## Topic 24 Heaps

"You think you know when you can **learn**, are more sure when you can **write**, even more when you can **teach**, but certain when you can **program**."

- Alan Perlis



## **Priority Queue**

- Recall priority queue
  - elements enqueued based on priority
  - dequeue removes the highest priority item
- Options?
  - List? Binary Search Tree? Clicker 1

Linked List enqueue BST enqueue

A. O(N) O(1)

B. O(N) O(log N)

C. O(N) O(N)

 $\mathsf{D.} \ \ \mathsf{O}(\mathsf{log}\mathsf{N}) \qquad \qquad \mathsf{O}(\mathsf{log}\mathsf{N})$ 

E. O(1) O(logN)

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# **Another Option**

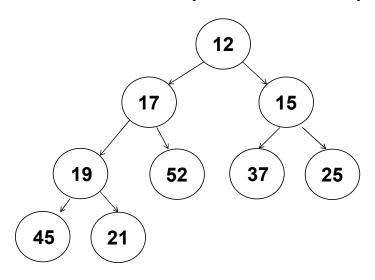
- The *heap* data structure
  - not to be confused with the runtime heap (portion of memory for dynamically allocated variables)
- A complete binary tree
  - all levels have maximum number of nodes except deepest where nodes are filled in from left to right
- Maintains the *heap order property* 
  - in a min heap the value in the root of any subtree is less than or equal to all other values in the subtree

#### Clicker 2

- In a max heap with no duplicates where is the largest value?
- A. the root of the tree
- B. in the left-most node
- C. in the right-most node
- D. a node in the lowest level
- E. none of these

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# **Example Min Heap**



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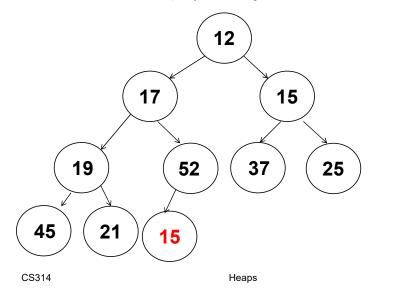
## **Add Operation**

- Add new element to next open spot in array
- Swap with parent if new value is less than parent
- Continue back up the tree as long as the new value is less than new parent node

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# Add Example

Add 15 to heap (initially next left most node)

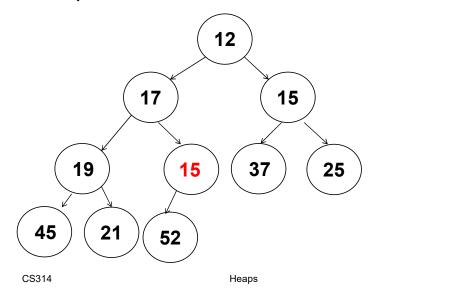


## Add Example

Swap 15 and 52

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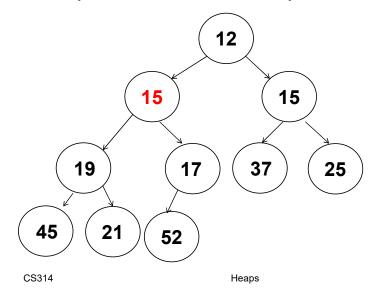
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# **Enqueue Example**

Swap 15 and 17, then stop



### Add Example

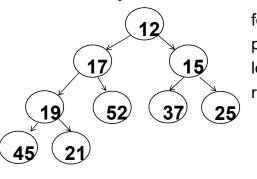
Insert the following values 1 at a time into a min heap:

16 9 5 8 13 8 8 5 5 19 27 9 3

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# Internal Storage

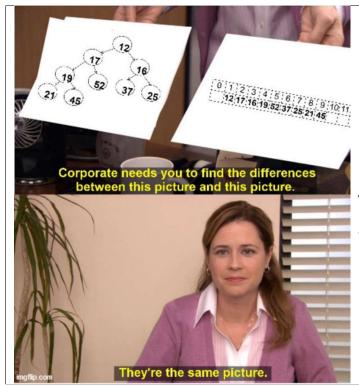
Interestingly heaps are often implemented with an array instead of nodes



for element at index i: parent index: i / 2left child index: i \* 2right child index: i \* 2 + 1 9

