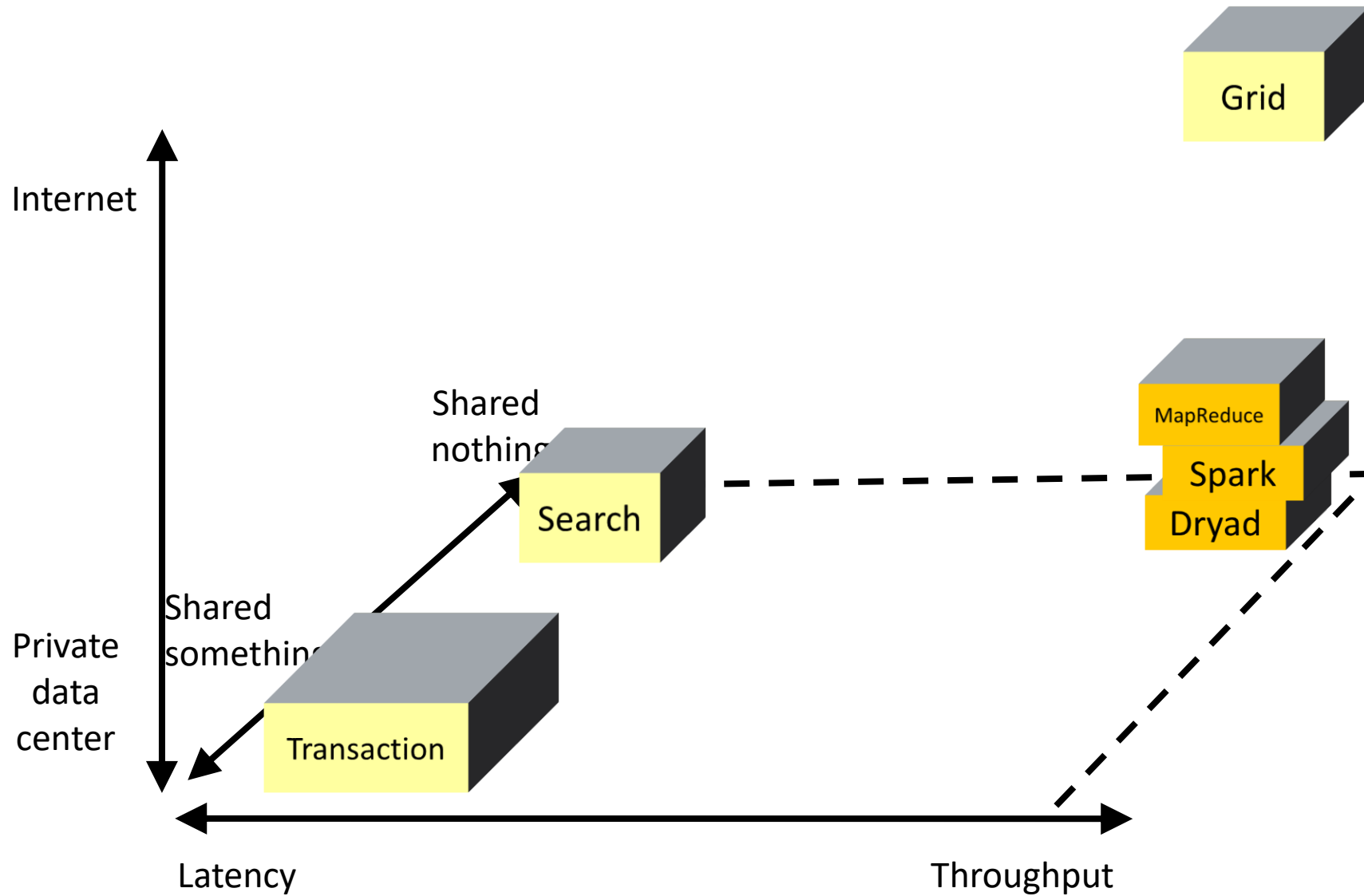
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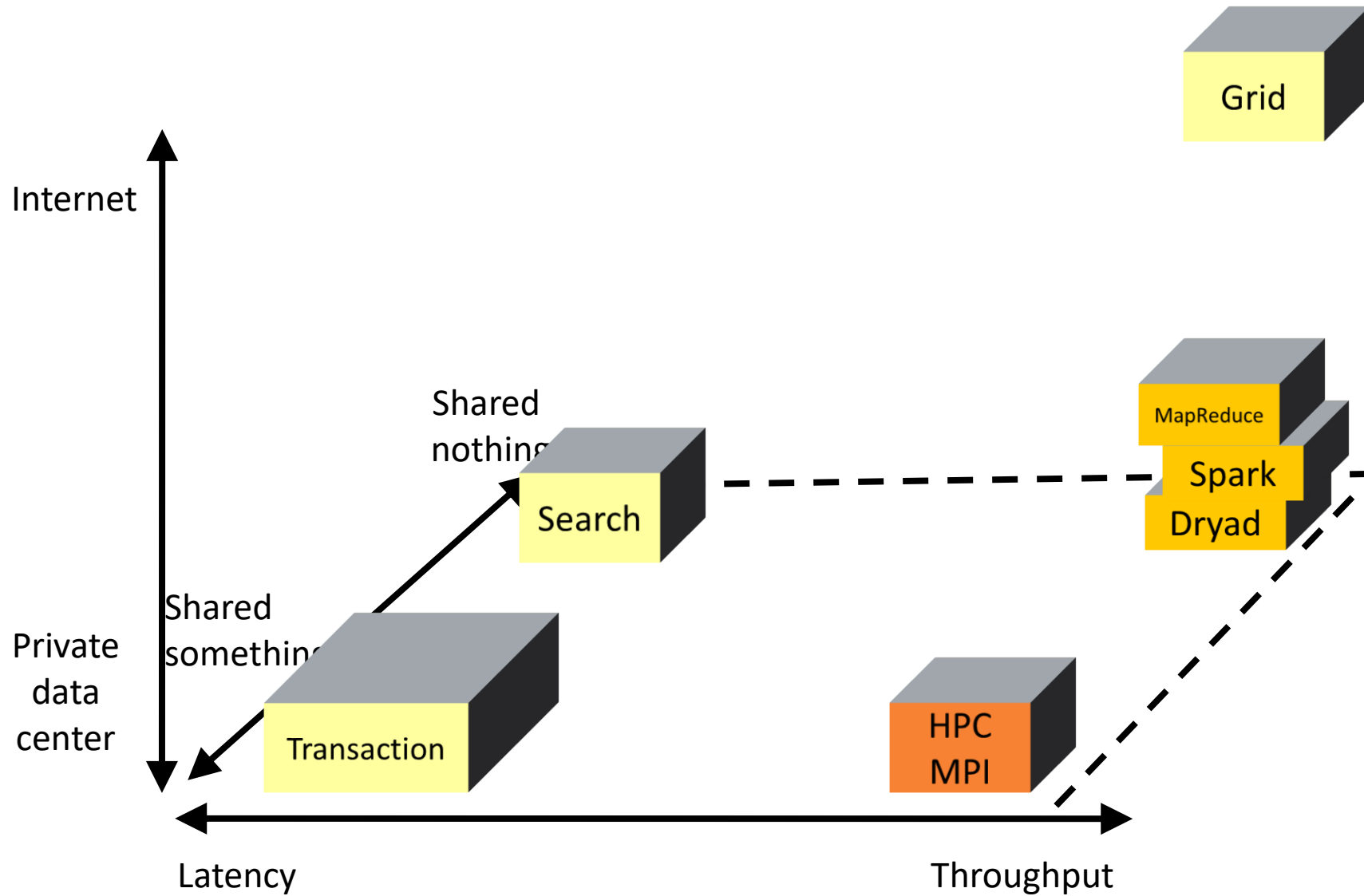
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Review: Design Space



You are an engineer at:
Hare-brained-scheme.com

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Input: <search_term>, <files>

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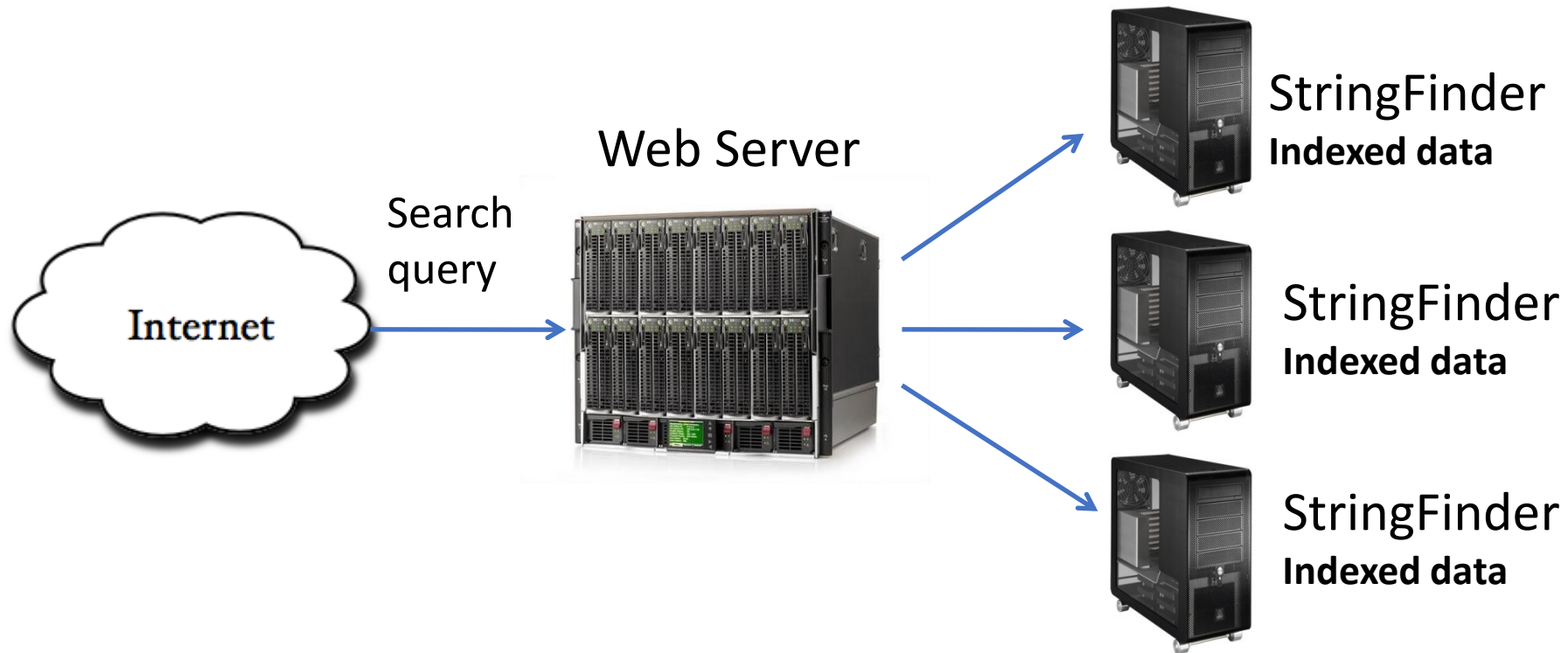
Input: <search_term>, <files>

Output: list of files containing <search_term>

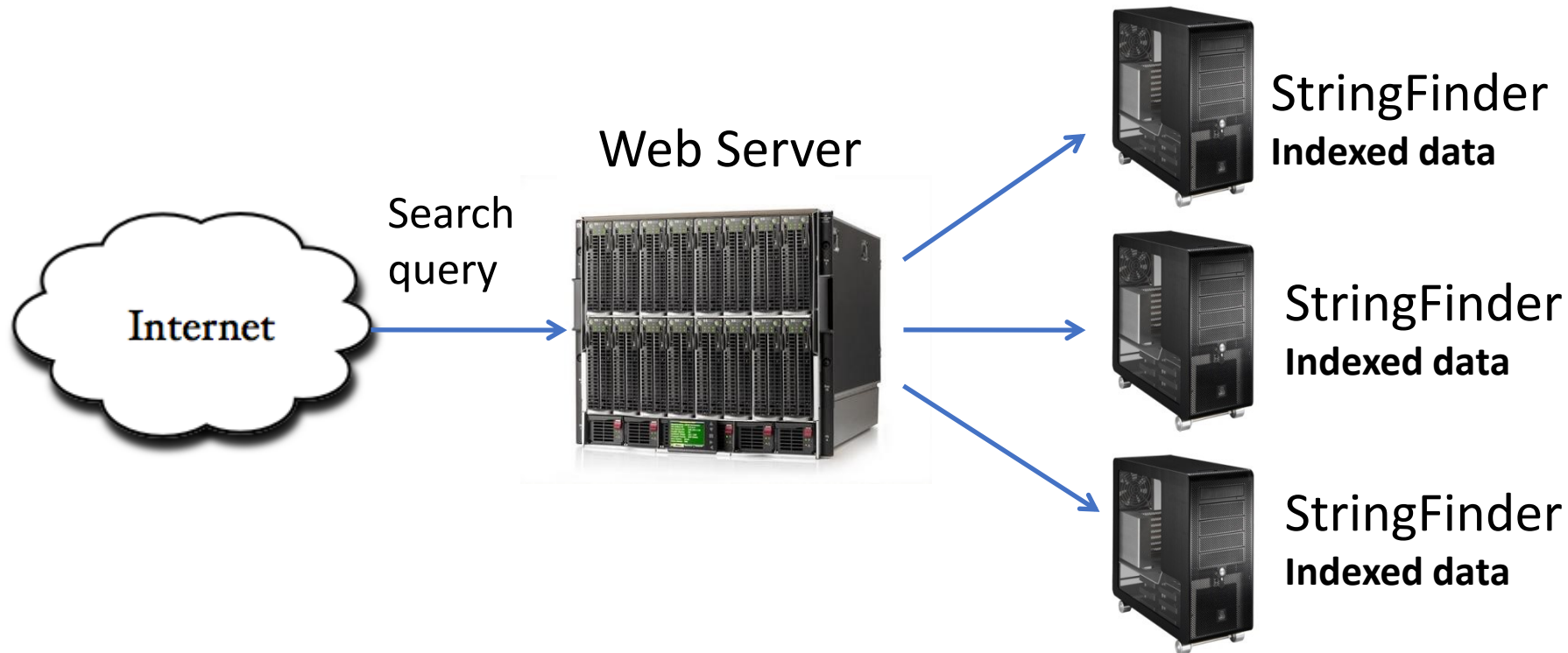
One Solution

```
public class StringFinder {  
    int main(...) {  
        foreach(File f in getInputFiles()) {  
            if(f.contains(searchTerm))  
                results.add(f.getFileName());  
        }  
        System.out.println("Files:" + results.toString());  
    }  
}
```

Infrastructure is hard to get right

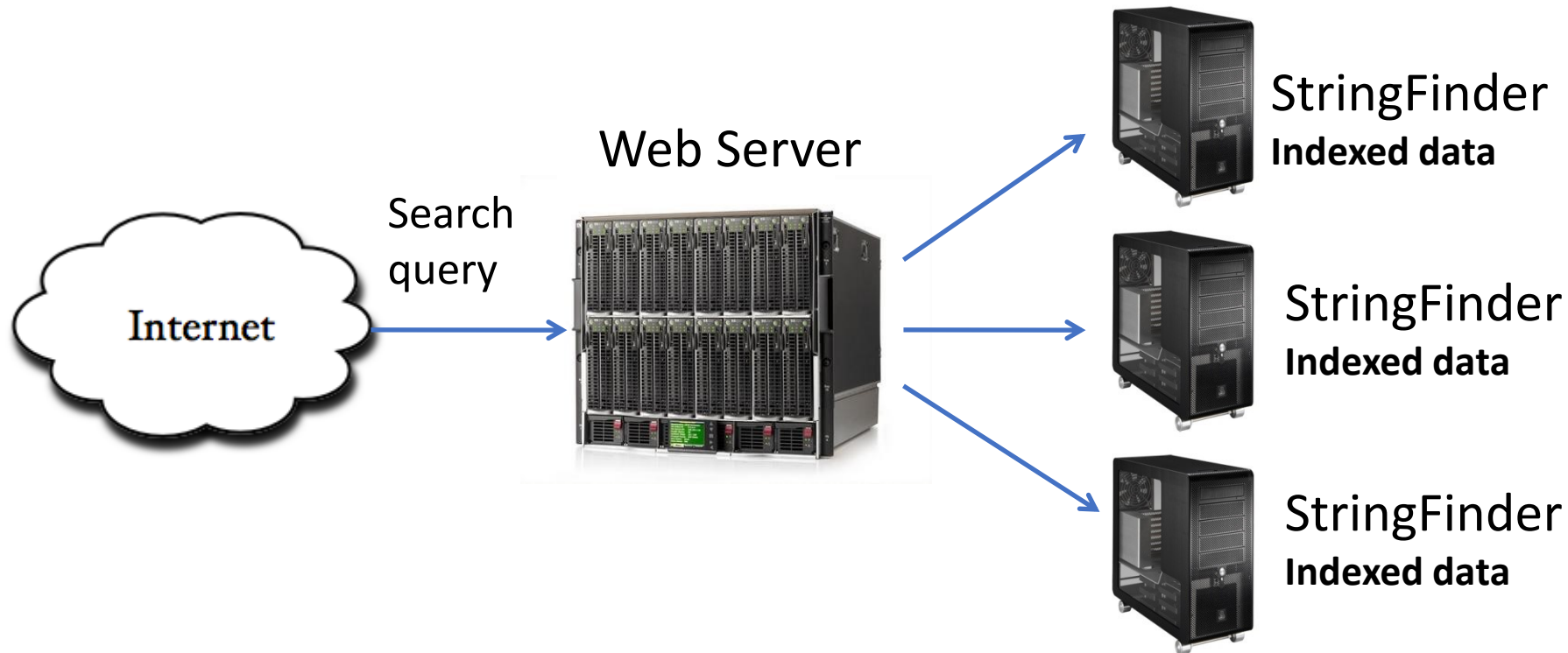


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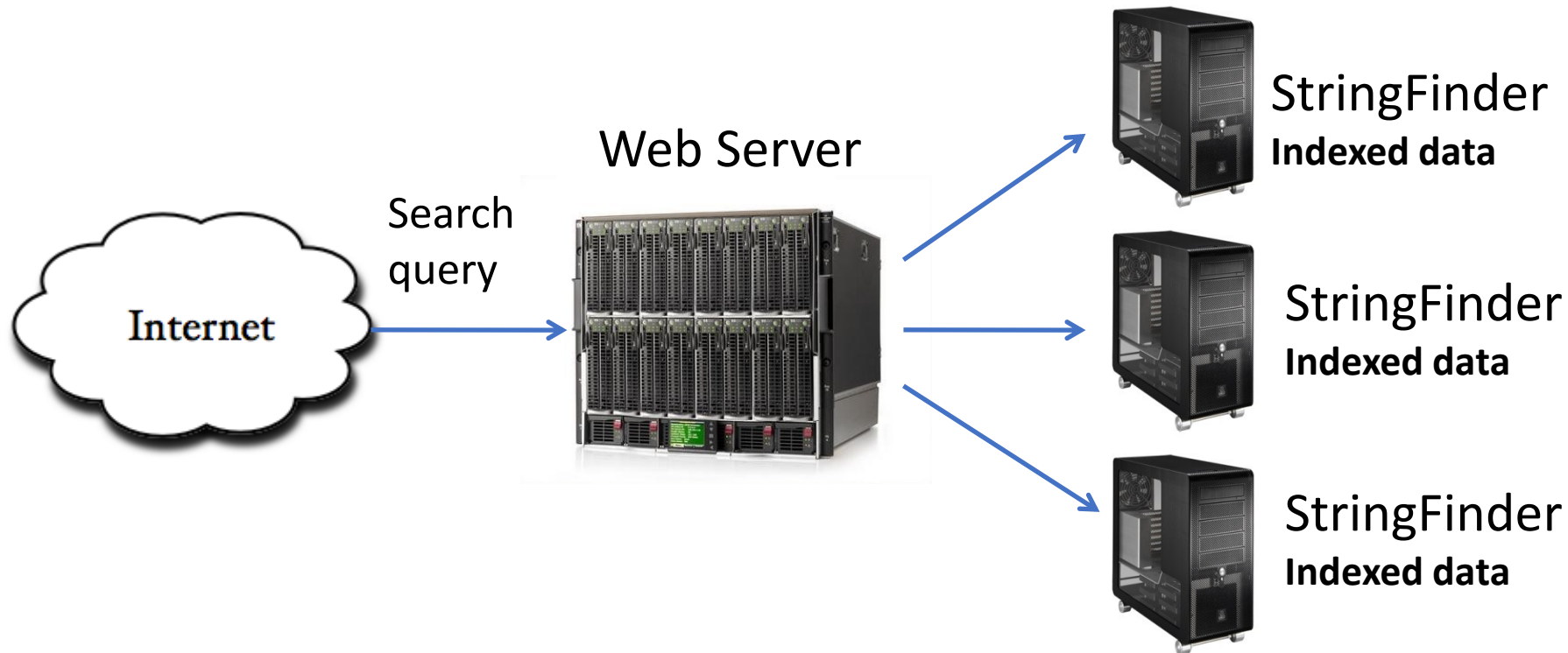
1. How do we distribute the searchable files on our machines?

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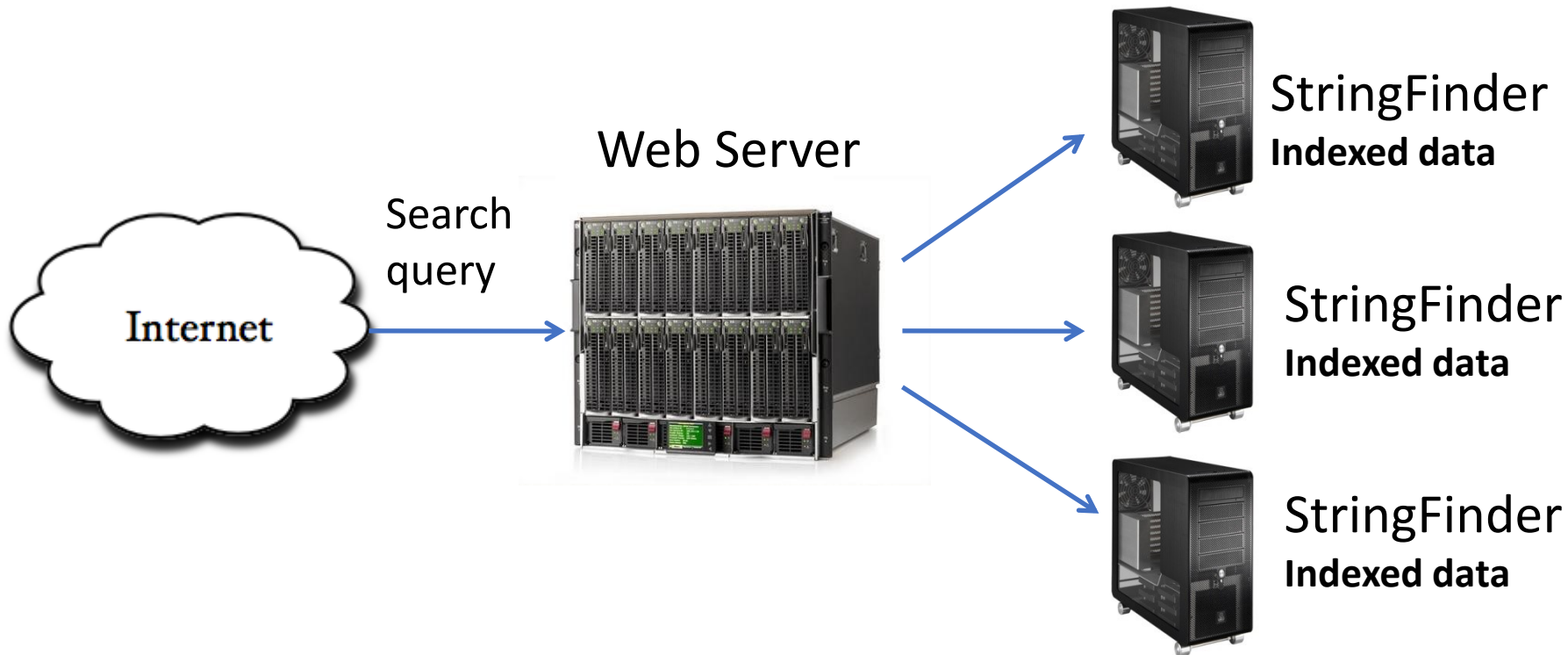
1. How do we distribute the searchable files on our machines?
2. What if our webserver goes down?

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2. What if our webserver goes down?
3. What if a StringFinder machine dies? How would you know it was dead?

Infrastructure is hard to get right



1. How do we distribute the searchable files on our machines?
2. What if our webserver goes down?
3. What if a StringFinder machine dies? How would you know it was dead?
4. **What if marketing comes and says, “well, we also want to show pictures of the earth from space too! Ooh..and the moon too!”**

StringFinder was the easy part!

You really need general infrastructure.

Many different tasks

Want to use hundreds or thousands of PC's

Continue to function if something breaks

Must be easy to program...

Dataflow Engines

Programming model + infrastructure

Write programs that run on lots of machines

Automatic parallelization and distribution

Fault-tolerance

I/O and jobs Scheduling

Status and monitoring

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Key Ideas:

*All modern “big data” platforms are **dataflow engines!***

Differences:

1. what graph structures are allowed?
2. How does this impact programming model?