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
	12	17	15	19	<b>52</b>	37	25	45	21						

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In Honor of Elijah, The Meme King, Spring 2020

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## **PriorityQueue Class**

```
public class PriorityQueue<E extends Comparable<? super E>>
{
    private int size;
    private E[] con;

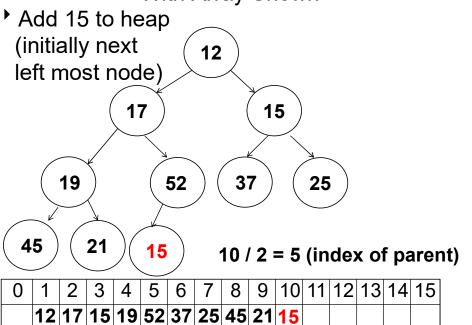
public PriorityQueue() {
        con = getArray(2);
    }

    private E[] getArray(int size) {
        return (E[]) (new Comparable[size]);
    }

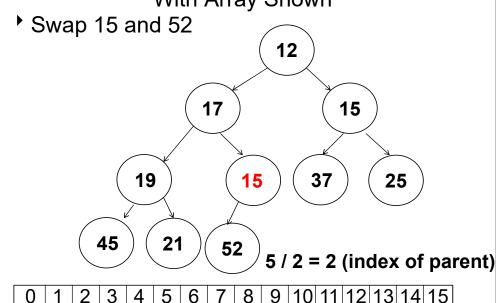
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```

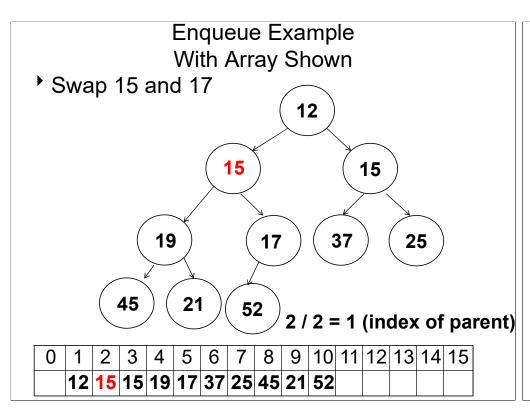
## PriorityQueue enqueue / add

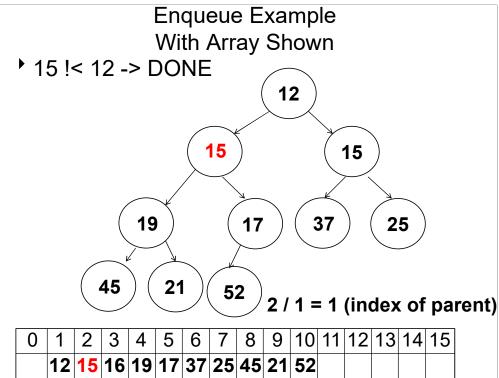
# Enqueue / add Example With Array Shown

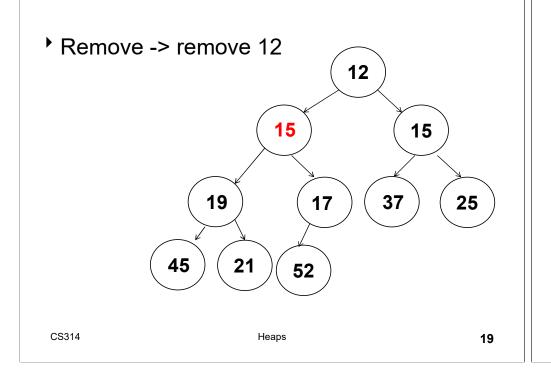


# Enqueue Example With Array Shown





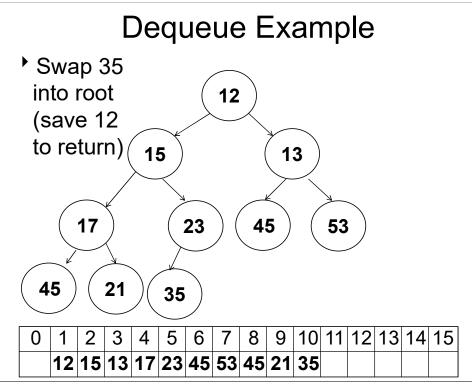


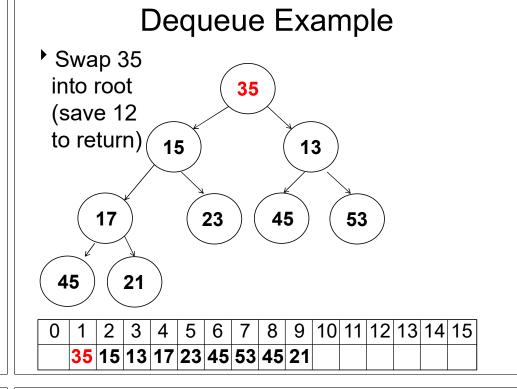


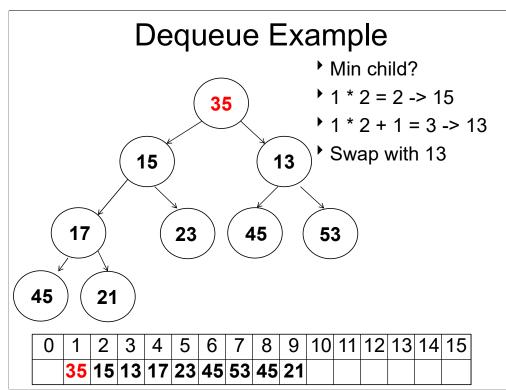
## Remove / Dequeue

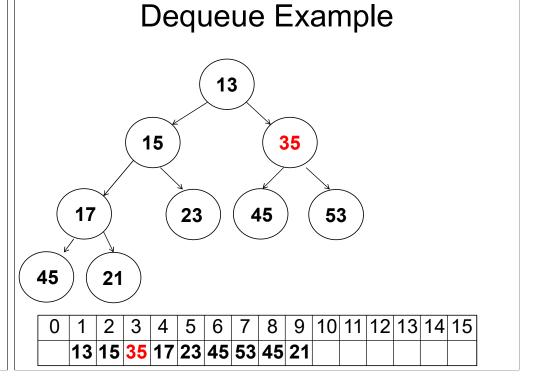
- min value / front of queue is in root of tree