MapReduce

- Input & Output: sets of <key, value> pairs
- Programmer writes 2 functions:

```
map (in_key, in_value) -> list(out_key,  
    intermediate_value)
```

- Processes <k,v> pairs
- Produces intermediate pairs

```
reduce (out_key, list(interm_val)) ->  
    list(out_value)
```

- Combines intermediate values for a key
- Produces a merged set of outputs

Indexing (1)

```
public void map() {  
    String line = value.toString();  
    StringTokenizer itr = new StringTokenizer(line);  
    if(itr.countTokens() >= N) {  
        while(itr.hasMoreTokens()) {  
            word = itr.nextToken()+"|"+key.getFileName();  
            output.collect(word, 1);  
        }  
    }  
}
```

Input: a line of text, e.g. **“mistakes were made”** from **myfile.txt**
Output:
mistakes|myfile.txt
were|myfile.txt
made|myfile.txt

Indexing (2)

```
public void reduce() {  
    int sum = 0;  
    while(values.hasNext()) {  
        sum += values.next().get();  
    }  
    output.collect(key, sum);  
}
```

Input: a <term,filename> pair, list of occurrences (e.g. {1, 1,..1})

Output:

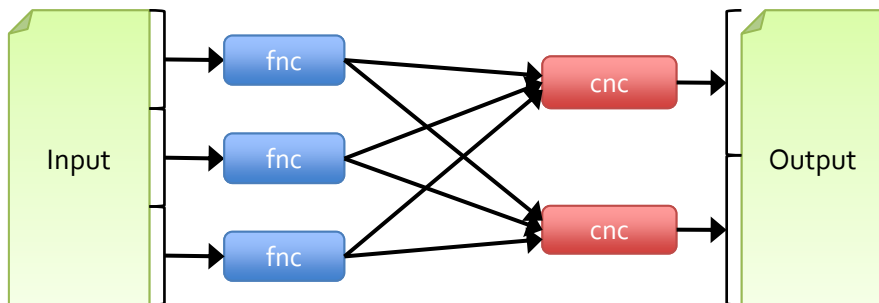
mistakes myfile.txt	10
were myfile.txt	45
made myfile.txt	2

Review: K-Means

```
public void kmeans() {  
    while (...) {  
        for each point  
            find_nearest_center(point) ;  
        for each center  
            compute_new_center(center)  
    }  
}
```

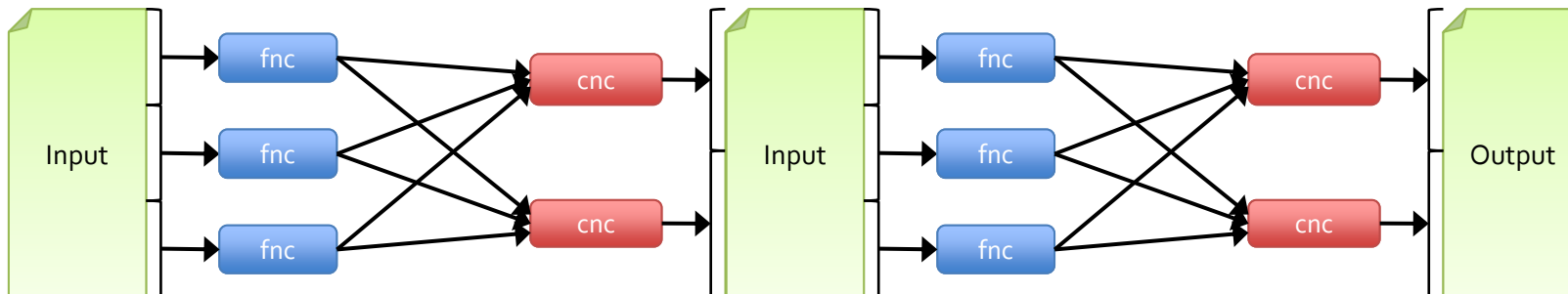
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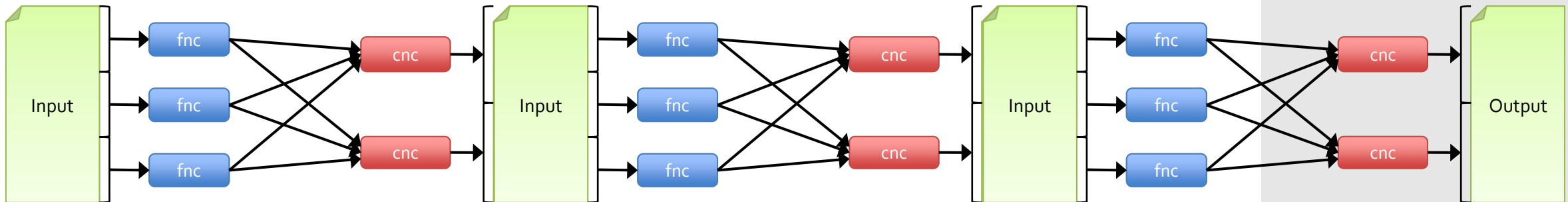
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```



Example: K-Means Mapper

```
/*
 * Map: find minimum distance center for point, emit to reducer
 */
@Override
public void map(LongWritable key, Text value,
                OutputCollector<DoubleWritable, DoubleWritable> output,
                Reporter reporter) throws IOException {
    String line = value.toString();
    double point = Double.parseDouble(line);
    double min1, min2 = Double.MAX_VALUE, nearest_center = mCenters.get(0);
    // Find the minimum center from a point
    for (double c : mCenters) {
        min1 = c - point;
        if (Math.abs(min1) < Math.abs(min2)) {
            nearest_center = c;
            min2 = min1;
        }
    }
    // Emit the nearest center and the point
    output.collect(new DoubleWritable(nearest_center),
                  new DoubleWritable(point));
}
```

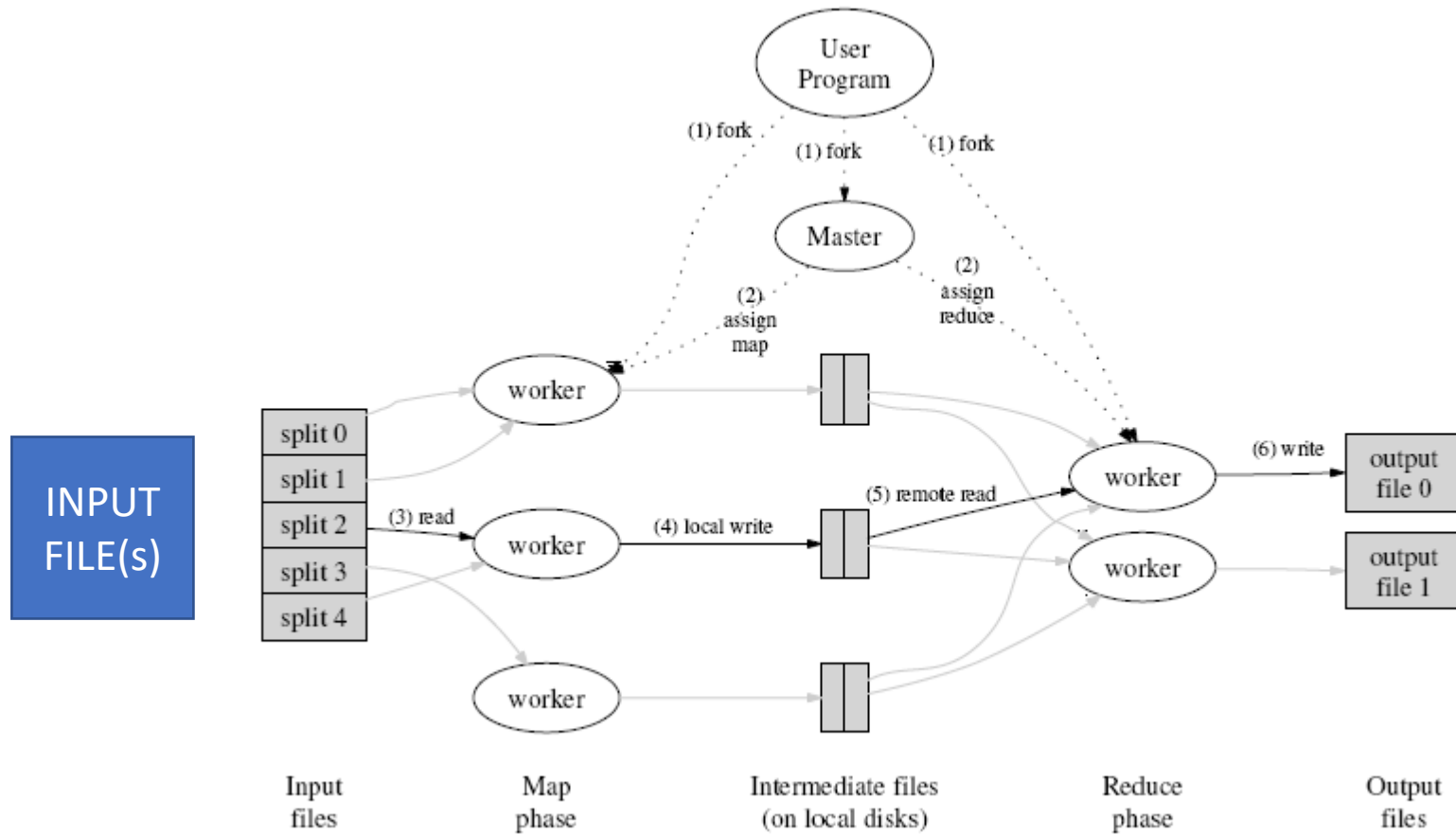
Example: K-Means Reducer

```
/*
 * Reduce: collect all points per center and calculate
 * the next center for those points
 */
@Override
public void reduce(
    DoubleWritable key, Iterator<DoubleWritable> values,
    OutputCollector<DoubleWritable, Text> output, Reporter reporter)
    throws IOException {
    double newCenter;
    double sum = 0;
    int no_elements = 0;
    String points = "";
    while (values.hasNext()) {
        double d = values.next().get();
        points = points + " " + Double.toString(d);
        sum = sum + d;
        ++no_elements;
    }

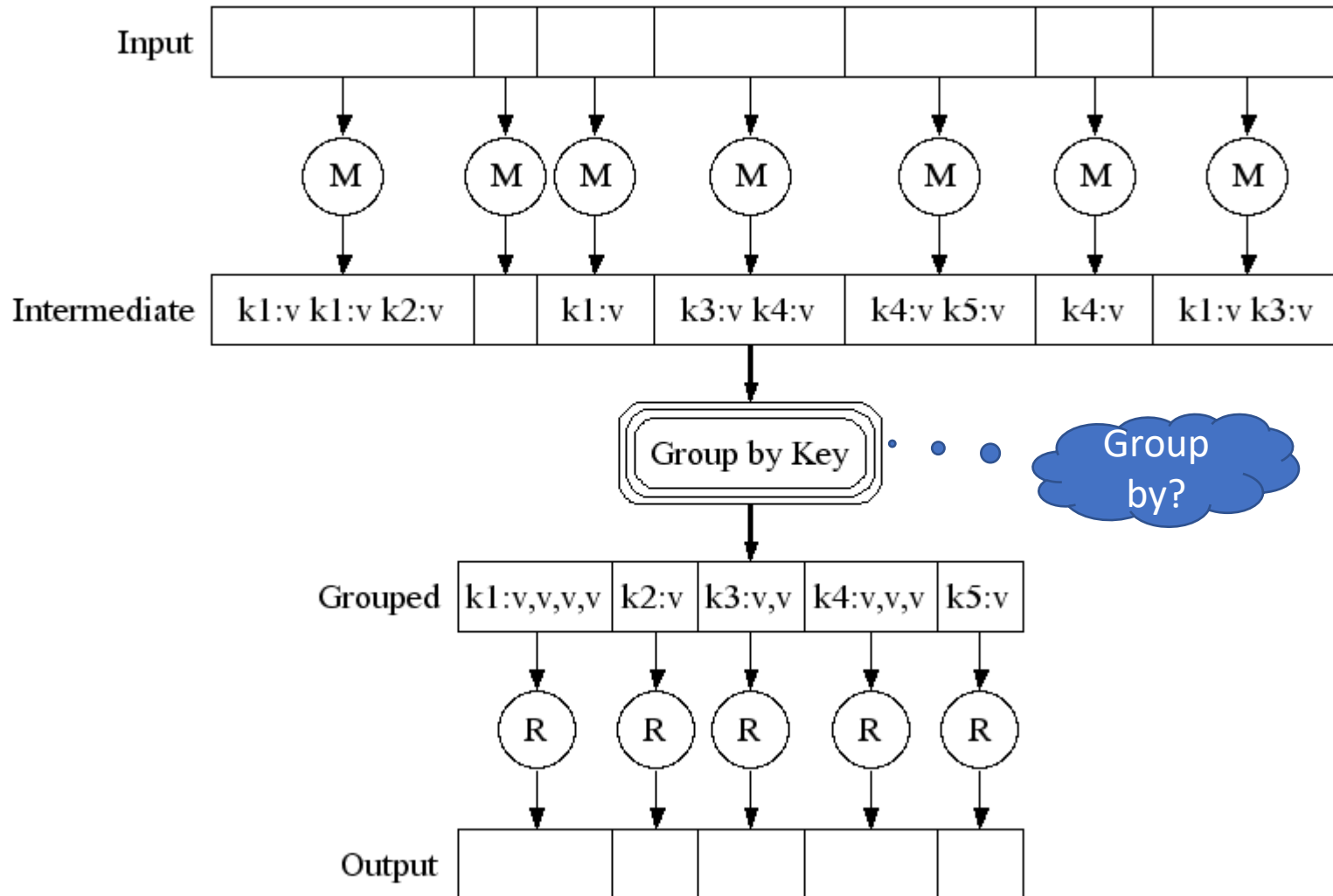
    // We have a new center now
    newCenter = sum / no_elements;

    // Emit new center and point
    output.collect(new DoubleWritable(newCenter), new Text(points));
}
```

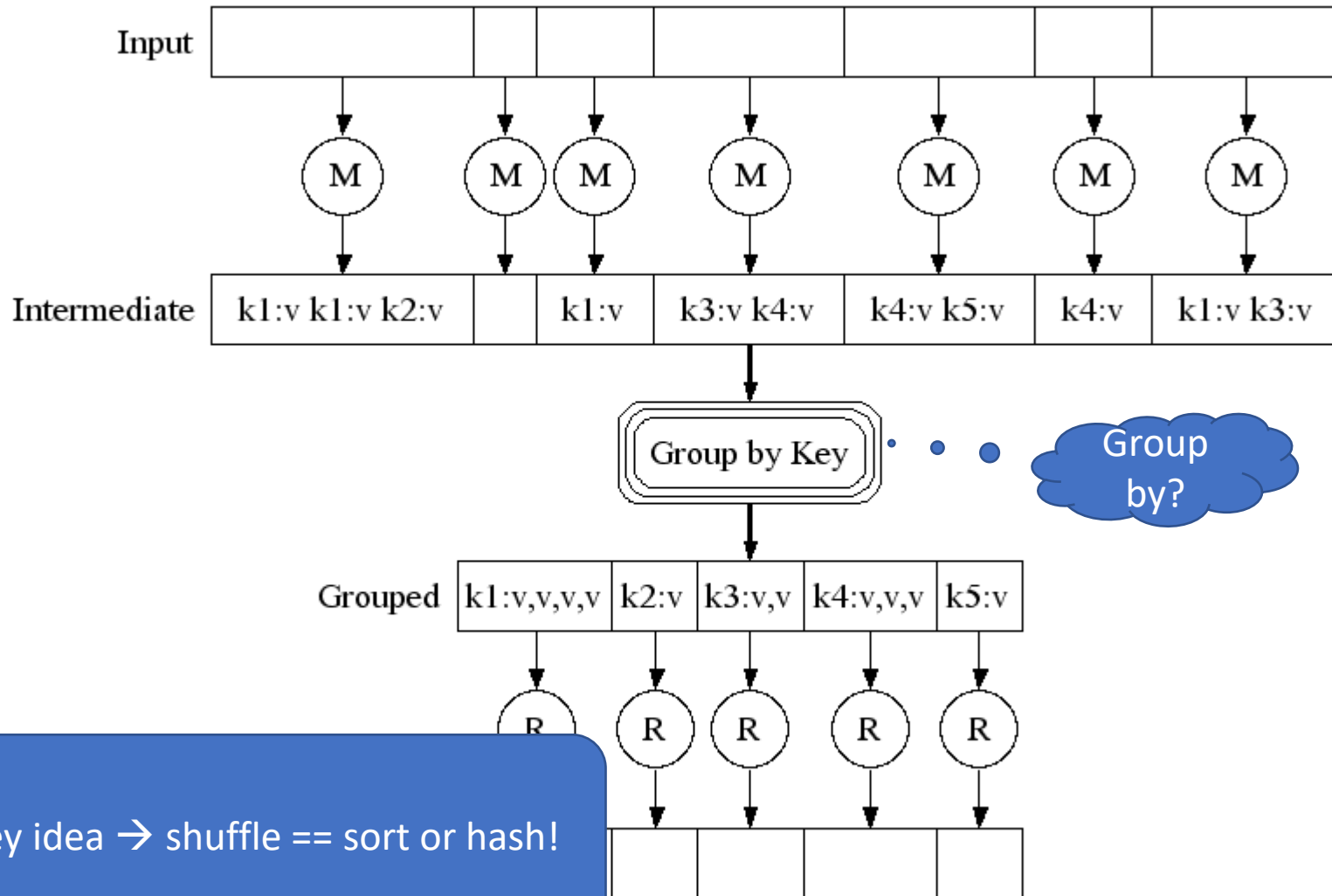

How Does Parallelization Work?



Execution



Execution



Key idea → shuffle == sort or hash!

Task Granularity And Pipelining

|map tasks| >> |machines| -- why?

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Minimize fault recovery time

Pipeline map with other tasks

Easier to load balance dynamically

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Why might he say this?

MapReduce: not without Controversy

MapReduce: A major step backwards | The Database Column

<http://databasecolumn.vertica.com/database-innovation/mapreduce-a-major-step-backwards/>

September 6, 2011

on Jan 17 in [Database architecture](#), [Database history](#), [Database innovation](#) posted by [DeWitt](#)

[Note: Although the system attributes this post to a single author, it was written by David J. DeWitt and Michael Stonebraker]

On January 8, a Database Column reader asked for our views on new distributed database research efforts, and we'll begin here with our views on [MapReduce](#). This is a good time to discuss it, since the recent trade press has been filled with news of the revolution of so-called "cloud computing." This paradigm entails harnessing large numbers of (low-end) processors working in parallel to solve a computing problem. In effect, this suggests constructing a data center by lining up a large number of "jelly beans" rather than utilizing a much smaller number of high-end servers.

For example, IBM and Google have announced plans to make a 1,000 processor cluster available to a few select universities to teach students how to program such clusters using a software tool called MapReduce [1]. Berkeley has gone so far as to plan on teaching their freshman how to program using the MapReduce framework.

As both educators and researchers, we are amazed at the hype that the MapReduce proponents have spread about how it represents a paradigm shift in the development of scalable, data-intensive applications. MapReduce may be a good idea for writing certain types of general-purpose computations, but to the database community, it is:

1. A giant step backward in the programming paradigm for large-scale data intensive applications

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Missing most DBMS features

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So why is it such a big success?

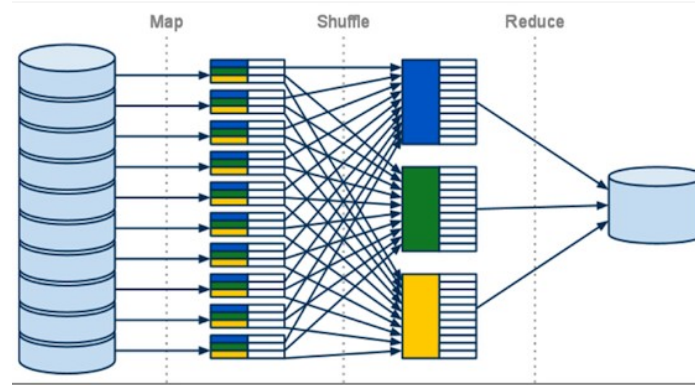
MapReduce and Dataflow

MapReduce and Dataflow

- MR is a *dataflow* engine

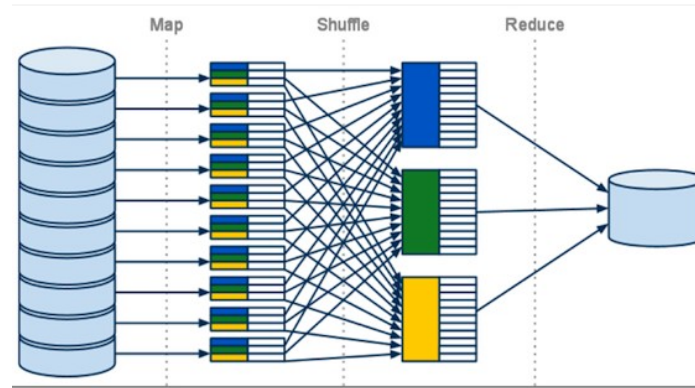
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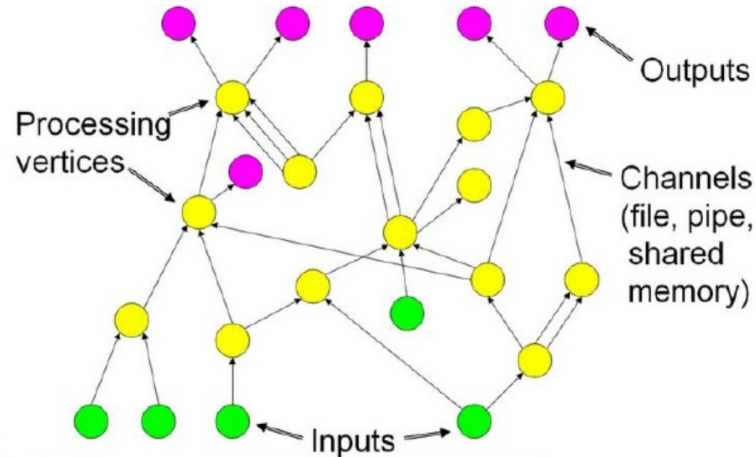
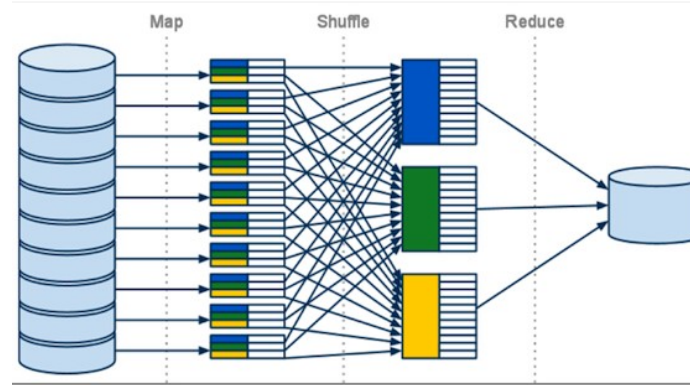
MapReduce and Dataflow

- MR is a ***dataflow*** engine
- Lots of others
 - Dryad
 - DryadLINQ
 - Dandelion
 - CIEL
 - GraphChi/Pregel
 - Spark



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MapReduce vs Dryad (and others...)

DAG instead of BSP

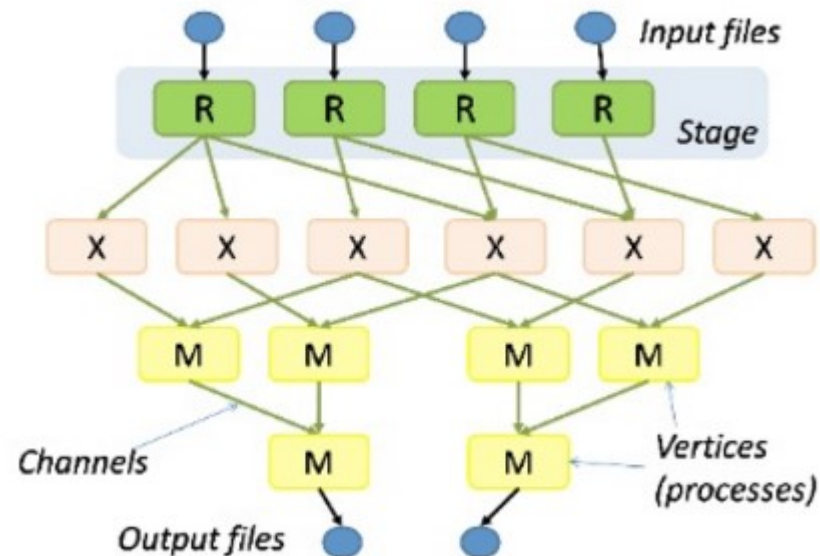
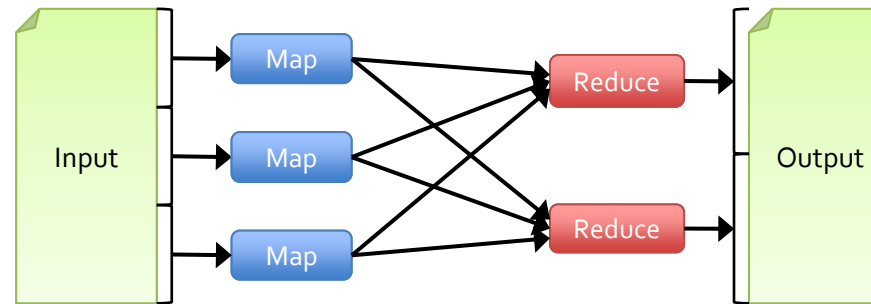
Interface variety

Memory FIFO

Disk

Network

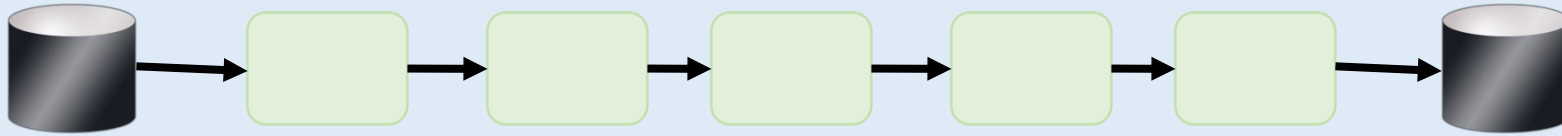
Flexible Modular Composition



Dryad (2007): 2-D Piping

- Unix Pipes: 1-D

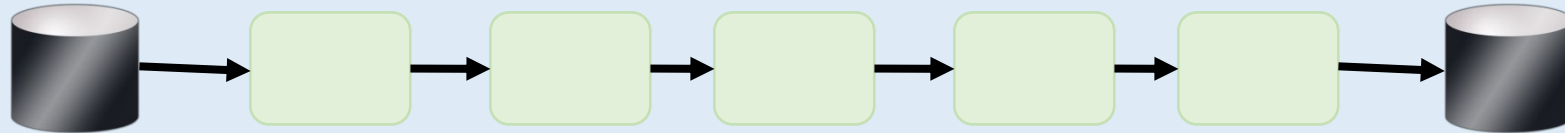
grep | sed | sort | awk | perl



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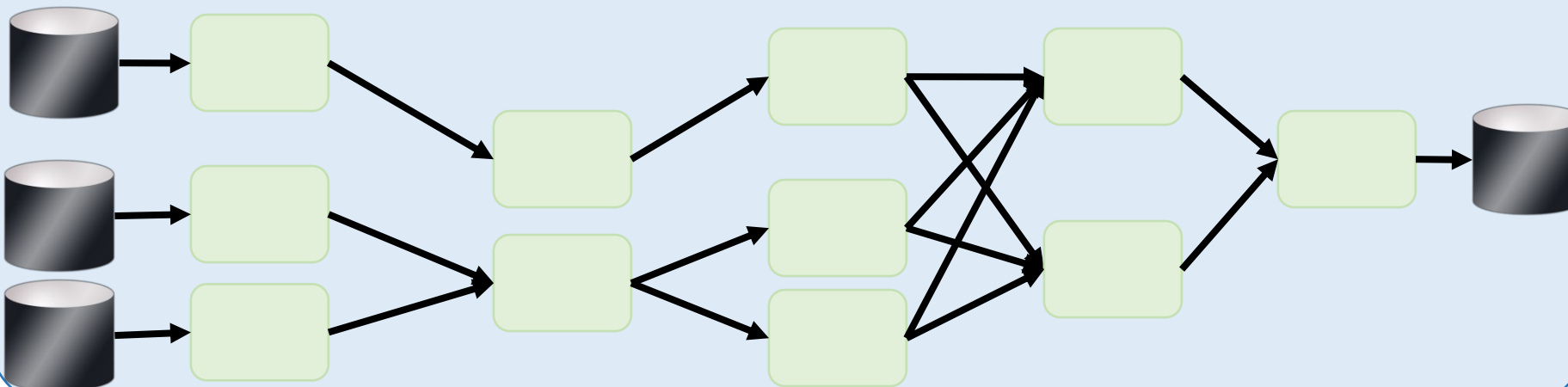
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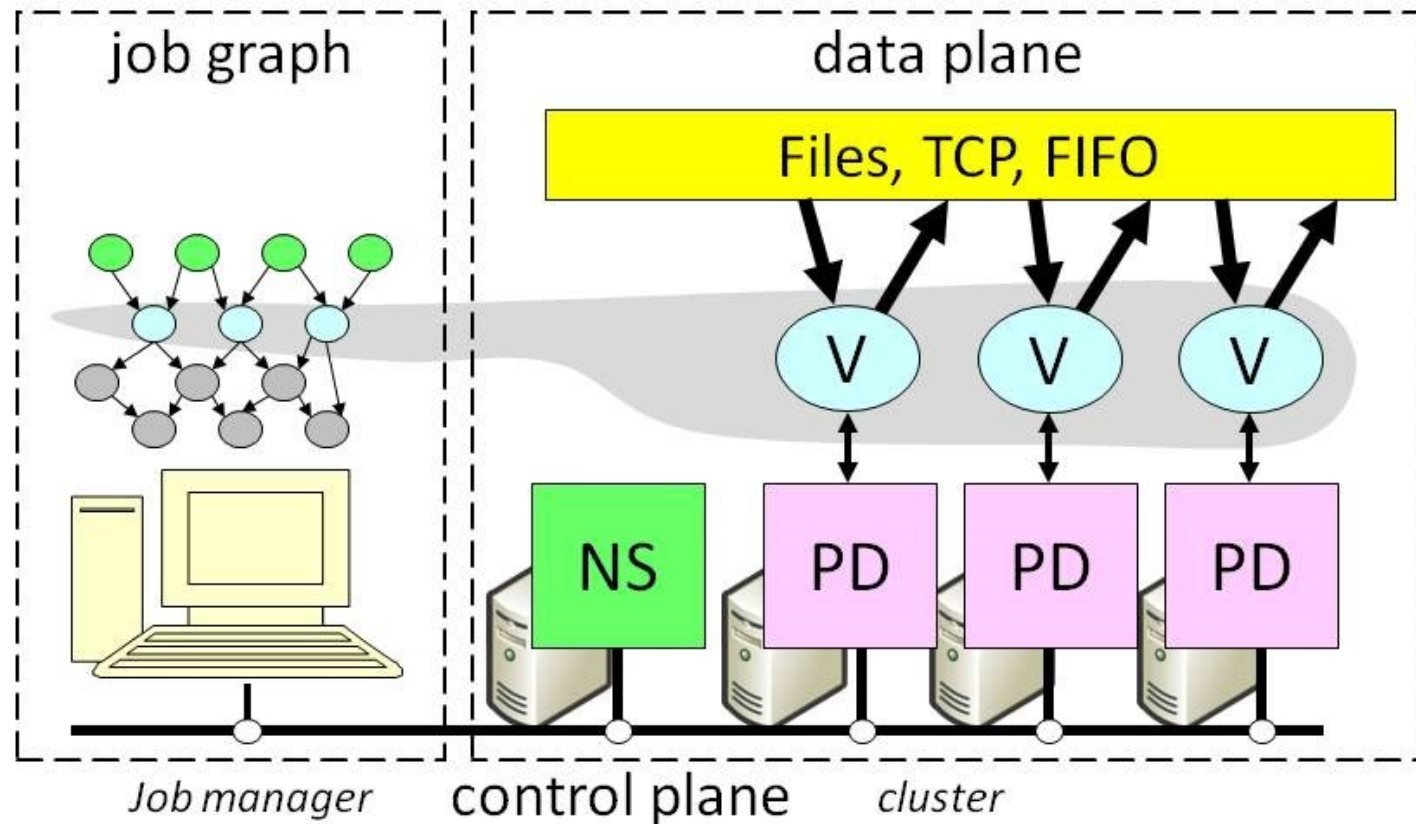


- Dryad: 2-D

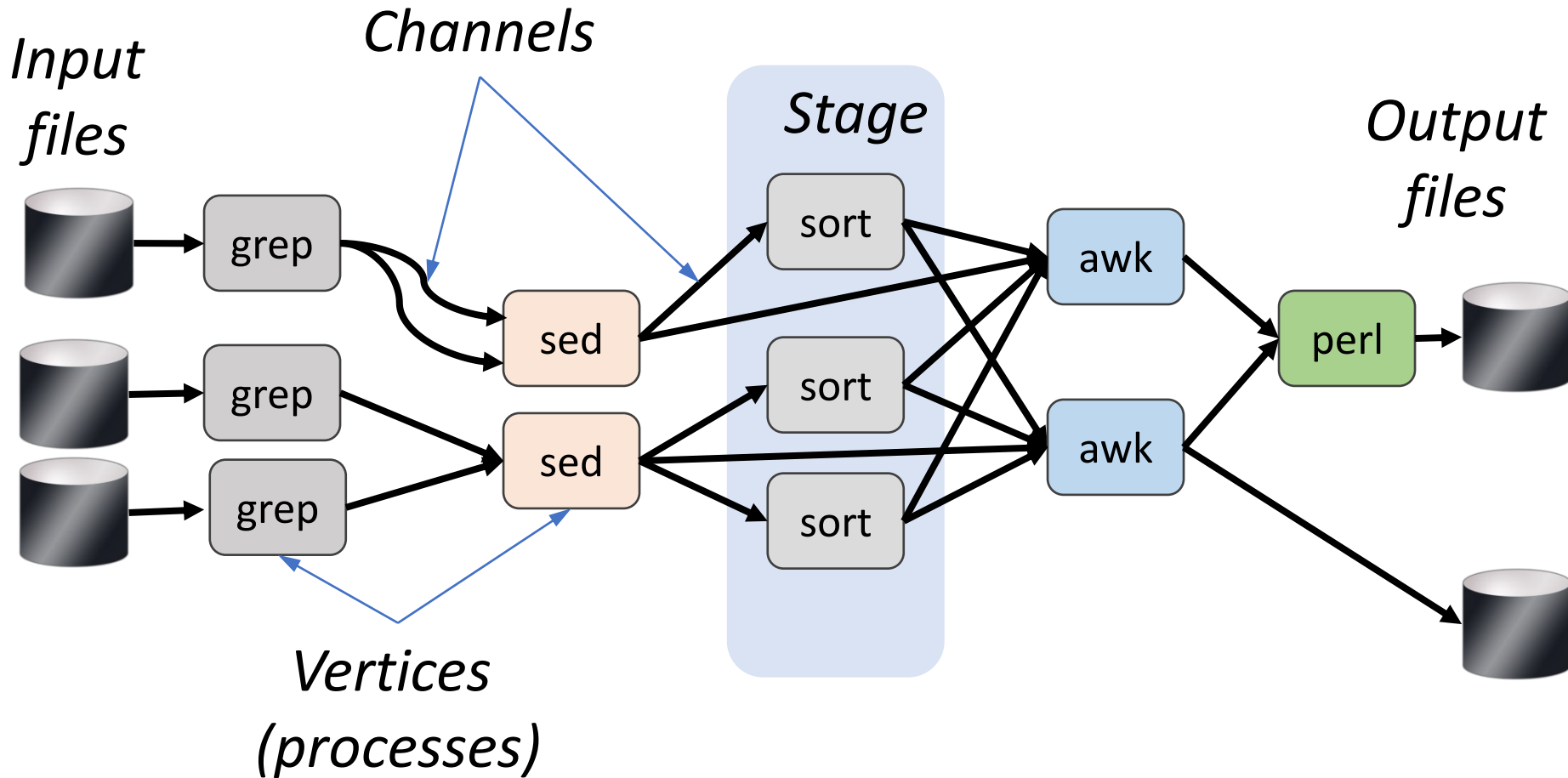
grep¹⁰⁰⁰ | sed⁵⁰⁰ | sort¹⁰⁰⁰ | awk⁵⁰⁰ | perl⁵⁰



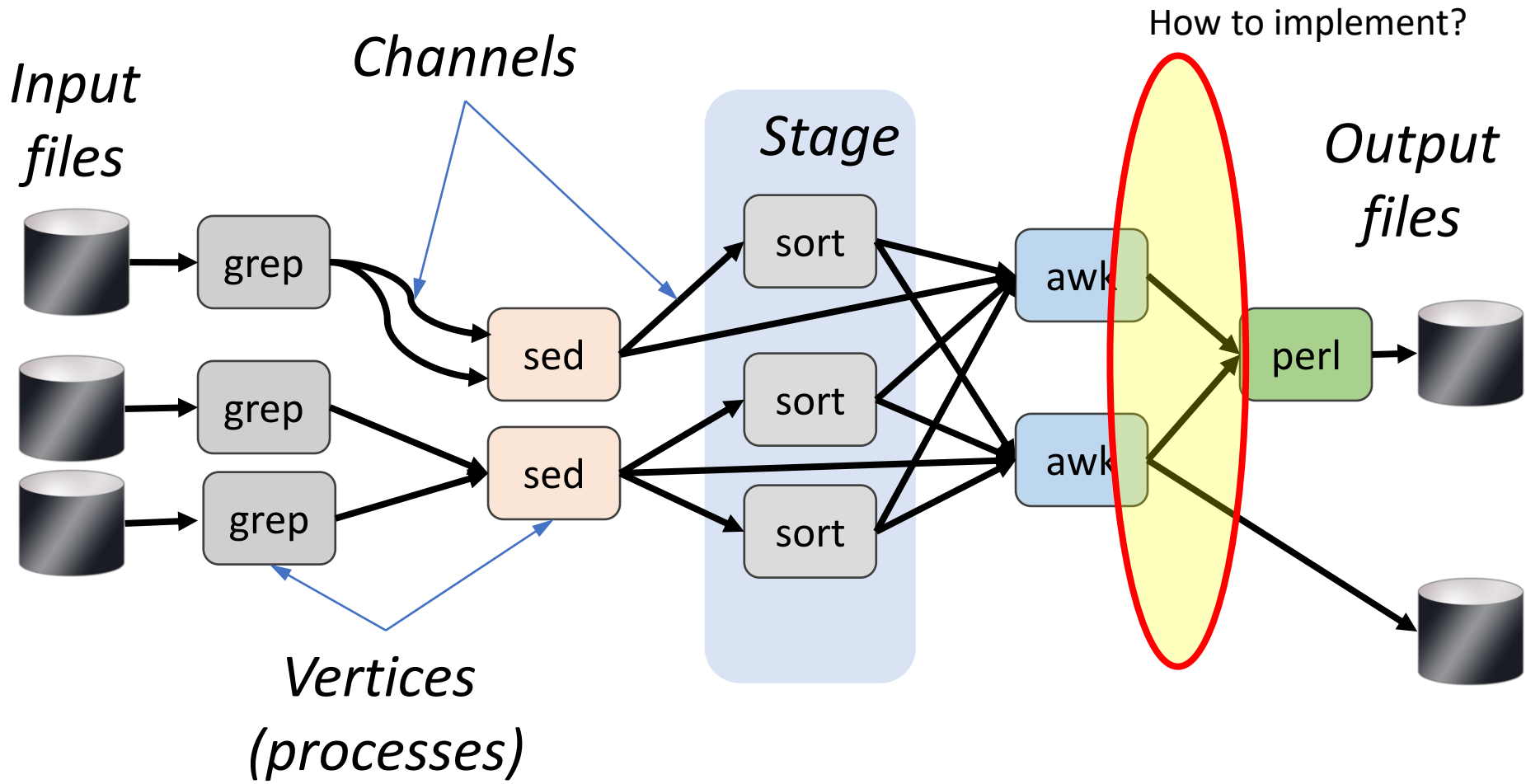
Dataflow Engines



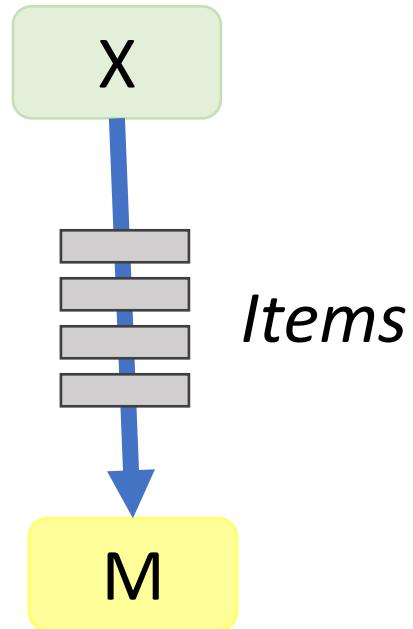
Dataflow Job Structure



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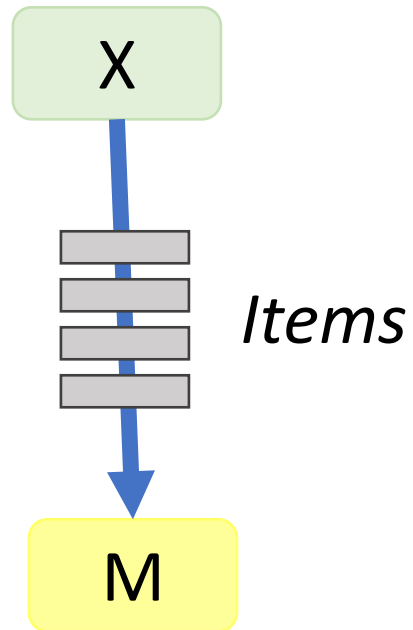
Channels



Finite streams of items

- distributed filesystem files
(persistent)
- SMB/NTFS files
(temporary)
- TCP pipes
(inter-machine)
- memory FIFOs
(intra-machine)

Channels



Key idea:
Encapsulate data movement behind
channel abstraction → gets
programmer out of the picture

Finite streams of items

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Spark (2012) Background

Commodity clusters: important platform

In industry: search, machine translation, ad targeting, ...

In research: bioinformatics, NLP, climate simulation, ...

Cluster-scale models (e.g. MR) de facto standard

Fault tolerance through replicated durable storage

Dataflow is the common theme

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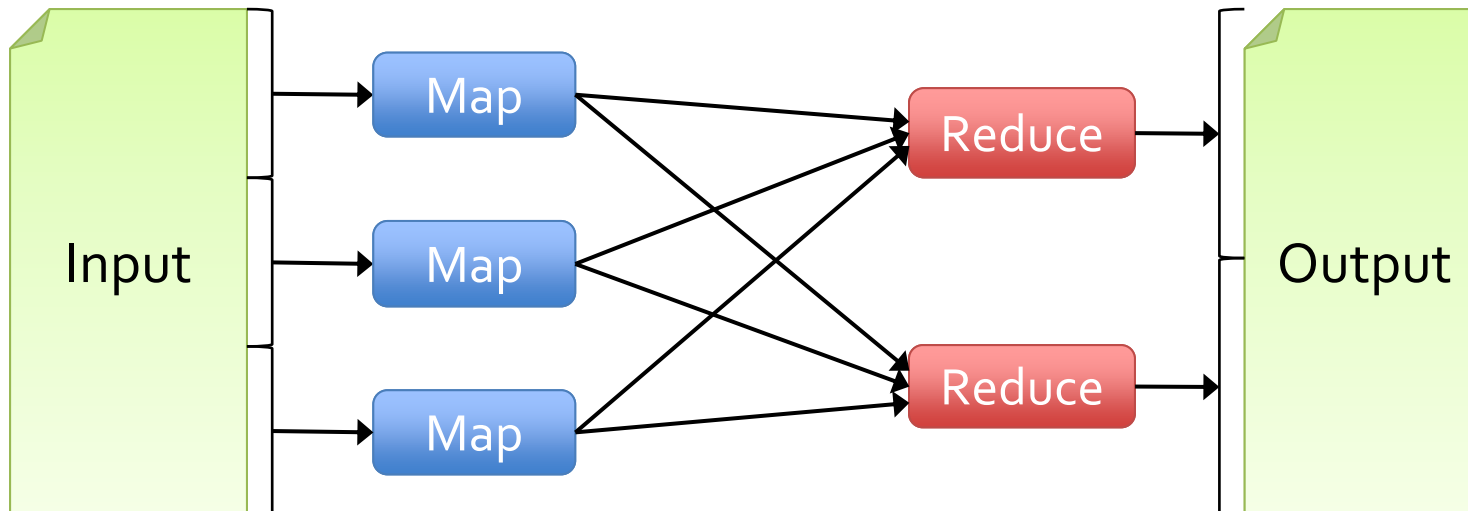
Multi-core

Iteration

Motivation

Programming models for clusters transform data flowing from stable storage to stable storage

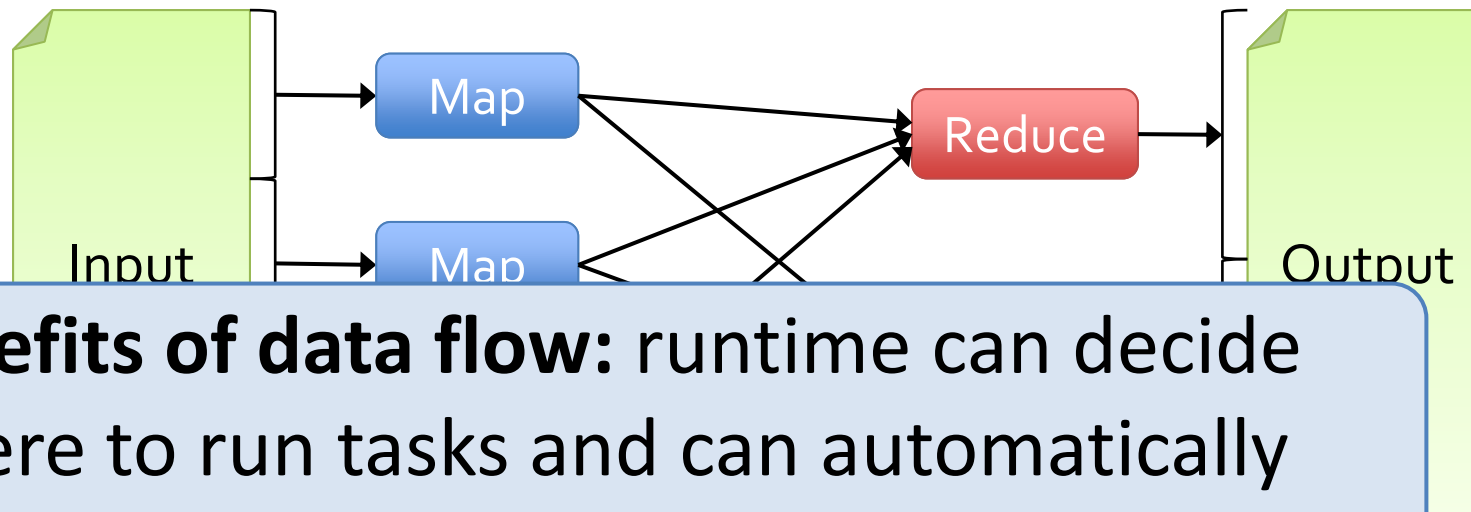
E.g., MapReduce:



Motivation

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Benefits of data flow: runtime can decide where to run tasks and can automatically recover from failures

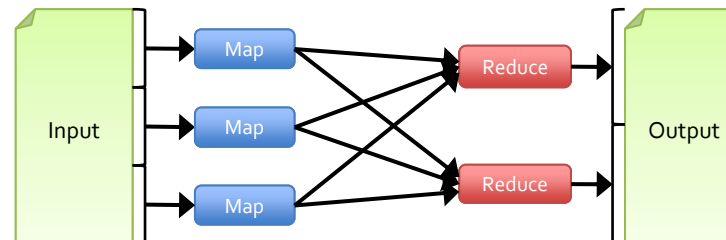
Iterative Computations: PageRank

1. Start each page with a rank of 1
2. On each iteration, update each page's rank to

$$\sum_{i \in \text{neighbors}} \text{rank}_i / |\text{neighbors}_i|$$

```
links = // RDD of (url, neighbors) pairs  
ranks = // RDD of (url, rank) pairs
```

```
for (i <- 1 to ITERATIONS) {  
  ranks = links.join(ranks).flatMap {  
    (url, (links, rank)) =>  
      links.map(dest => (dest, rank/links.size))  
  }.reduceByKey(_ + _)  
}
```



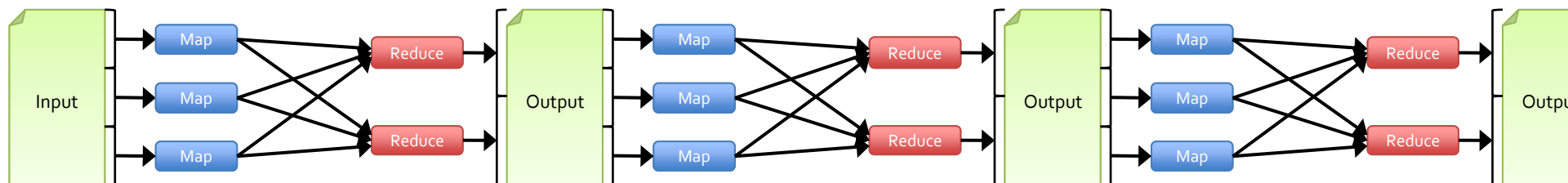
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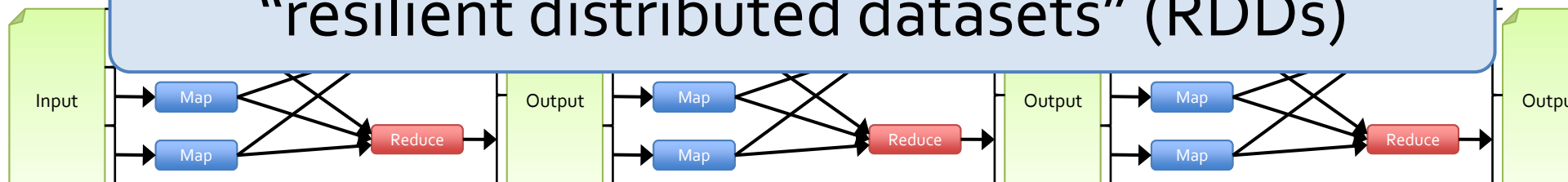
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```

Solution: augment data flow model with
“resilient distributed datasets” (RDDs)



Programming Model

- Resilient distributed datasets (RDDs)
 - Immutable collections partitioned across cluster that can be rebuilt if a partition is lost
 - Created by transforming data in stable storage using data flow operators (map, filter, group-by, ...)
 - Can be *cached* across parallel operations

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 - Reduce, collect, count, save, ...