- swap value from last node to root and move down swapping with smaller child unless values is smaller than both children

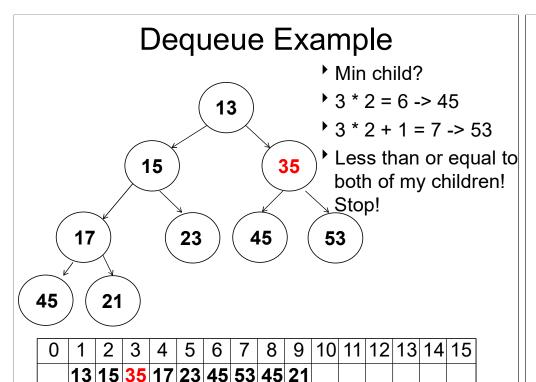
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## Clicker 3 - PriorityQueue Comparison

- Run a Stress test of PQ implemented with Heap and PQ implemented with BinarySearchTree
- What will result be?
- A. Heap takes half the time or less of BST
- B. Heap faster, but not twice as fast
- C. About the same
- D. BST faster, but not twice as fast
- E. BST takes half the time or less of Heap

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## Dequeue Code

```
public E dequeue() {
    E top = con[1];
    int hole = 1;
   boolean done = false;
   while ( hole * 2 < size && ! done ) {
        int child = hole * 2;
        // see which child is smaller
        if (con[child].compareTo(con[child + 1]) > 0)
            child++;
                        // child now points to smaller
        // is replacement value bigger than child?
        if (con[size].compareTo( con[child] ) > 0 ) {
            con[hole] = con[child];
            hole = child;
        else
            done = true;
    con[hole] = con[size];
    size--;
    return top;
```

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