Programming Model

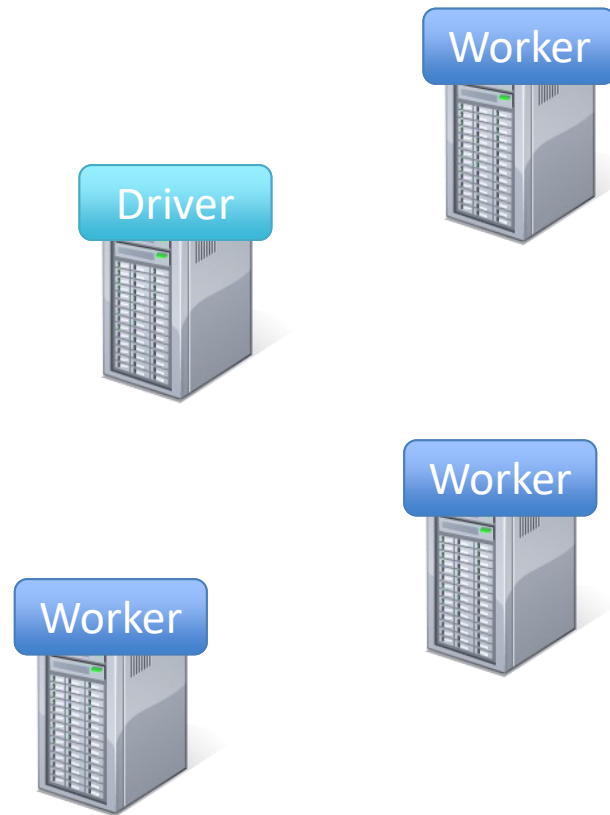
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 - Can be *cached* across parallel operations
- Parallel operations on RDDs
 - Reduce, collect, count, save, ...
- Restricted shared variables
 - Accumulators, broadcast variables

| Example: Log Mining

- Load error messages from a log into memory, then interactively search for various patterns

Example: Log Mining

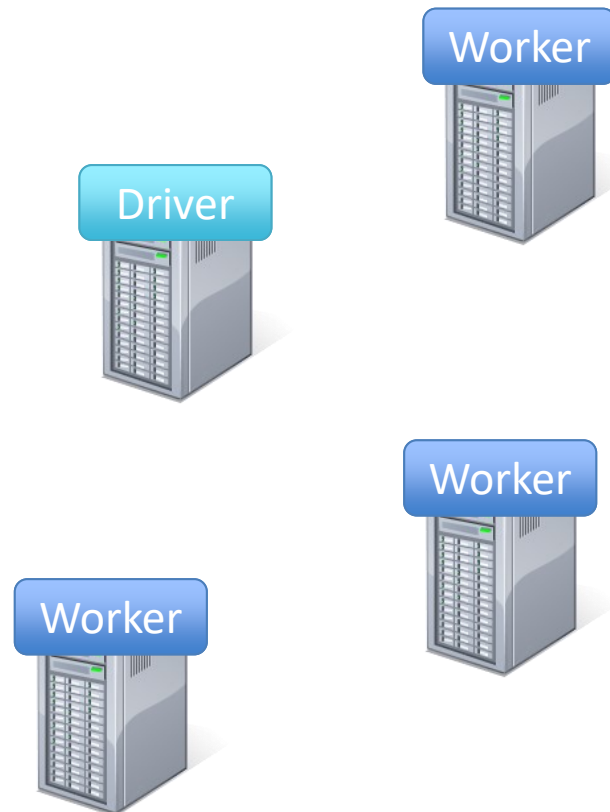
- Load error messages from a log into memory, then interactively search for various patterns



Example: Log Mining

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lines = spark.textFile("hdfs://...")
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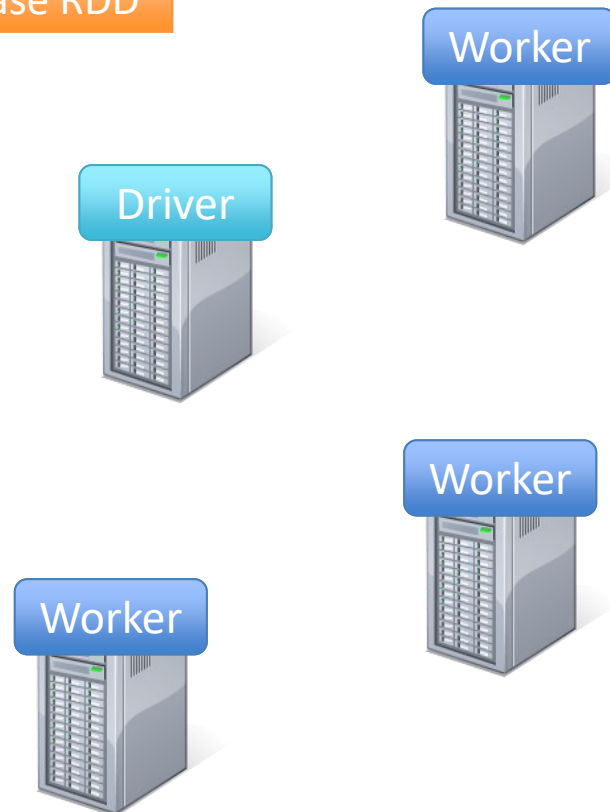


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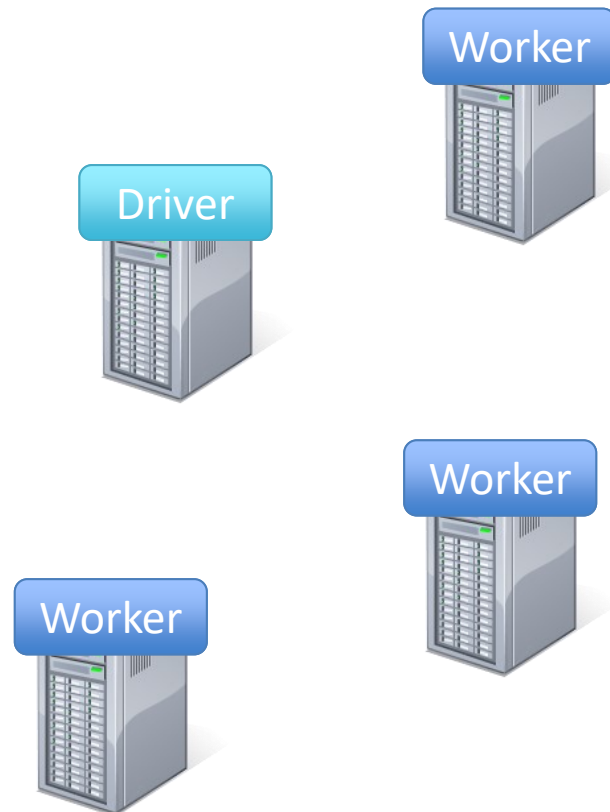
Base RDD



Example: Log Mining

- Load error messages from a log into memory, then interactively search for various patterns

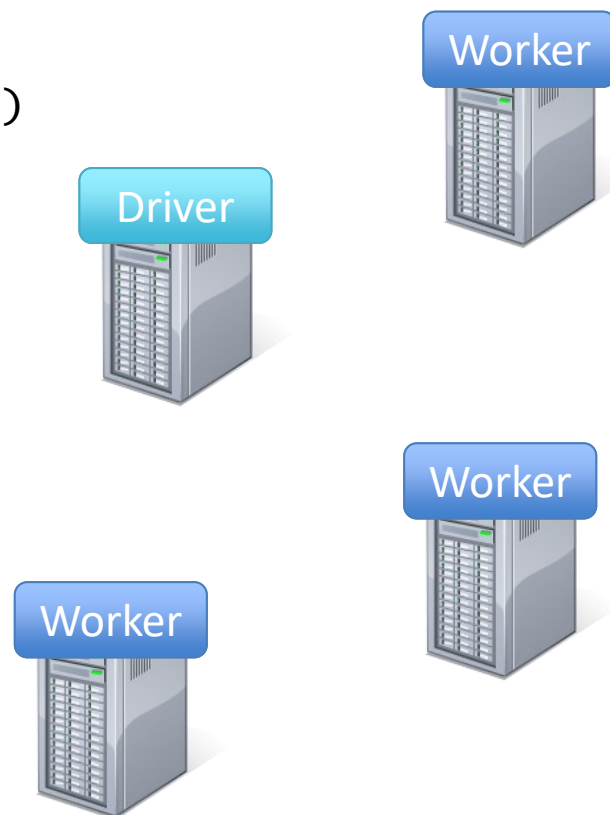
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```



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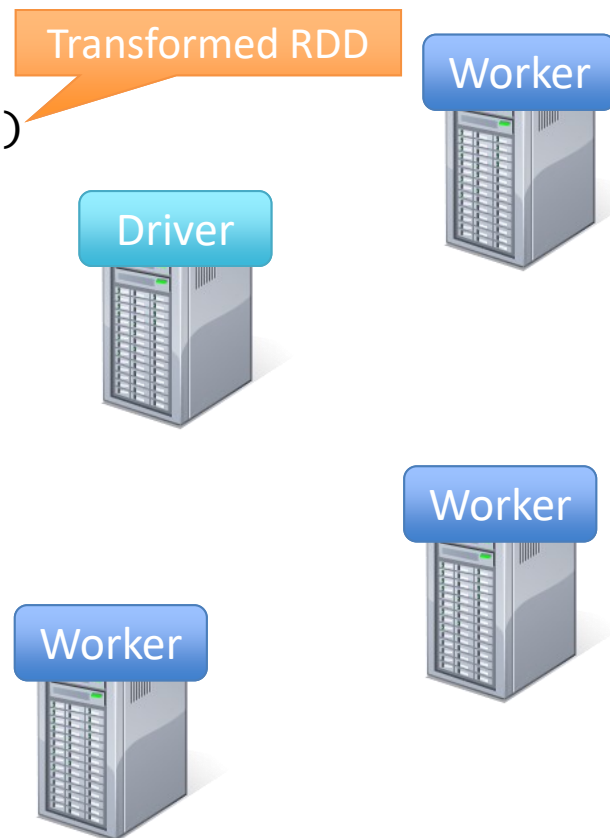
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```



Example: Log Mining

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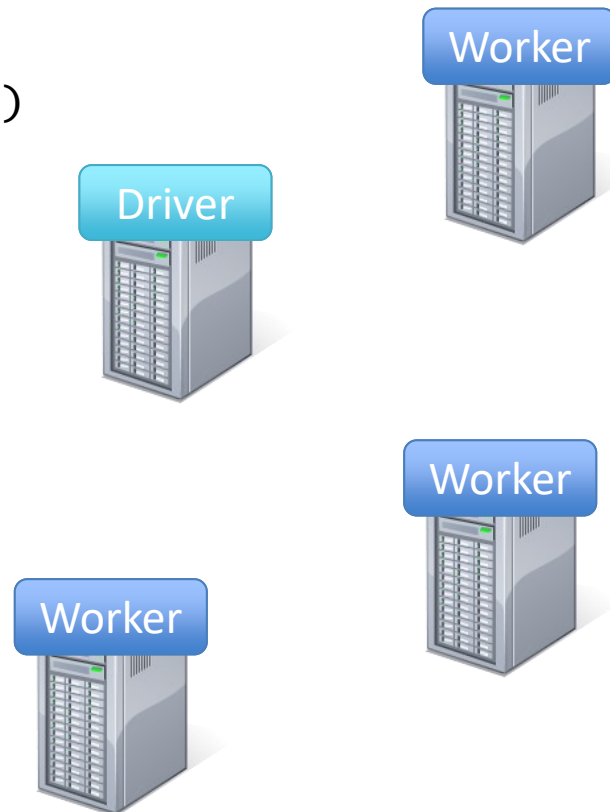
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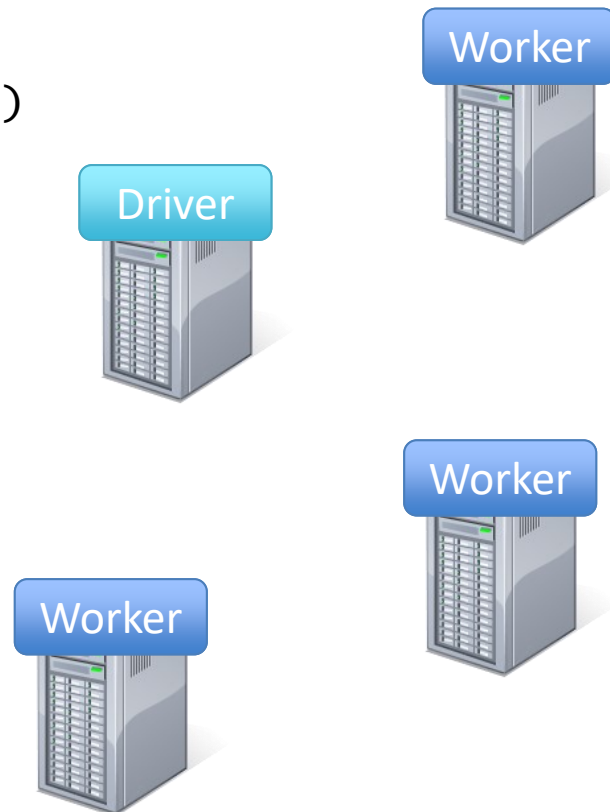
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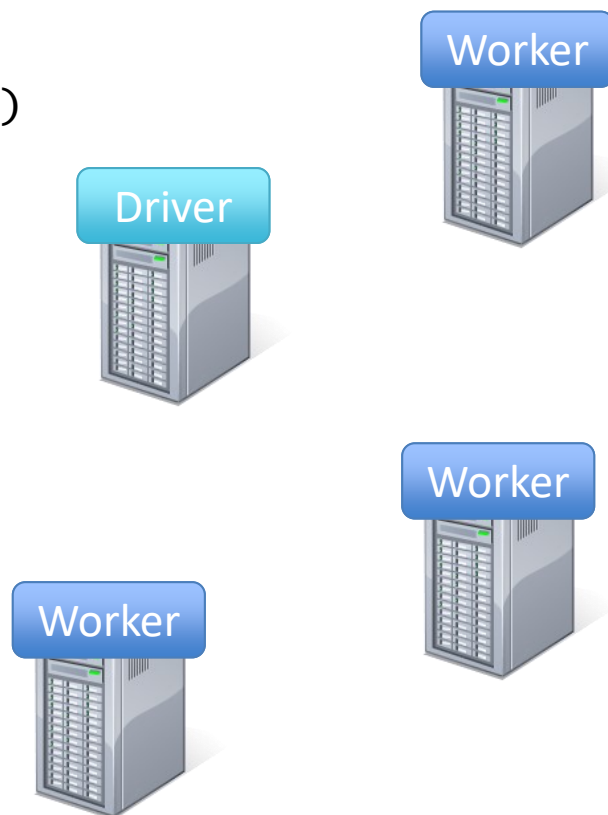
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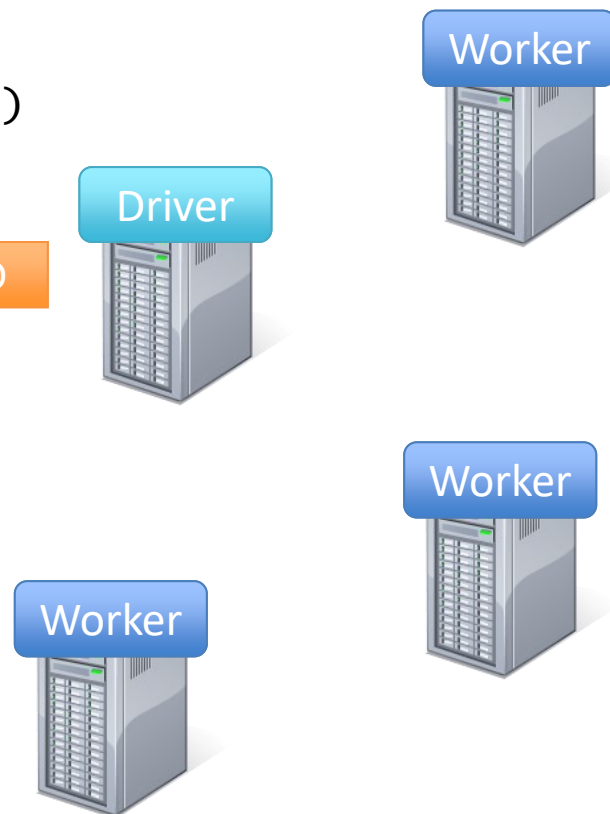


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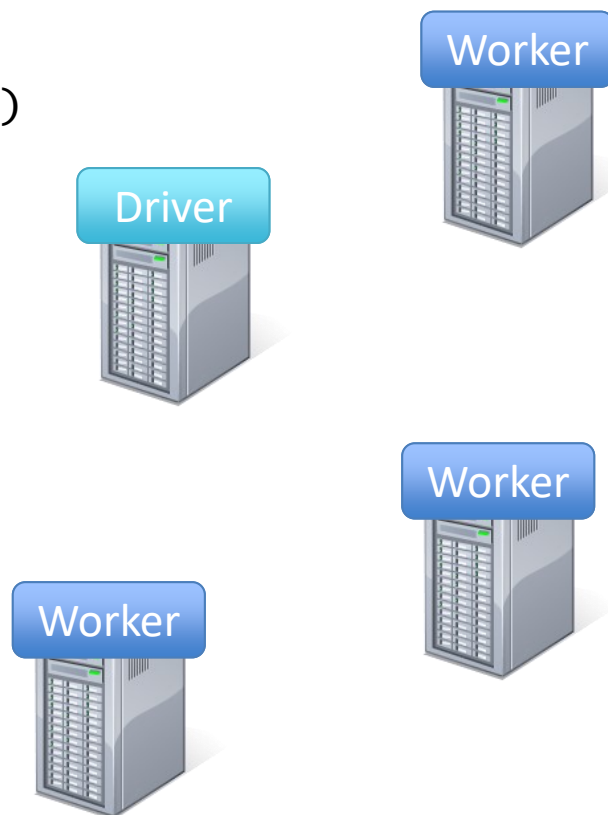
Cached RDD



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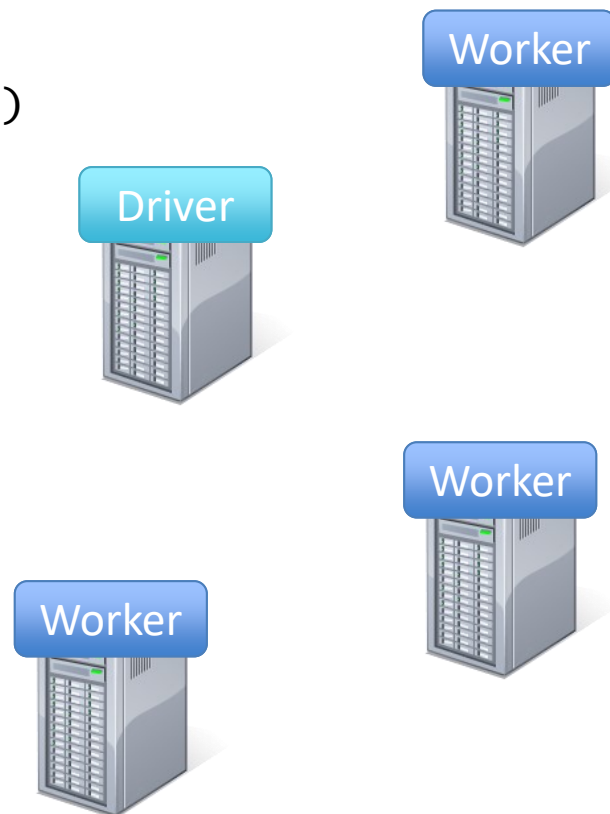


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errors = lines.filter(_.startsWith("ERROR"))
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cachedMsgs = messages.cache()

cachedMsgs.filter(_.contains("foo")).count
```

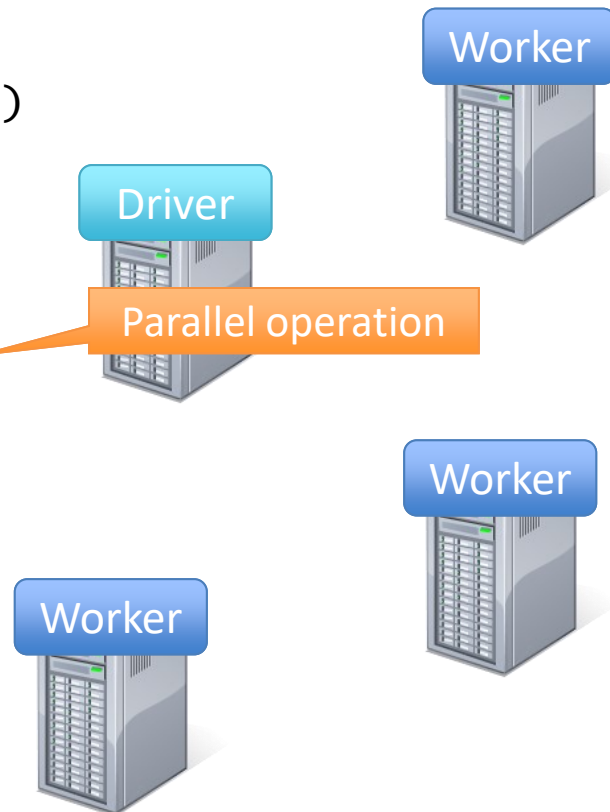


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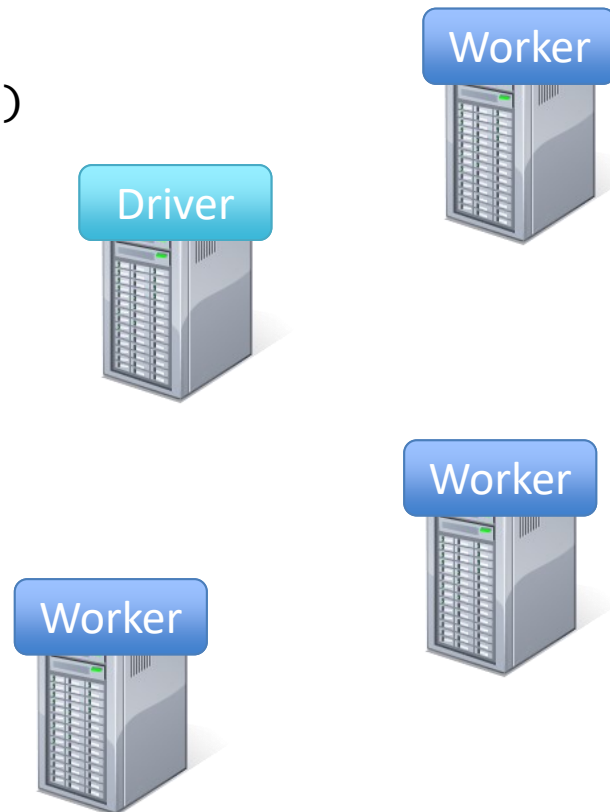


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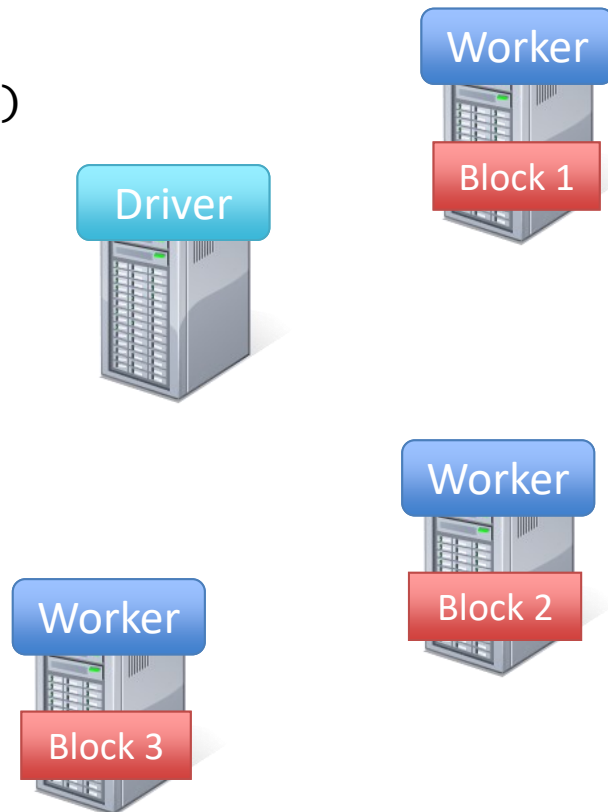
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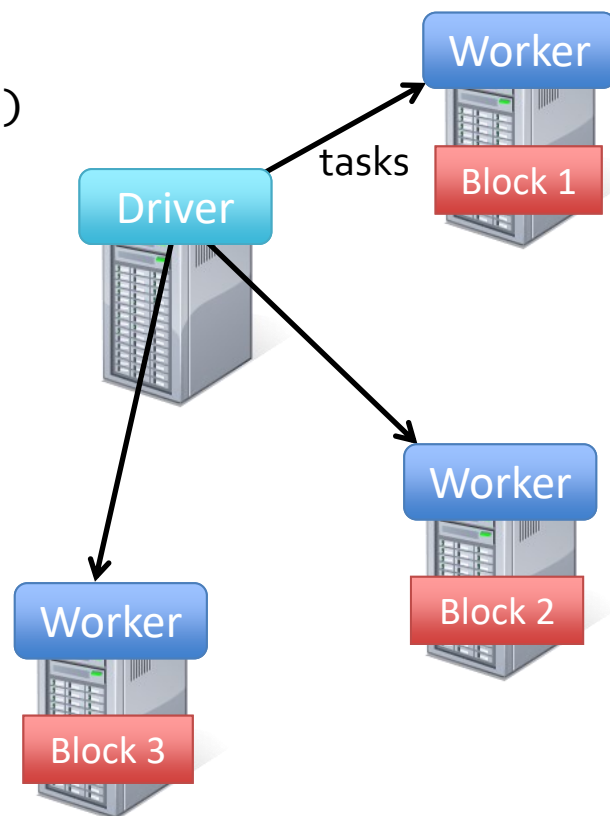


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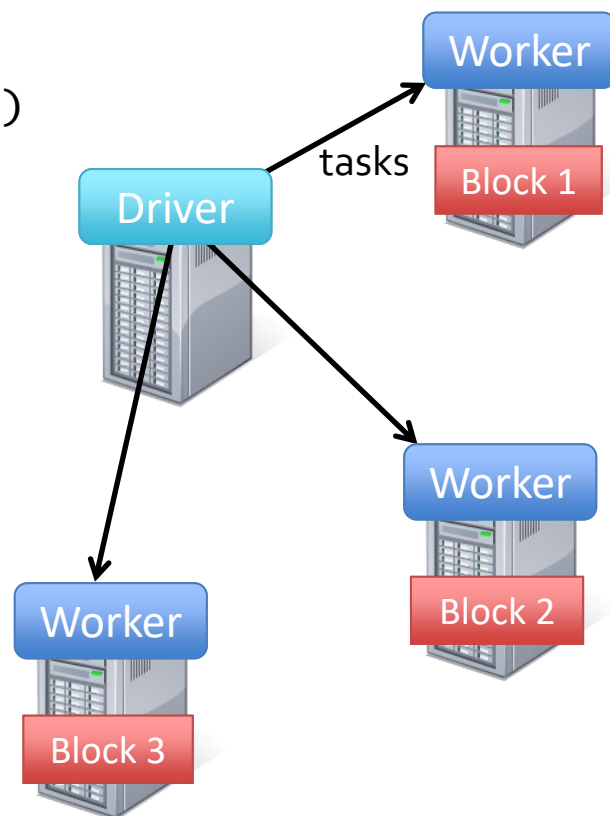


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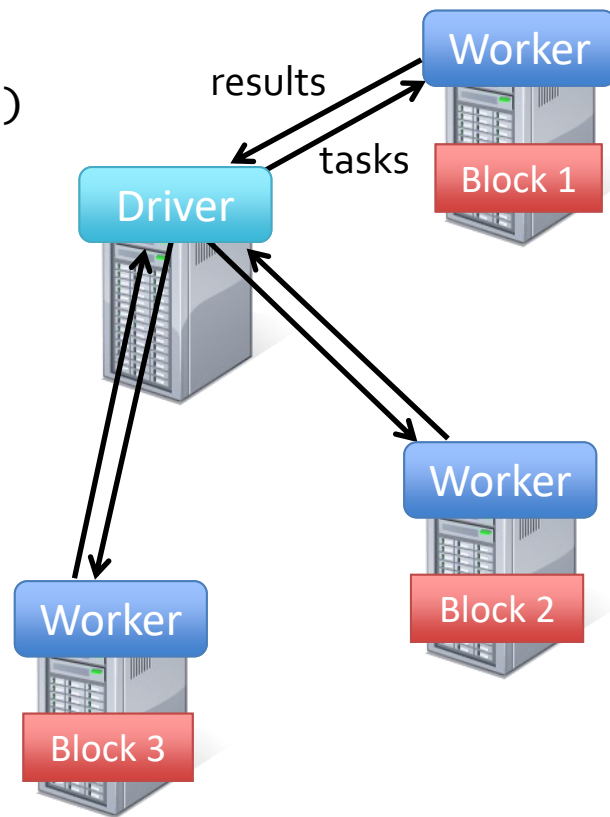


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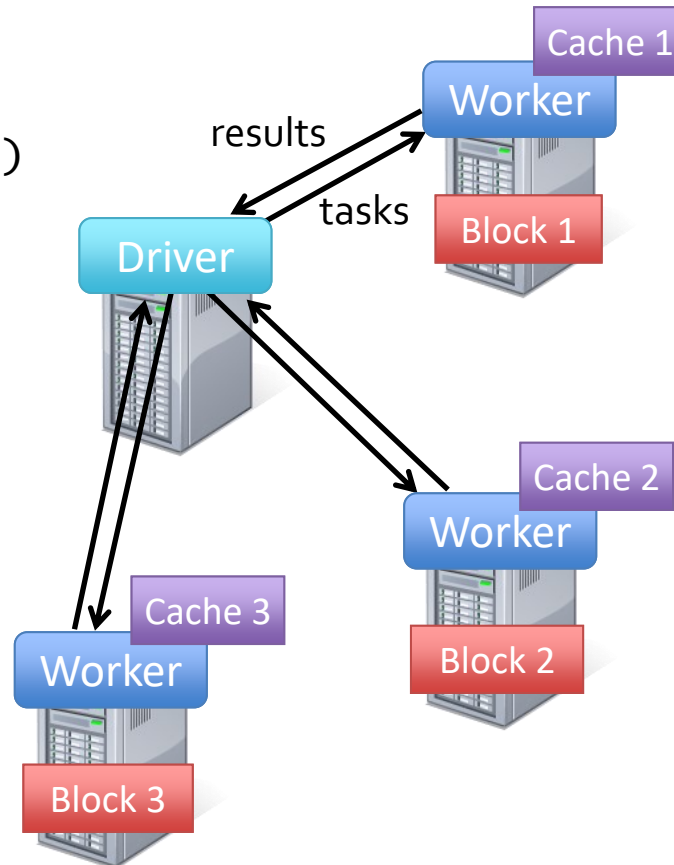


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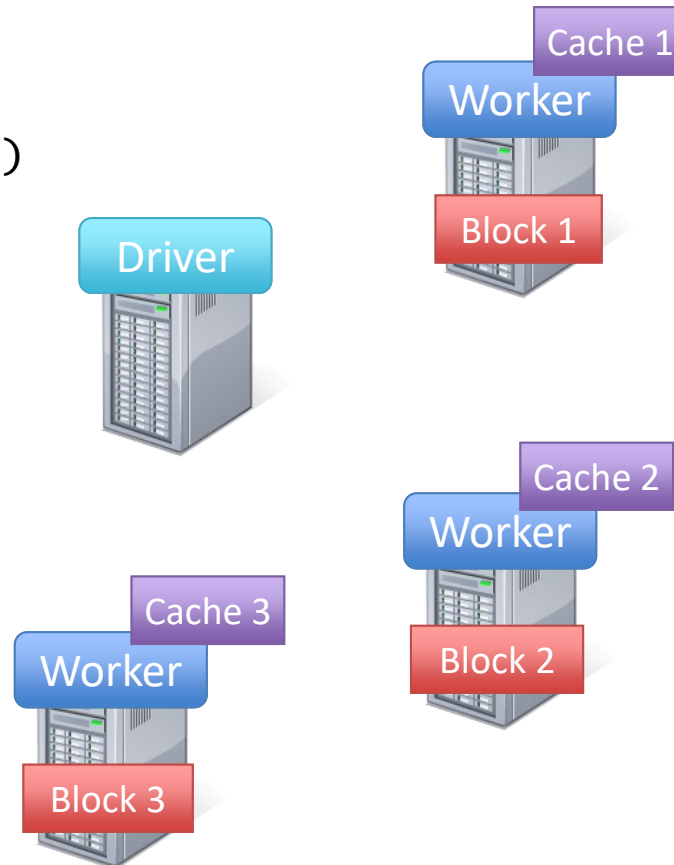


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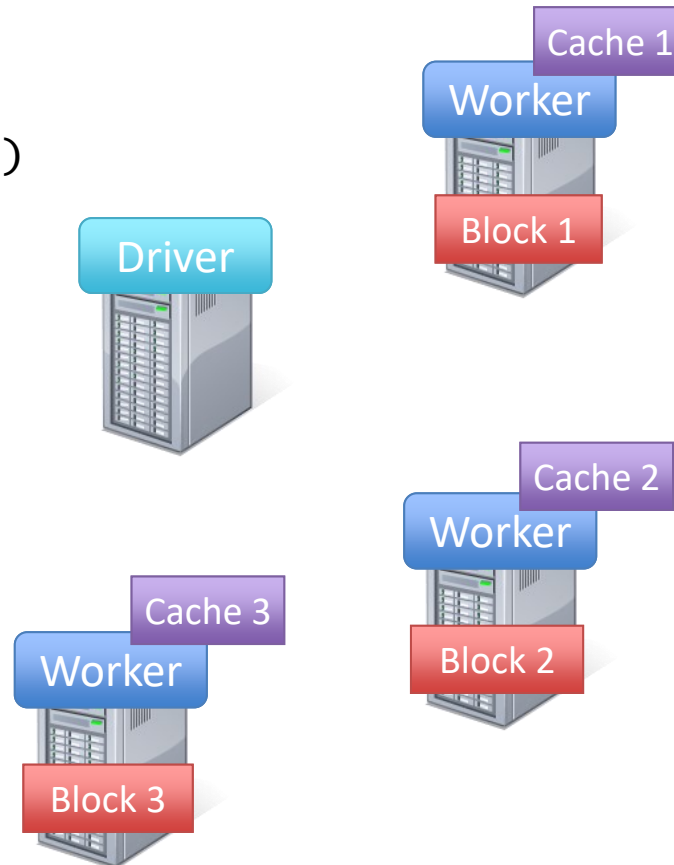


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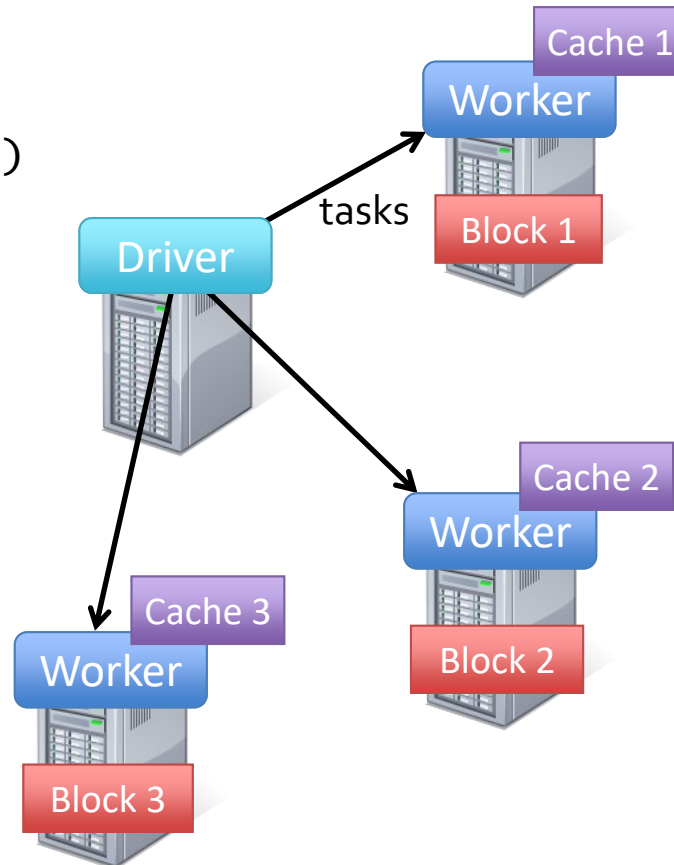


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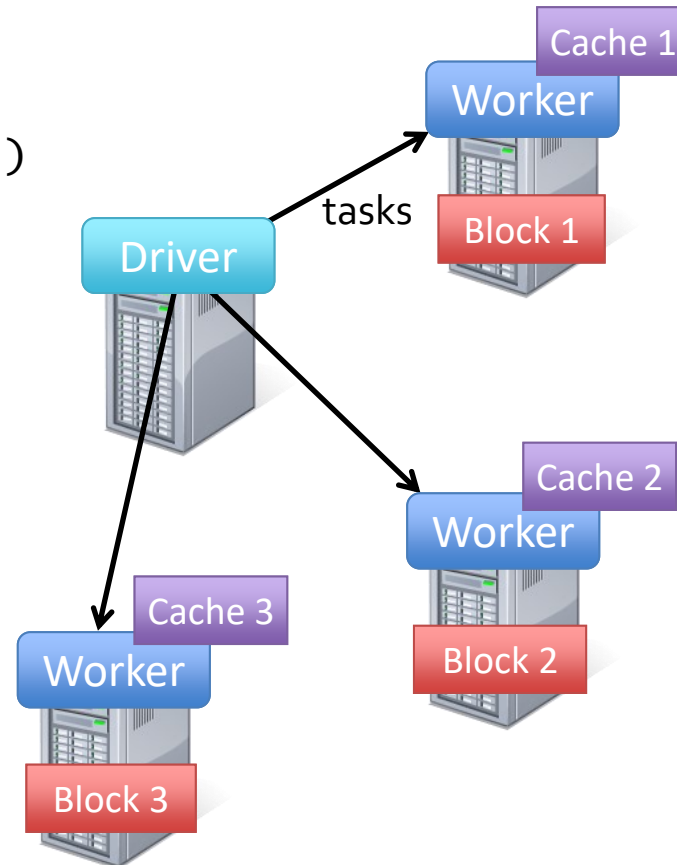


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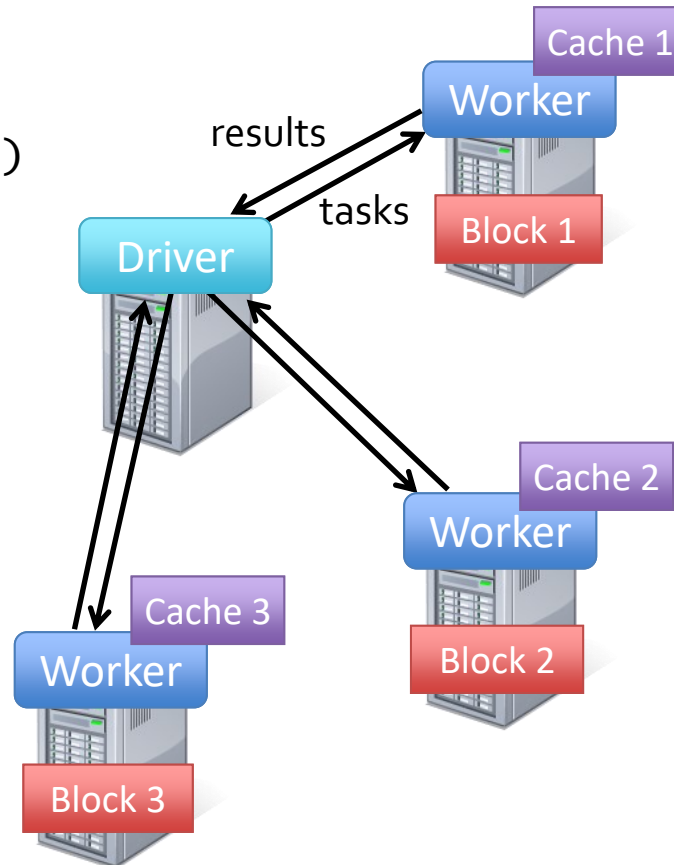


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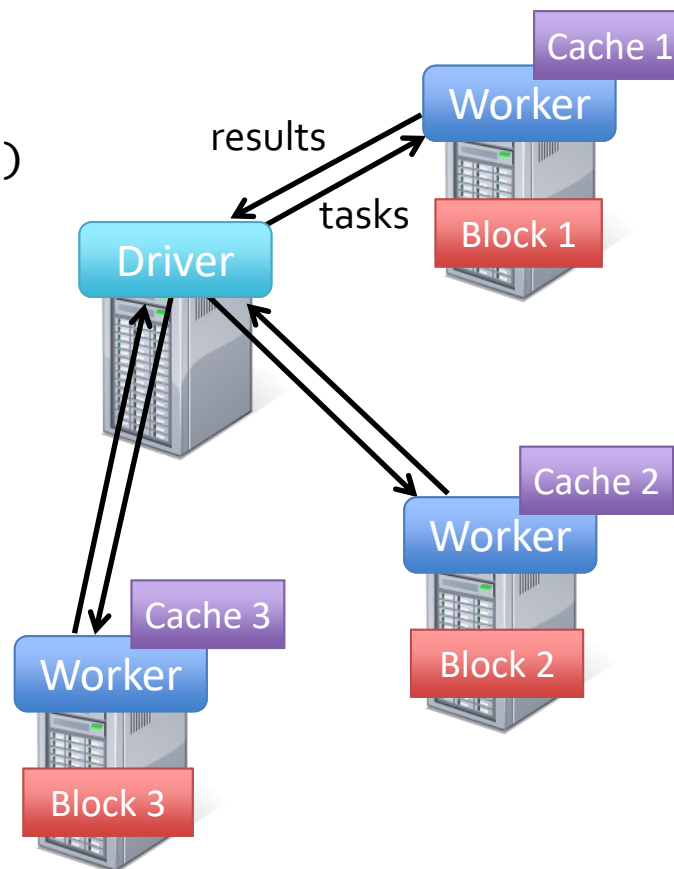


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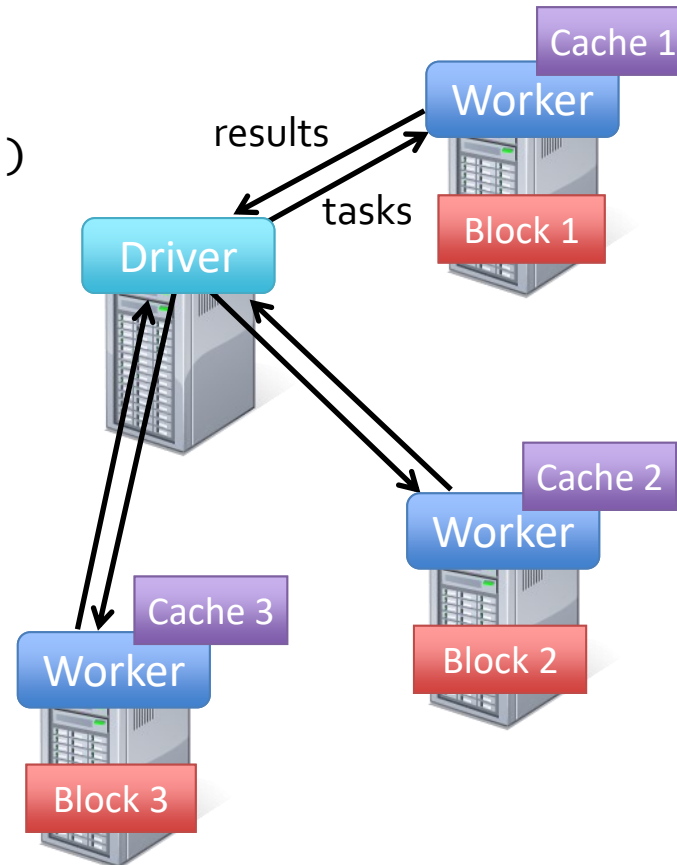
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. . .
```

Result: full-text search of Wikipedia in <1 sec (vs 20 sec for on-disk data)

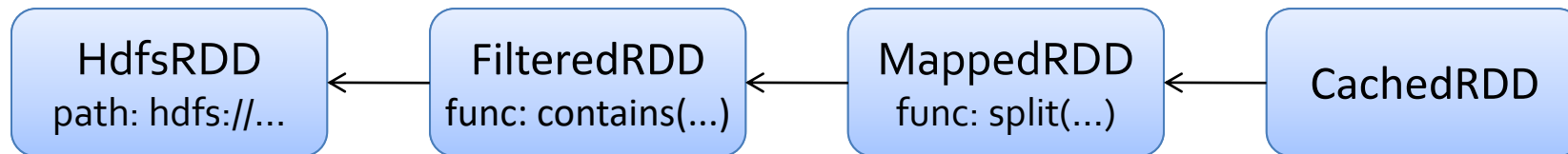


RDD Fault Tolerance

- RDDs maintain *lineage* information that can be used to reconstruct lost partitions

- Ex:

```
cachedMsgs = textFile(...).filter(_.contains("error"))  
                        .map(_.split('\t')(2))  
                        .persist()
```



Data-Parallel Computation Systems

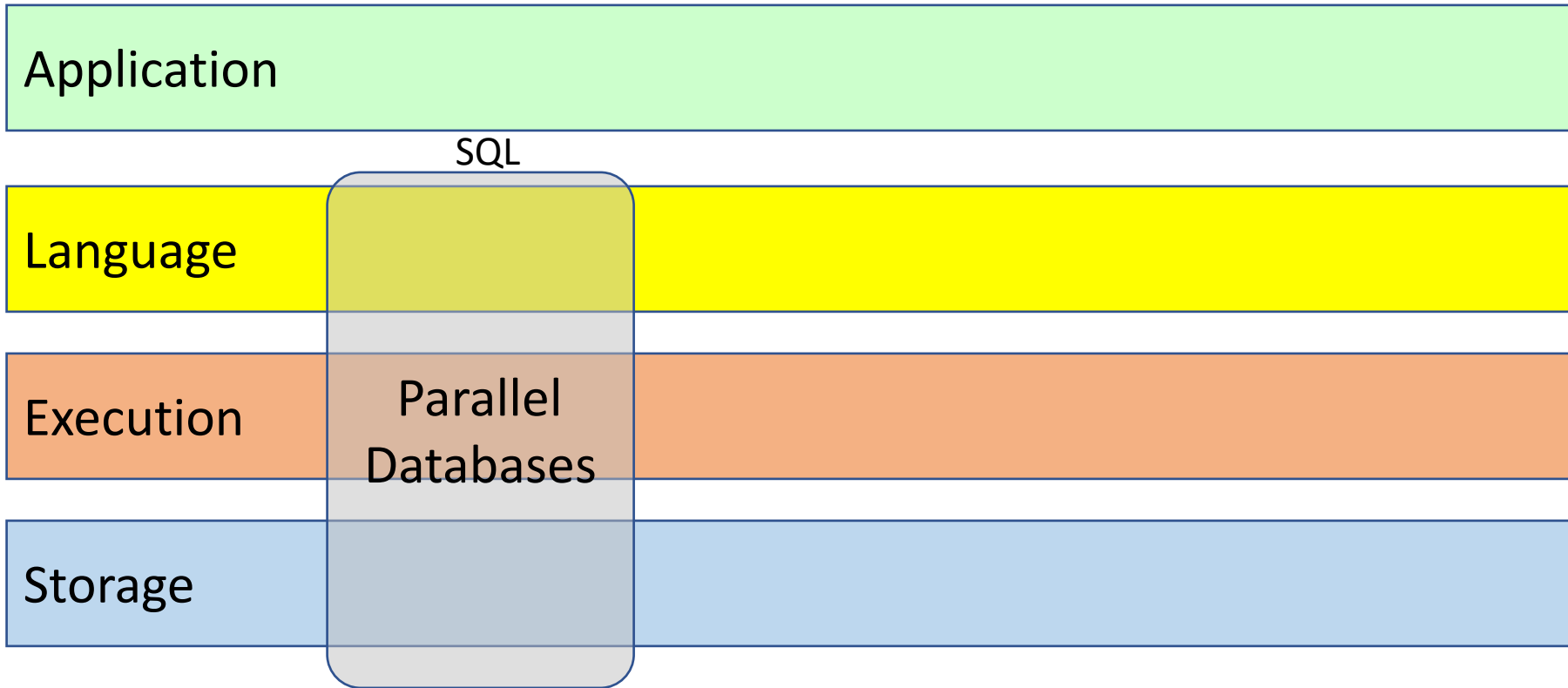
Application

Language

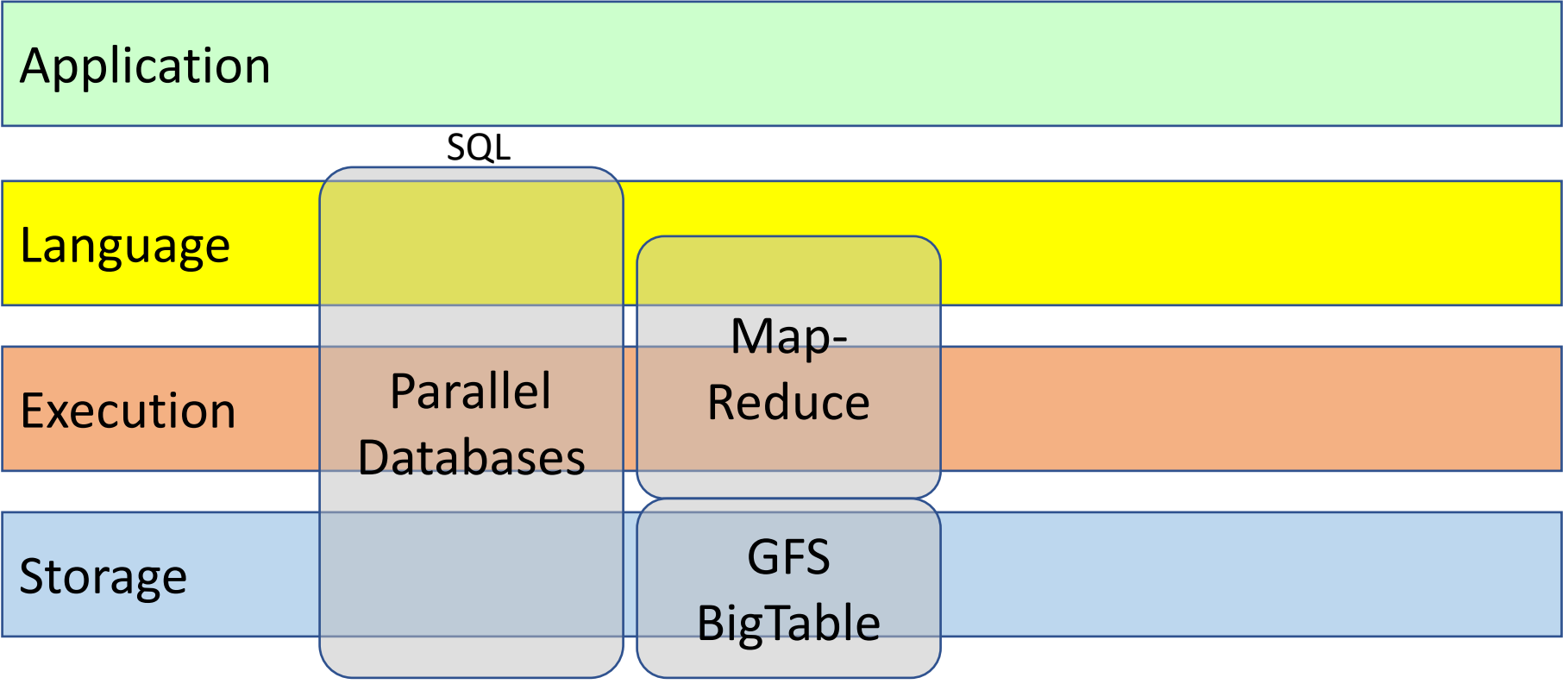
Execution

Storage

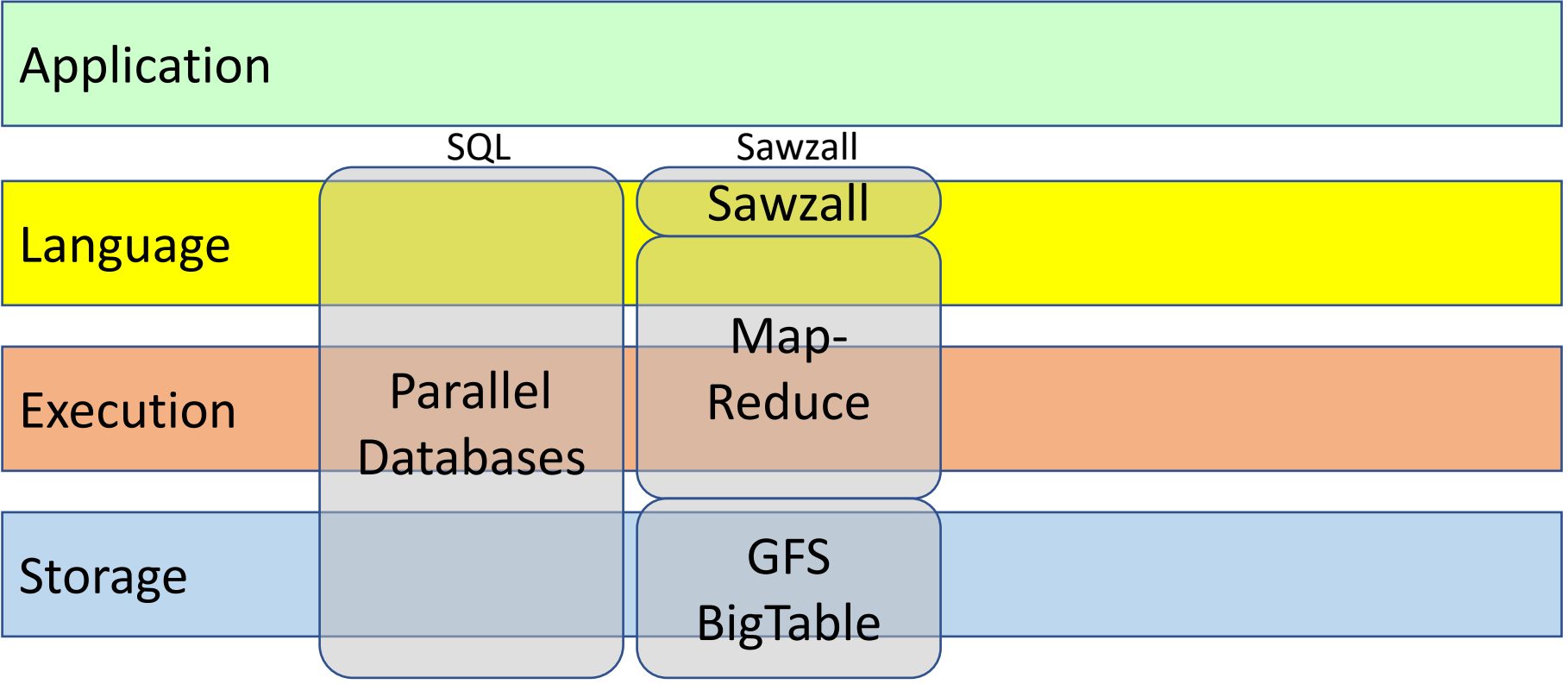
Data-Parallel Computation Systems



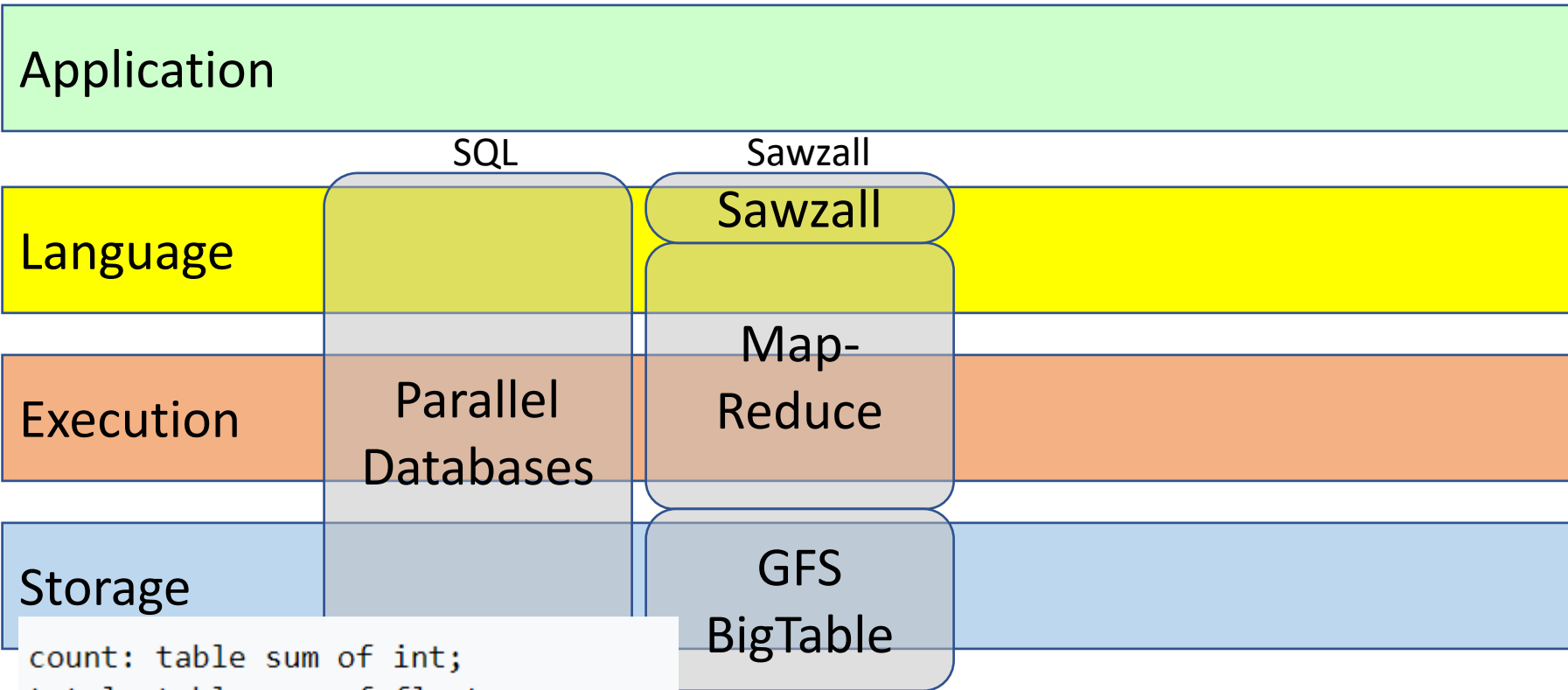
Data-Parallel Computation Systems



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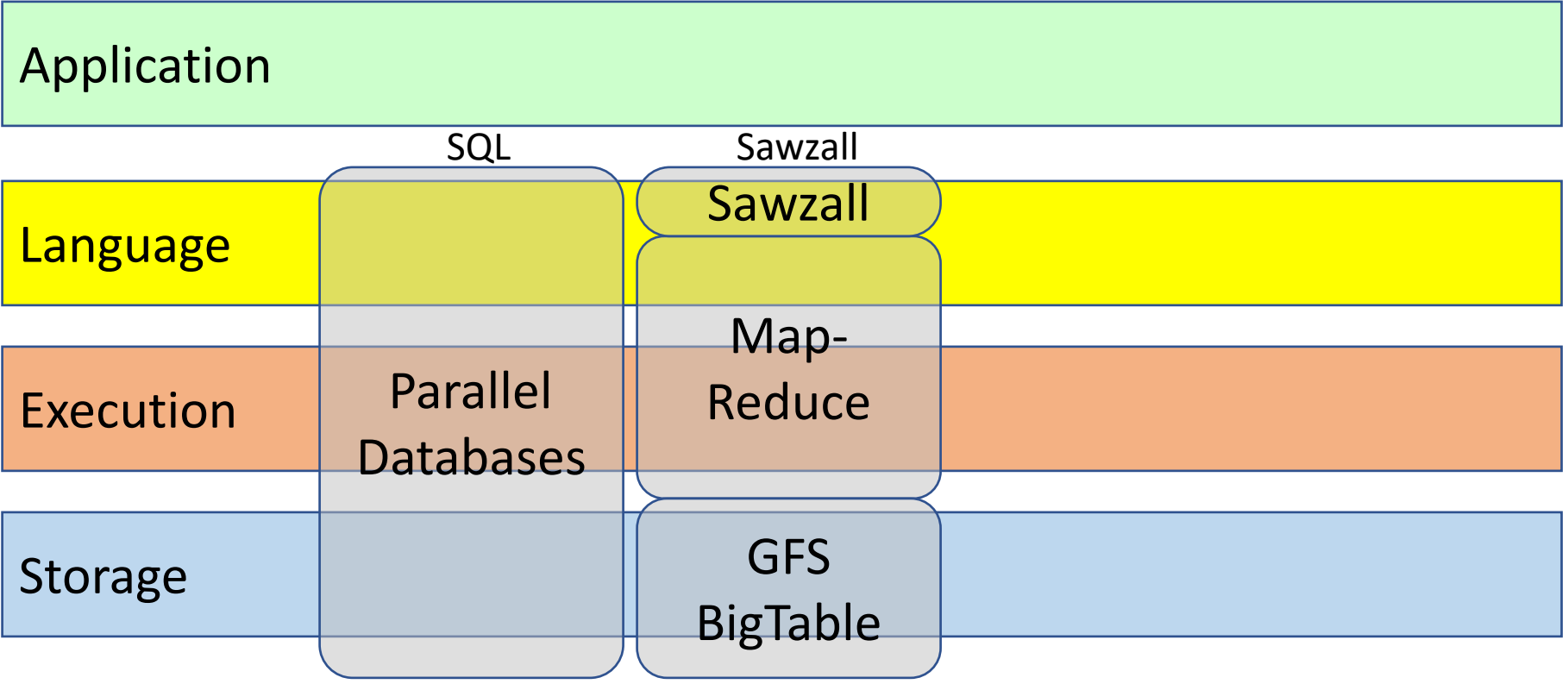


Data-Parallel Computation Systems

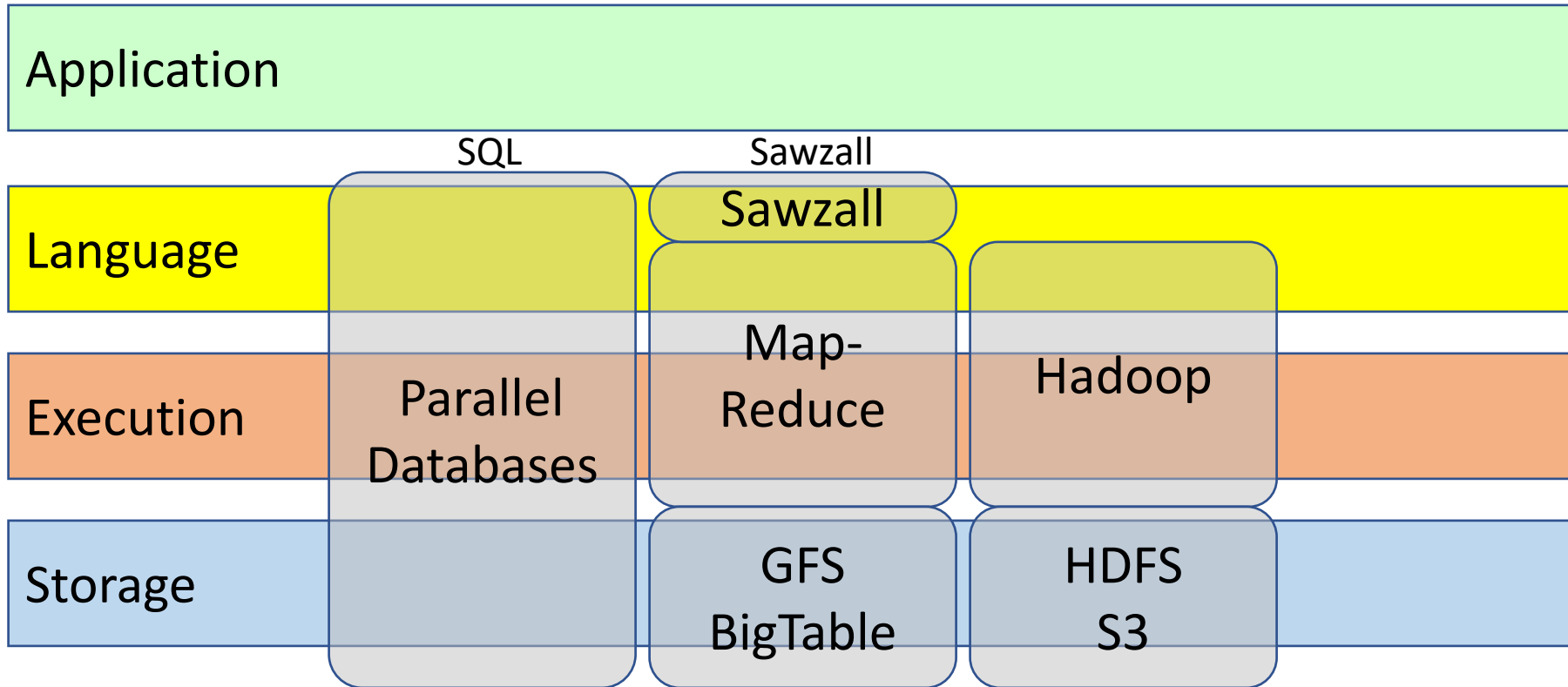


```
count: table sum of int;  
total: table sum of float;  
sum_of_squares: table sum of float;  
x: float = input;  
emit count <- 1;  
emit total <- x;  
emit sum_of_squares <- x * x;
```

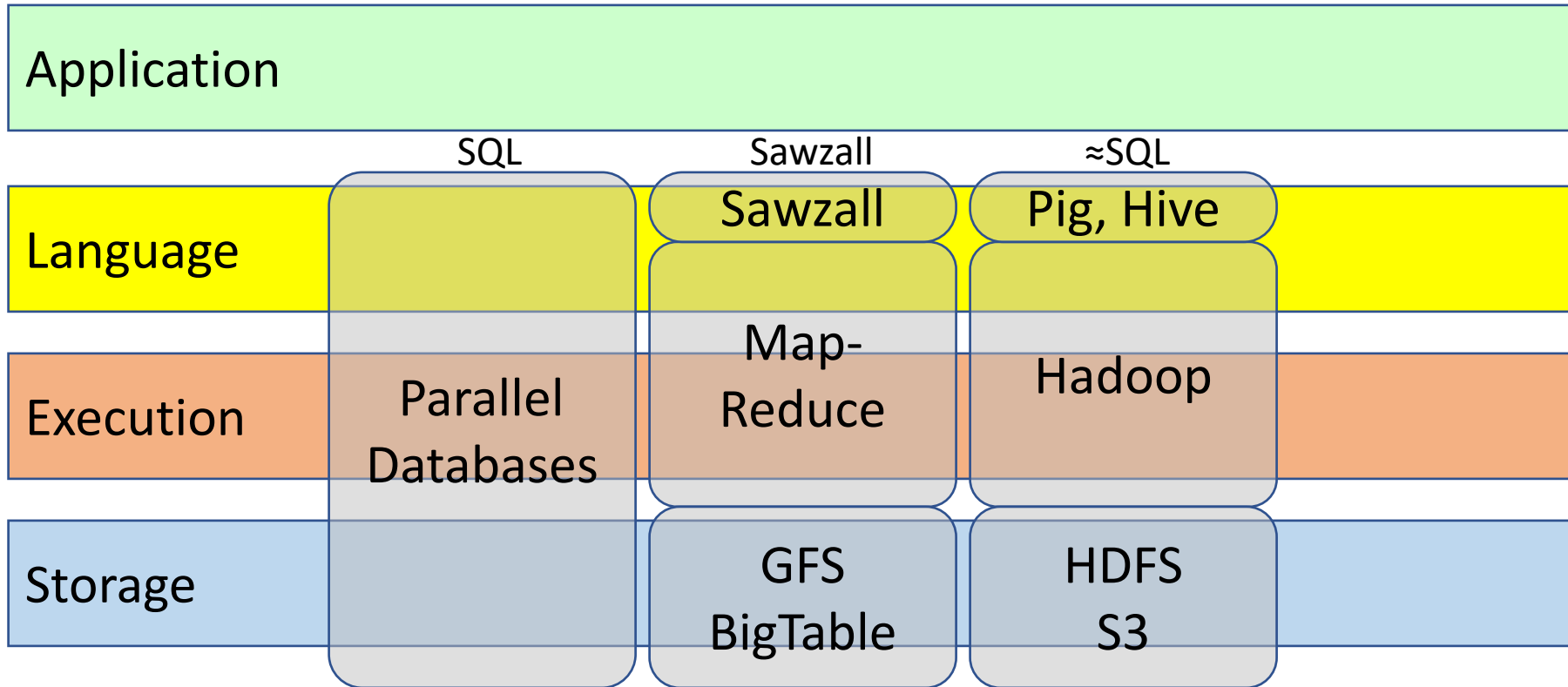
Data-Parallel Computation Systems



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Data-Parallel Computation Systems



Systems

```
lines = LOAD '/user/hadoop/HDFS_File.txt' AS (line:chararray);  
words = FOREACH lines GENERATE FLATTEN(TOKENIZE(line)) as word;  
grouped = GROUP words BY word;  
wordcount = FOREACH grouped GENERATE group, COUNT(words);  
DUMP wordcount;
```

App

SQL

Sawzall

≈SQL

Language

Sawzall

Pig, Hive

Execution

Parallel
Databases

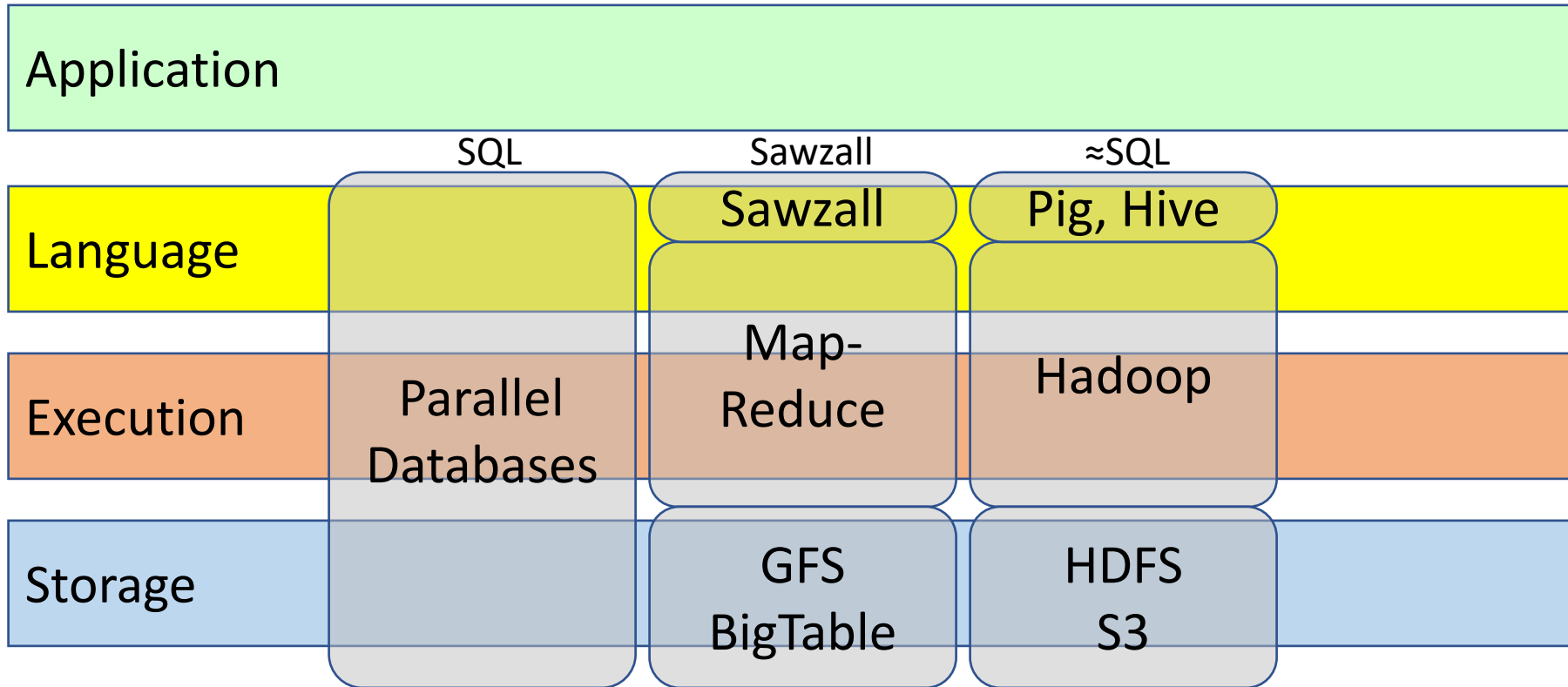
```
-- import the file as lines  
CREATE EXTERNAL TABLE lines(line string)  
LOAD DATA INPATH 'books' OVERWRITE INTO TABLE lines;  
  
-- create a virtual view that splits the lines  
SELECT word, count(*) FROM lines  
  
LATERAL VIEW explode(split(text, ' ')) lTable as word  
GROUP BY word;
```

Storage

BigTable

S3

Data-Parallel Computation Systems



Data-Parallel Computation Systems

Application	SQL	Sawzall	≈SQL	
Language		Sawzall	Pig, Hive	DryadLINQ Scope
Execution	Parallel Databases	Map- Reduce	Hadoop	Dryad
Storage		GFS BigTable	HDFS S3	Cosmos Azure SQL Server

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Summary

Dataflow key enabler for cluster-scale parallelism

Key issues become runtime's responsibility

- Data movement

- Scheduling

- Fault-tolerance

Example: Counting Words...

```
map(String input_key, String input_value):  
    // input_key: document name  
    // input_value: document contents  
    for each word w in input value:  
        EmitIntermediate(w, "1");  
  
reduce(String output_key,  
        Iterator intermediate_values):  
    // output_key: a word  
    // output_values: a list of counts  
    int result = 0;  
    for each v in intermediate_values:  
        result += ParseInt(v);  
    Emit(AsString(result));
```

MapReduce handles all the other details!

Redundant Execution

Slow worker can throttle performance: why?

What makes a worker slow?

Other Jobs on machine (how could we fix)

Bad disks, soft errors

Exotica (processor caches disabled!)

Solution: spawn backups near end of phase

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Materializing splits

$N=1000$ mappers \rightarrow $M=500$ keys = 500,000 local files

500 reducer instances “pull” these files

DBMSs push splits to sockets (no local temp files)

MapReduce: !novel && feature-poor

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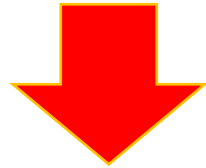
- Indexing
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- Integrity constraints, referential integrity
- Views

DryadLINQ = LINQ + Dryad

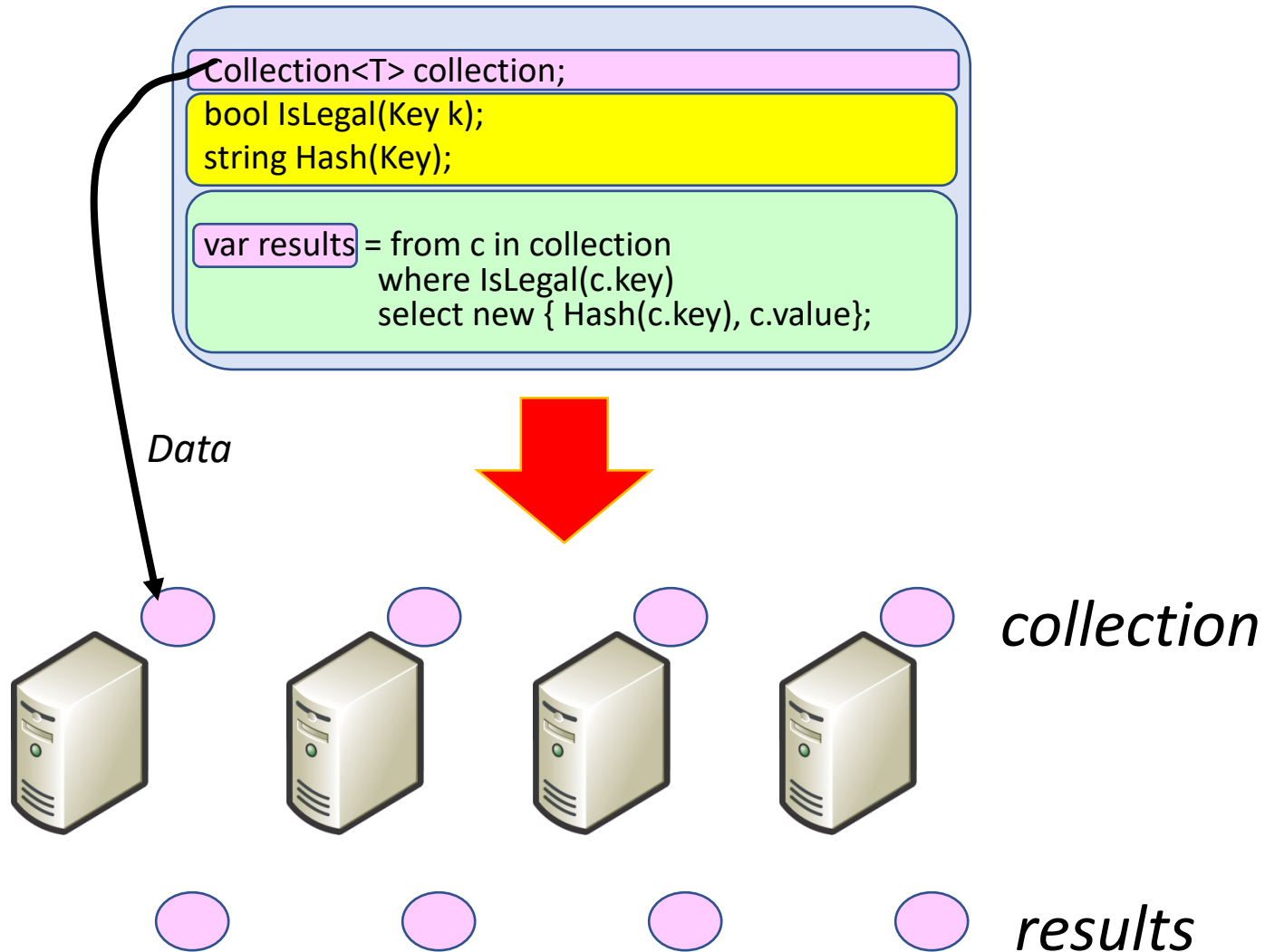
```
Collection<T> collection;
```

```
bool IsLegal(Key k);  
string Hash(Key);
```

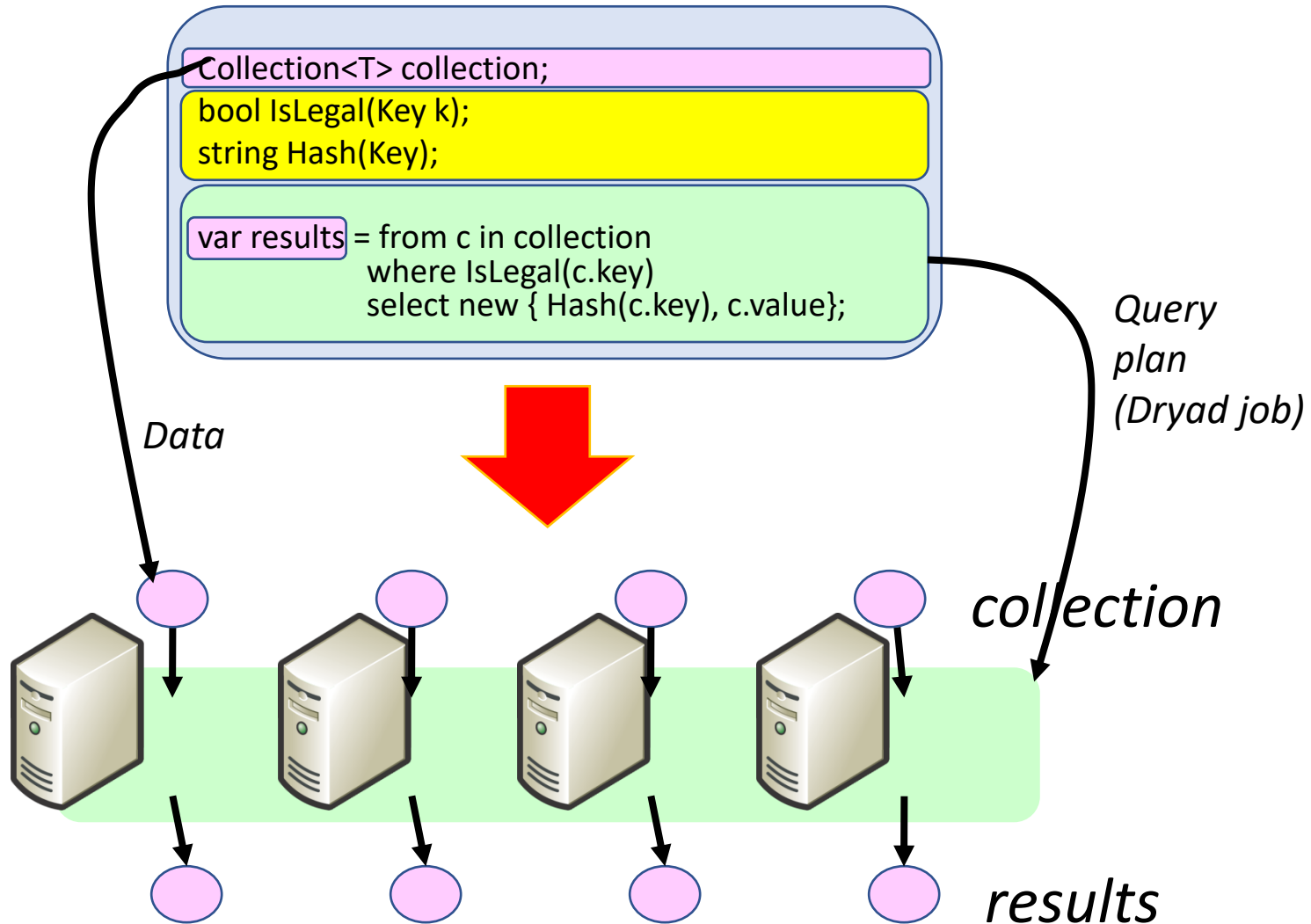
```
var results = from c in collection  
              where IsLegal(c.key)  
              select new { Hash(c.key), c.value};
```



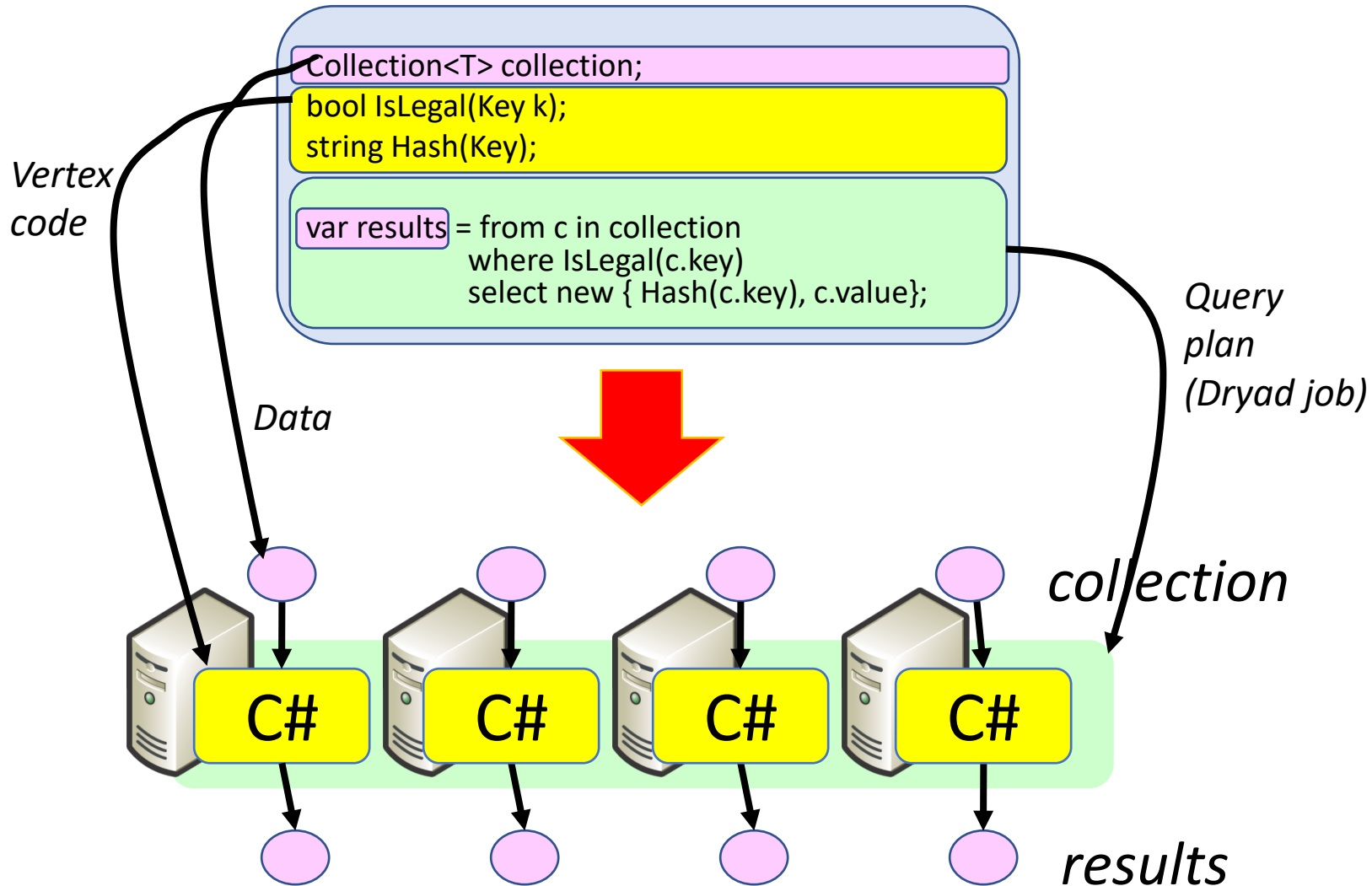
DryadLINQ = LINQ + Dryad



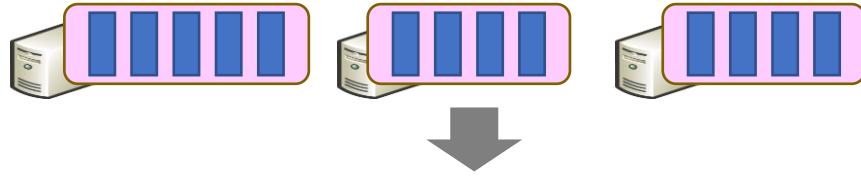
DryadLINQ = LINQ + Dryad



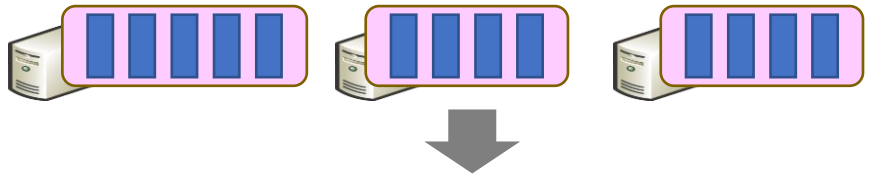
DryadLINQ = LINQ + Dryad



Programming Model

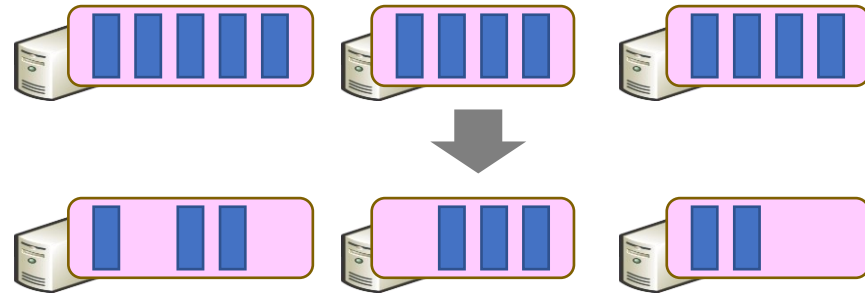


Programming Model



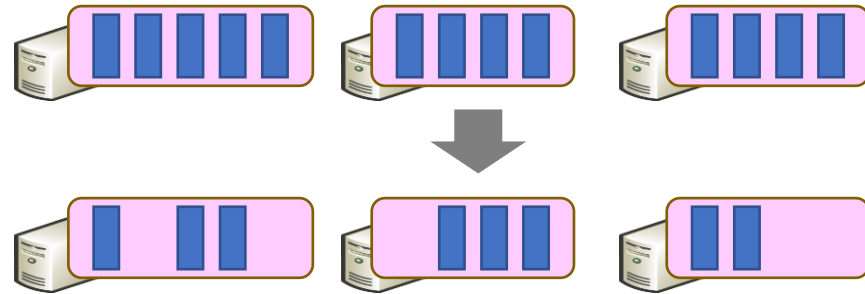
Where

Programming Model



Where

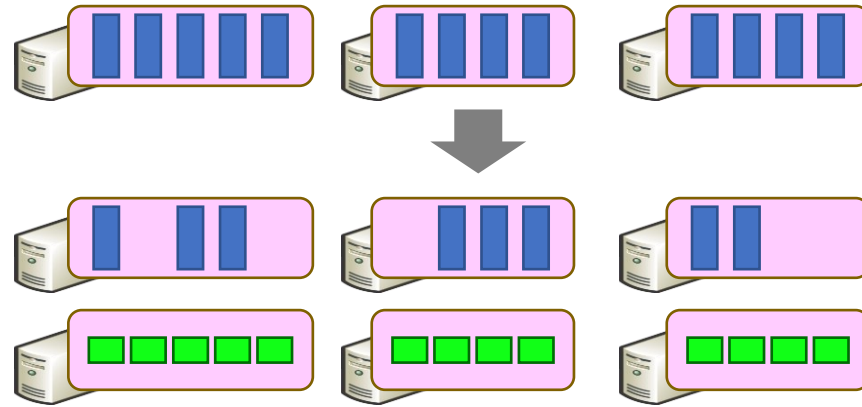
Programming Model



Where
Select

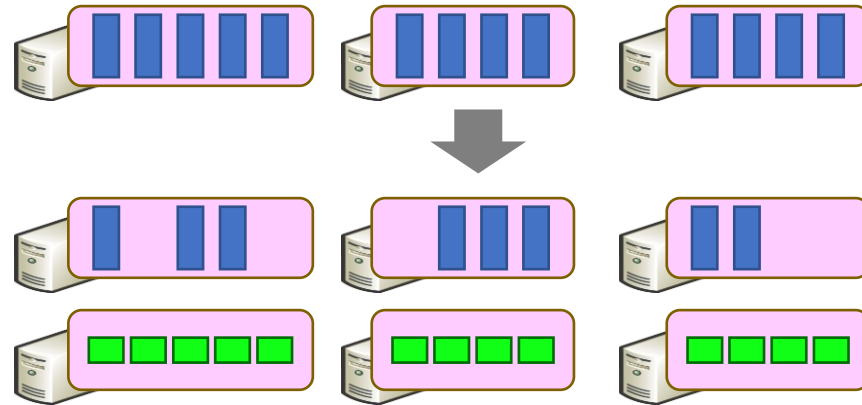
Programming Model

Where
Select



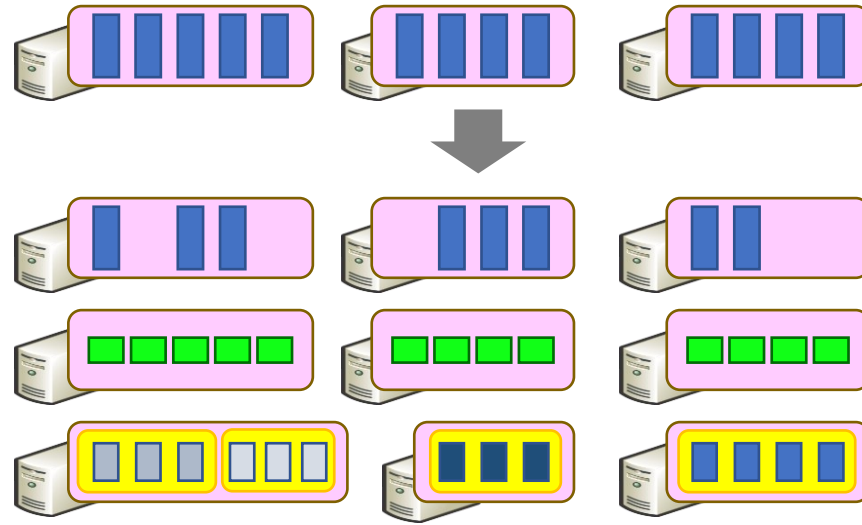
Programming Model

Where
Select
GroupBy



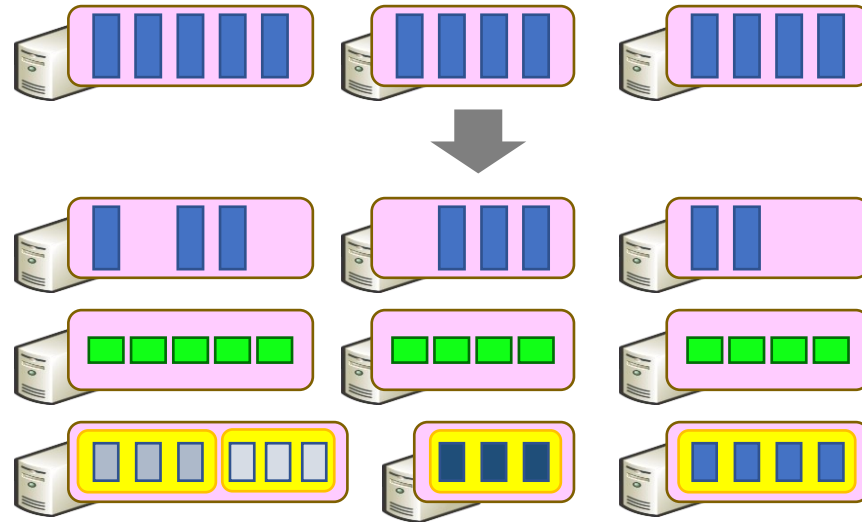
Programming Model

Where
Select
GroupBy



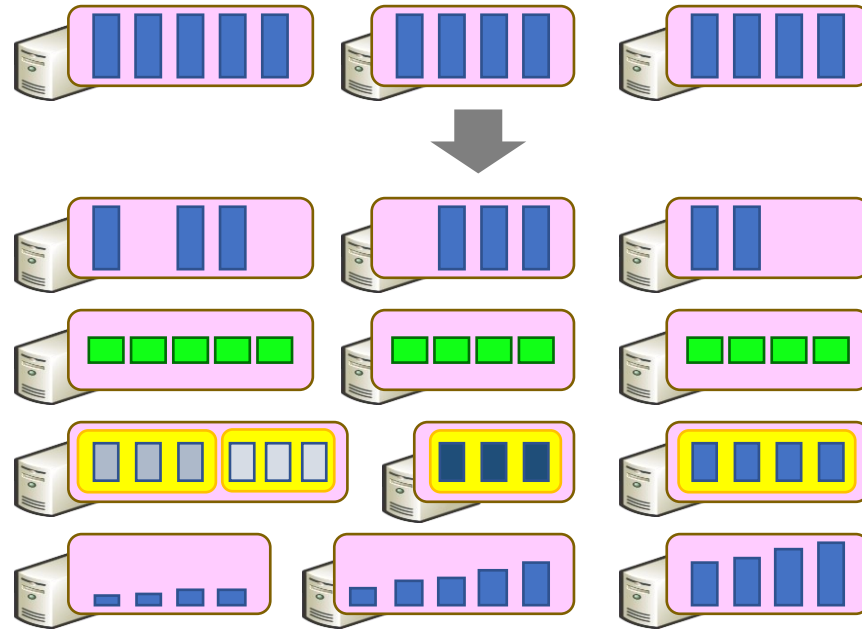
Programming Model

Where
Select
GroupBy
OrderBy



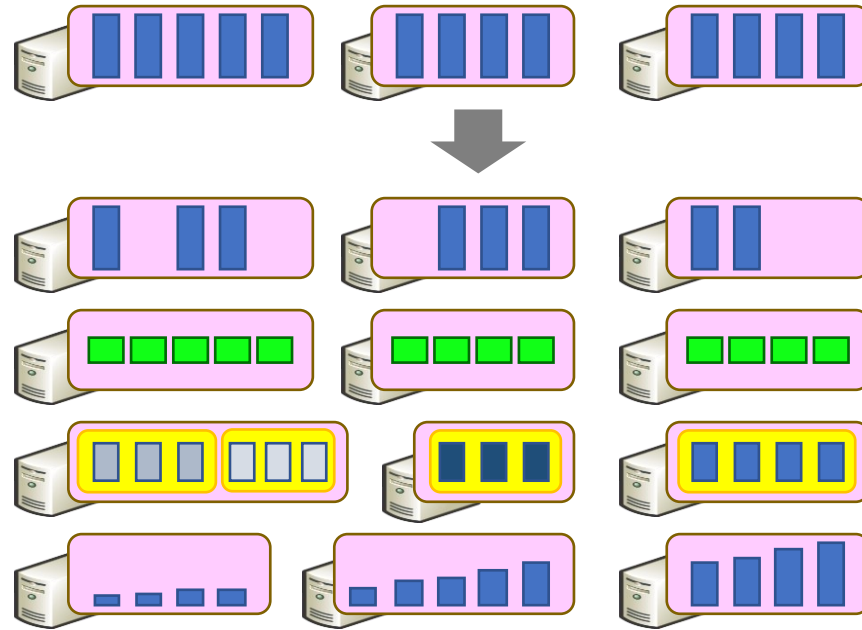
Programming Model

Where
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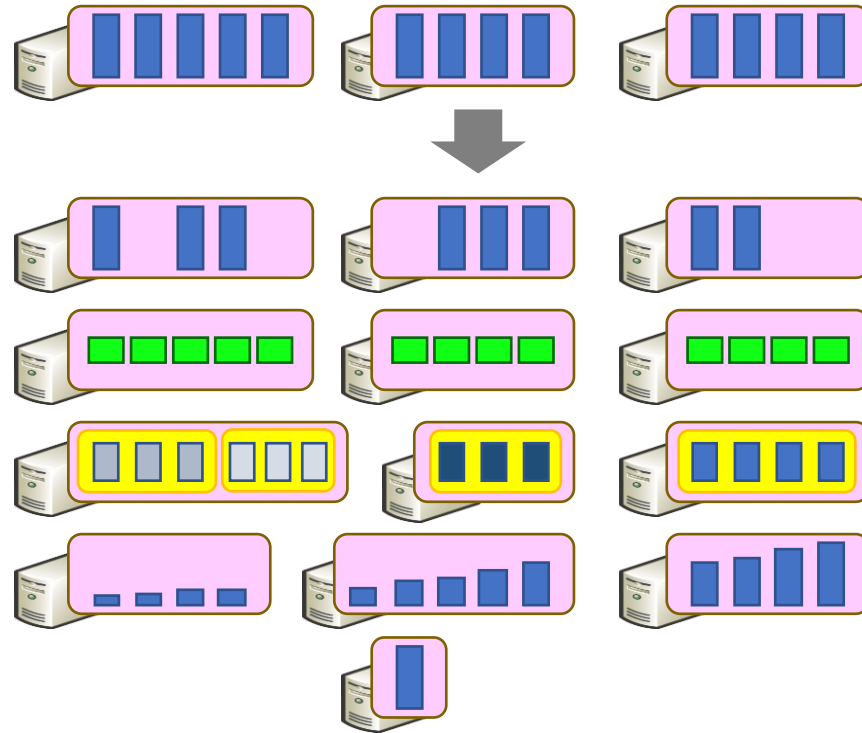
Programming Model

Where
Select
GroupBy
OrderBy
Aggregate



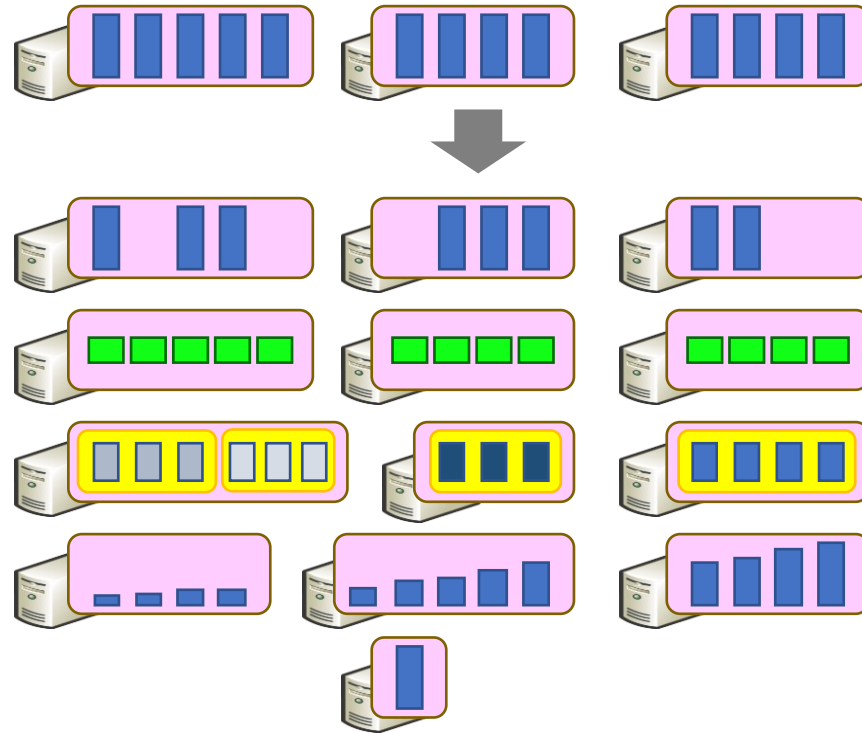
Programming Model

Where
Select
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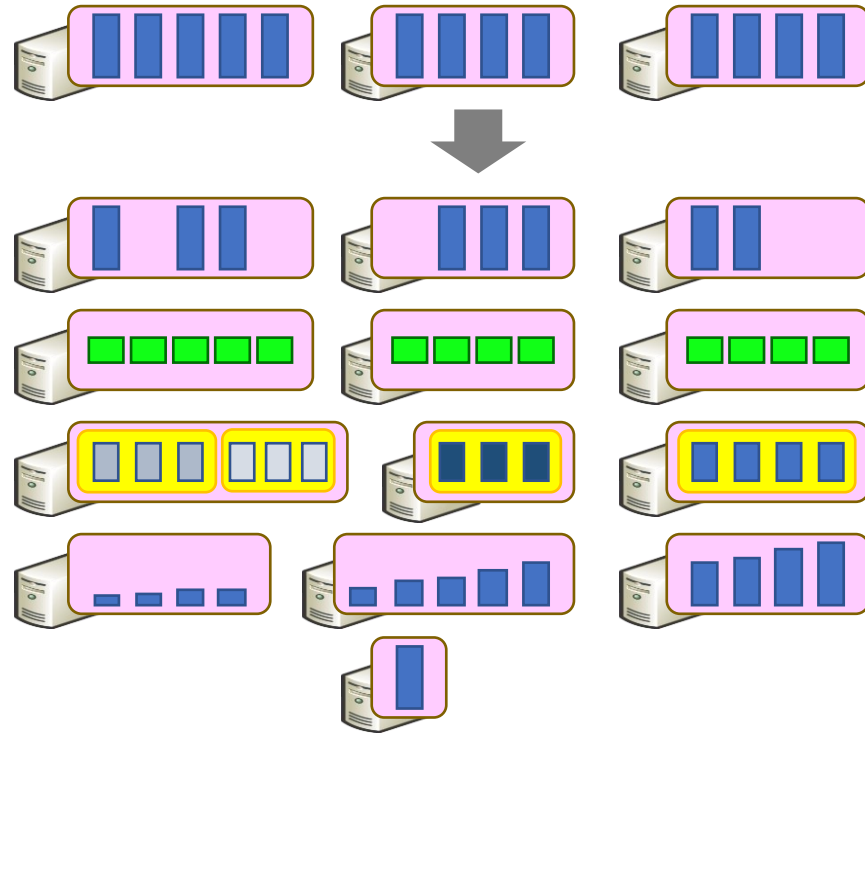
Programming Model

Where
Select
GroupBy
OrderBy
Aggregate
Join



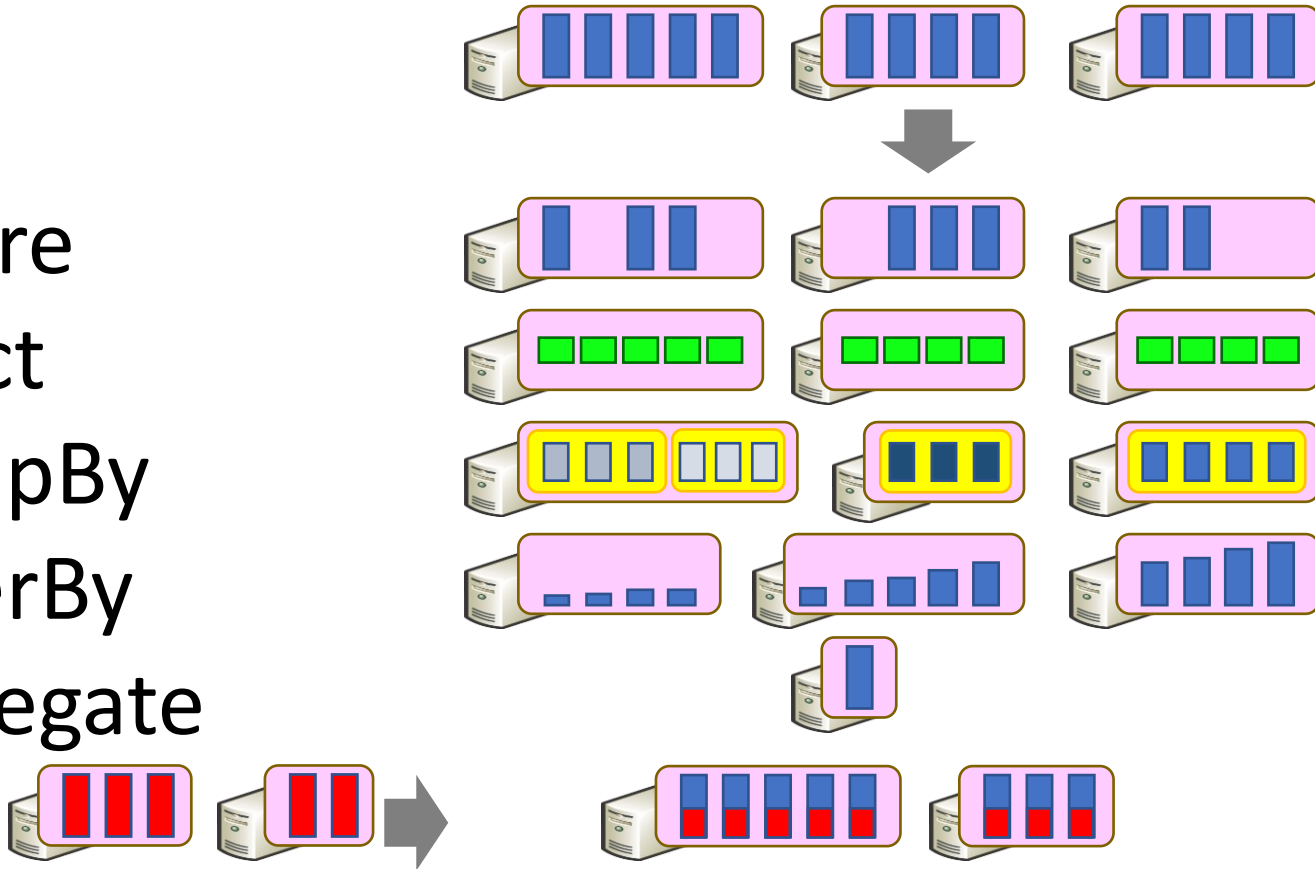
Programming Model

Where
Select
GroupBy
OrderBy
Aggregate
Join



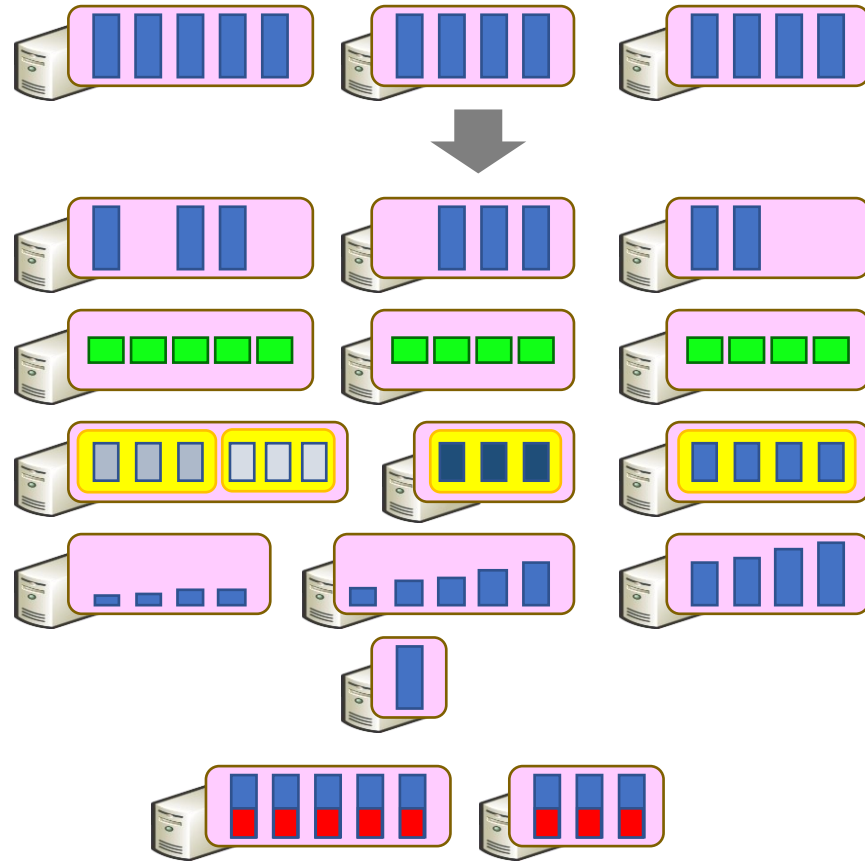
Programming Model

Where
Select
GroupBy
OrderBy
Aggregate
Join



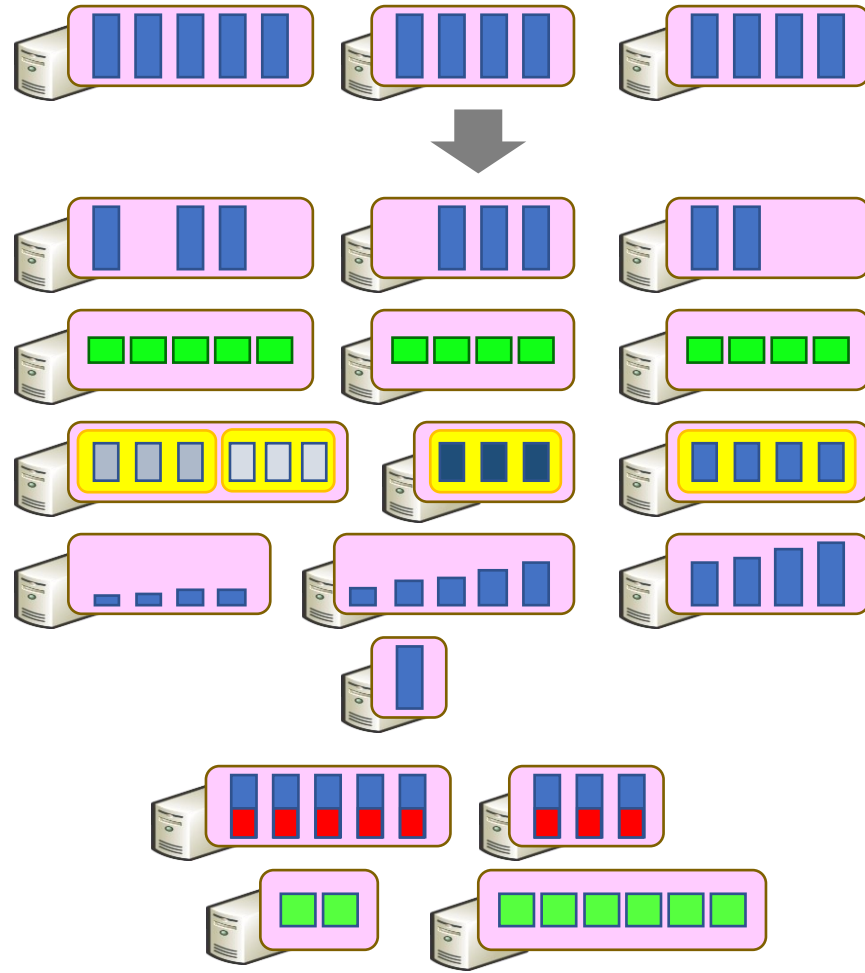
Programming Model

Where
Select
GroupBy
OrderBy
Aggregate
Join
Apply



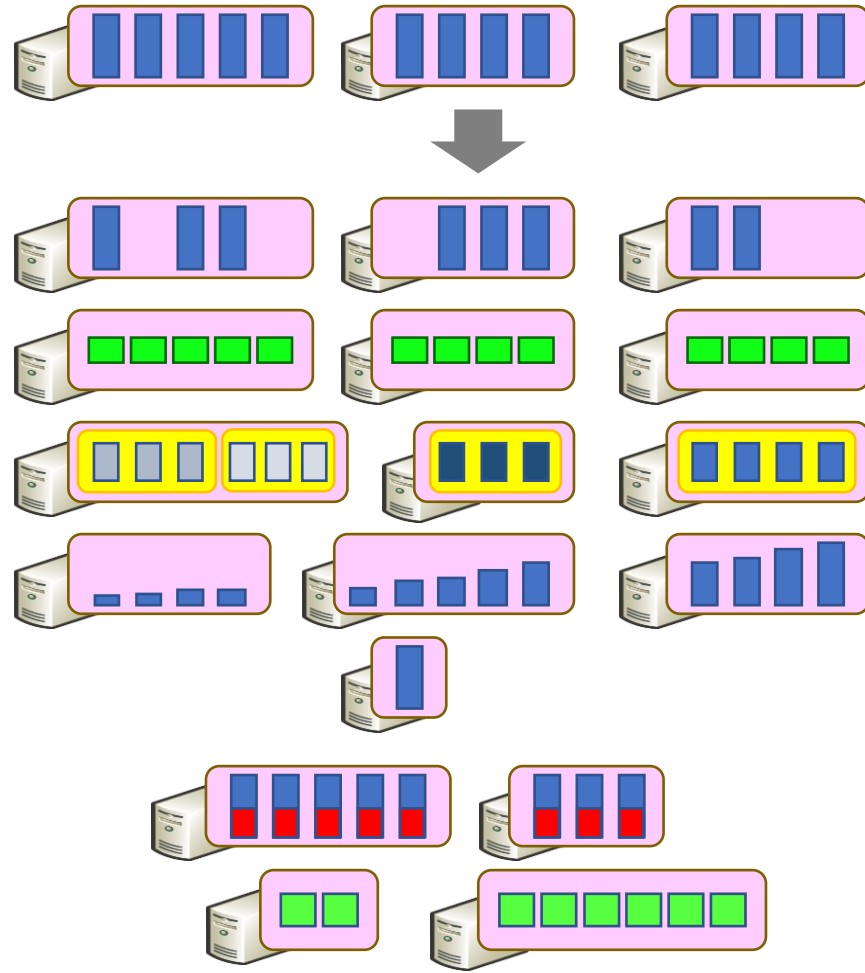
Programming Model

Where
Select
GroupBy
OrderBy
Aggregate
Join
Apply



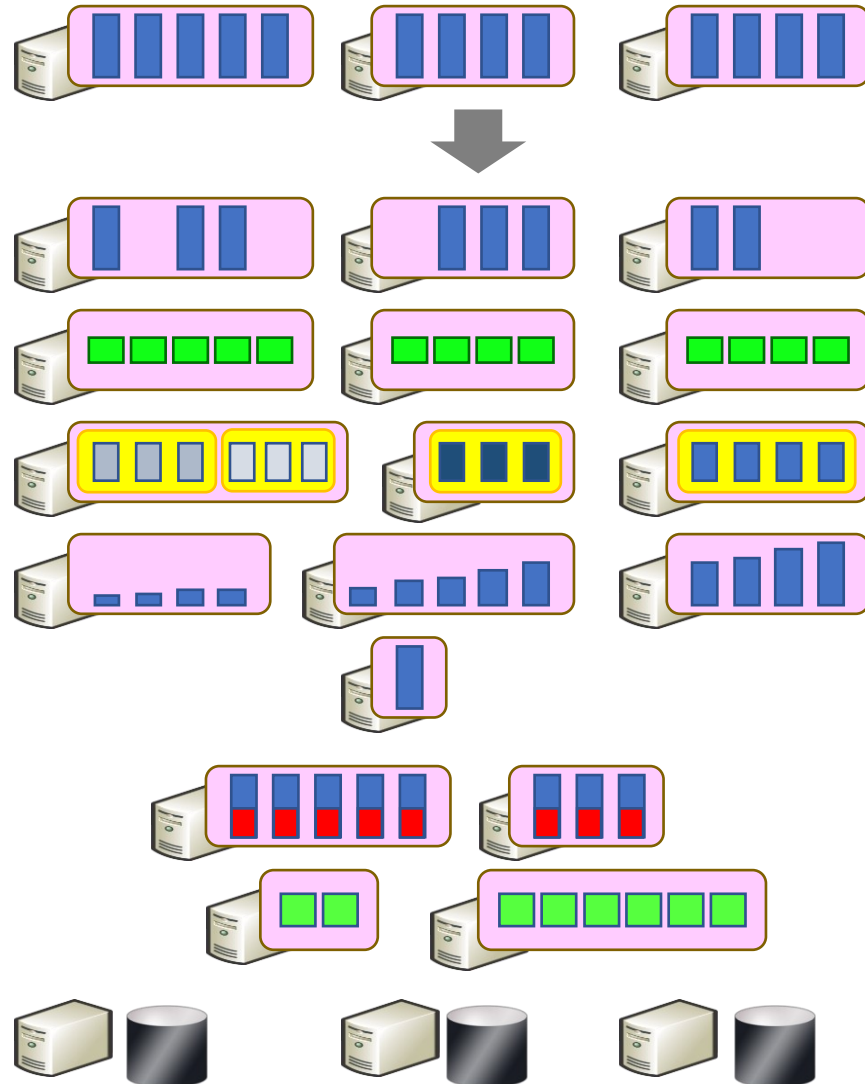
Programming Model

Where
Select
GroupBy
OrderBy
Aggregate
Join
Apply
Materialize



Programming Model

Where
Select
GroupBy
OrderBy
Aggregate
Join
Apply
Materialize



Example: Histogram

```
public static IQueryable<Pair> Histogram(
    IQueryable<LineRecord> input, int k)
{
    var words = input.SelectMany(x => x.line.Split(' '));
    var groups = words.GroupBy(x => x);
    var counts = groups.Select(x => new Pair(x.Key, x.Count()));
    var ordered = counts.OrderByDescending(x => x.count);
    var top = ordered.Take(k);
    return top;
}
```

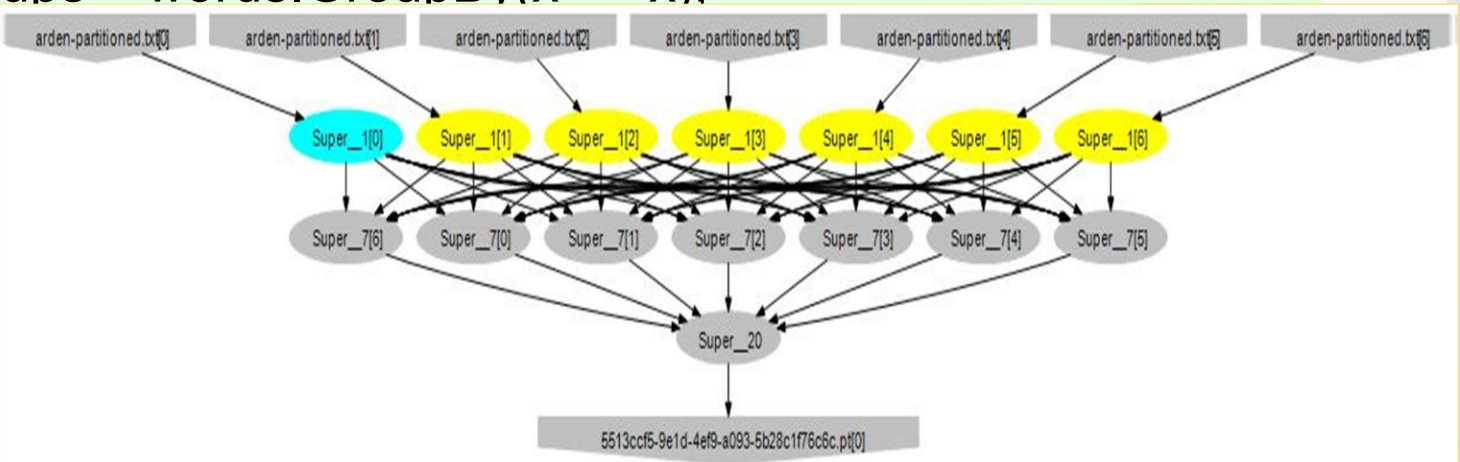
"A line of words of wisdom"
["A", "line", "of", "words", "of", "wisdom"]
[["A"], ["line"], ["of", "of"], ["words"], ["wisdom"]]
[{"A", 1}, {"line", 1}, {"of", 2}, {"words", 1}, {"wisdom", 1}]
[{"of", 2}, {"A", 1}, {"line", 1}, {"words", 1}, {"wisdom", 1}]
[{"of", 2}, {"A", 1}, {"line", 1}]

Example: Histogram

```
public static IQueryable<Pair> Histogram(
    IQueryable<LineRecord> input, int k)
{
    var words = input.SelectMany(x => x.line.Split(' '));
    var groups = words.GroupBy(x => x);

```

- SelectMany
- Sort
- GroupBy+Select
- HashDistribute
- MergeSort
- GroupBy
- Select
- Sort
- Take
- MergeSort
- Take



[A , line , of , words , of , wisdom]
[["A"], ["line"], ["of", "of"], ["words"], ["wisdom"]]
[{ "A", 1 }, { "line", 1 }, { "of", 2 }, { "words", 1 }, { "wisdom", 1 }]
[{ "of", 2 }, { "A", 1 }, { "line", 1 }, { "words", 1 }, { "wisdom", 1 }]
[{ "of", 2 }, { "A", 1 }, { "line", 1 }]

RDDs

- Immutable, partitioned, logical collection of records
 - Need not be materialized
 - contains information to rebuild a dataset
- Partitioning can be based on a key
- Built using bulk transformations on other RDDs
- Can be cached for future reuse

RDDs

- Immutable, partitioned, logical collection of records
 - Need not be materialized

Transformations
(define a new RDD)

build a dataset

on a key

operations on other RDDs

reuse

map
filter
sample
union
groupByKey
reduceByKey
join
persist/cache

...

RDDs

- Immutable, partitioned, logical collection of records
 - Need not be materialized

Transformations
(define a new RDD)

map
filter
sample
union
groupByKey
reduceByKey
join
persist/cache
...

Parallel operations
(return a result to driver)

reduce
collect
count
save
lookupKey
...

RDDs

- Immutable, partitioned, logical collection of records
 - Need not be materialized

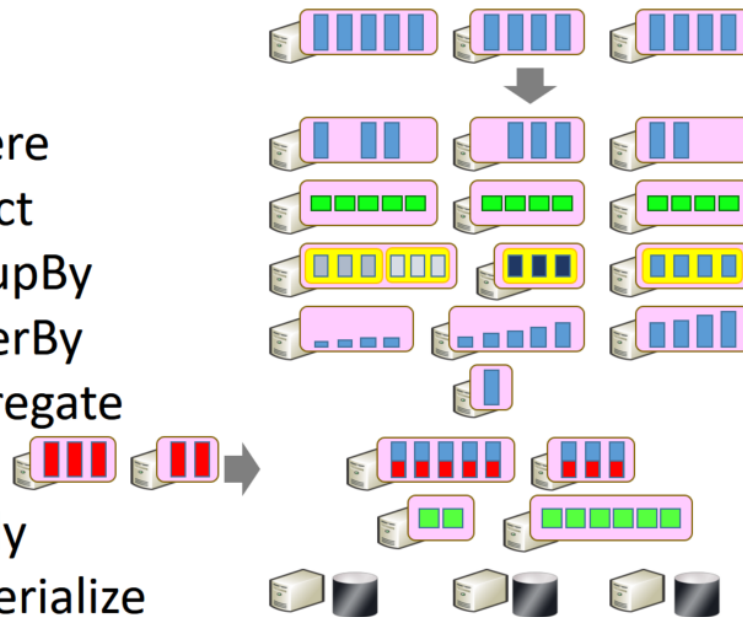
Transformations
(define a new RDD)

map
filter
sample
union
groupByKey
reduceByKey
join
persist/cache
...

Parallel operations
(return a result to driver)

reduce

Where
Select
GroupBy
OrderBy
Aggregate
Join
Apply
Materialize



RDDs vs Distributed Shared Memory

Concern	RDDs	Distr. Shared Mem.
Reads	Fine-grained	Fine-grained
Writes	Bulk transformations	Fine-grained
Consistency	Trivial (immutable)	Up to app / runtime
Fault recovery	Fine-grained and low-overhead using lineage	Requires checkpoints and program rollback
Straggler mitigation	Possible using speculative execution	Difficult
Work placement	Automatic based on data locality	Up to app (but runtime aims for transparency)

Review: What is GroupBy?

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Group a collection by key

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Group a collection by key

Lambda function maps elements → key

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var res = ints.GroupBy(x => x);
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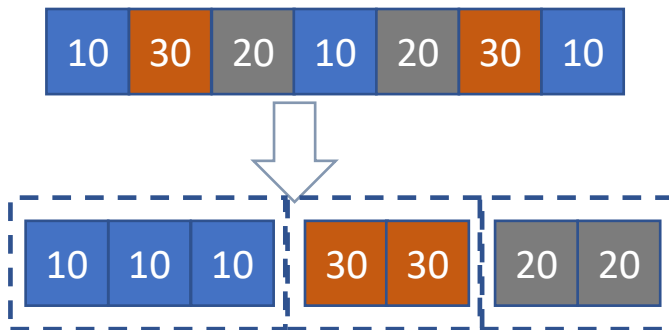
10	30	20	10	20	30	10
----	----	----	----	----	----	----

Review: What is GroupBy?

Group a collection by key

Lambda function maps elements → key

```
var res = ints.GroupBy(x => x);
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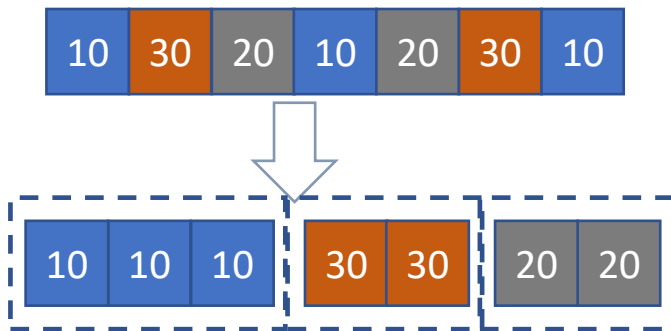


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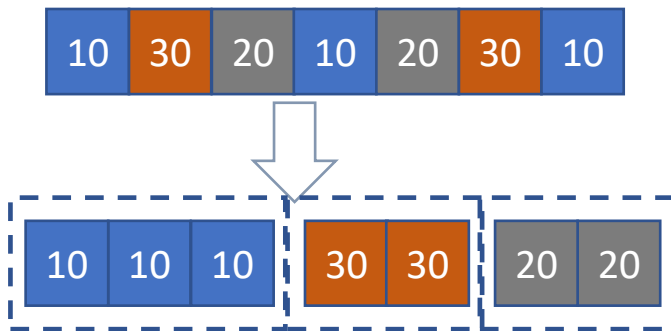
```
foreach(T elem in ints)
{
    key    = KeyLambda(elem);
    group = GetGroup(key);
    group.Add(elem);
}
```



Review: What is GroupBy?

Group a collection by key

Lambda function maps elements \rightarrow key

```
var res = ints.GroupBy(x => x);
```



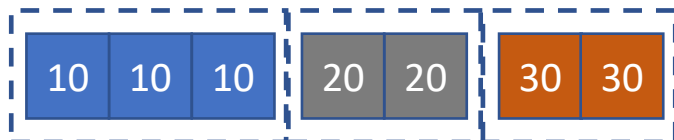
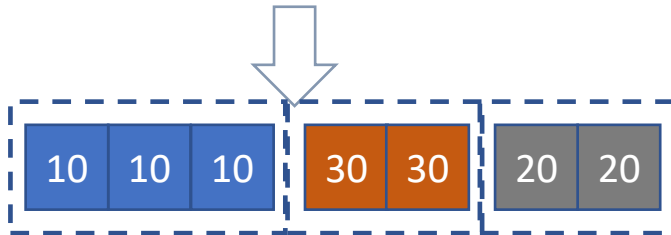
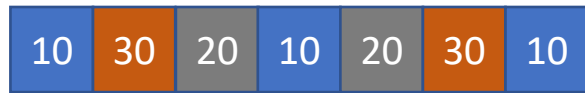
```
foreach(T elem in PF(ints))  
{  
    key = KeyLambda(elem);  
    group = GetGroup(key)   
    group.Add(elem);   
}
```



Review: What is GroupBy?

Group a collection by key

Lambda function maps elements \rightarrow key

```
var res = ints.GroupBy(x => x);
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```
foreach(T elem in PF(ints))  
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Why is MapReduce backwards?

Map == group-by

Reduce == aggregate

Why is MapReduce backwards?

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```
SELECT job, COUNT(*) as "numemps"  
FROM employees  
WHERE salary > 1000  
GROUP BY job;
```

Why is MapReduce backwards?

Map == group-by

Reduce == aggregate

```
SELECT job, COUNT(*) as "numemps"  
FROM employees  
WHERE salary > 1000  
GROUP BY job;
```

- Where is the aggregate in this example?
- Is the DBMS analogy clear?

Why is MapReduce backwards?

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Schemas are good (what's a schema?)

Why is MapReduce backwards?

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Separation of schema from app is good (why?)

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High-level access languages are good (why?)

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